Colorcue


Aug/Sept 1981

## Colorcue

## A bi-monthly publication by and for Intecolor and Compucolor Users

## Editors:

## August/September 1981

Ben Barlow
Volume 4 , Number 1

## 3 Editors' Notes

3 Compucolor Policy
5 The Serial Port, by Ben Barlow How to use the RS232 port.

13 Lower Case y's, by Tom Devlin A simple, cheap fix for a nagging problem

14 Screen Dump to MX-80, by Mark Fairbrother Subroutine to move graphic image from screen to printer

17 User Group Bulletin Board
19 Assembly Language Programming, by David B. Suits Part 1 of a series for the beginner

Advertisers: You will find our advertising policies attractive! Write for details.

Authors: This is a user-oriented and supported publication. Your articles/tips/hints are required to make it go. Send your articles or write for information.

Colorcue is published bi-monthly by Intelligent Systems Corporation, with editorial offices in Rochester, New York. Editorial and subscription correspondence should be addressed to the Editors, Colorcue, 161 Brookside Drive, Rochester, NY 14618. Product related correspondence should be addressed to ISC, 225 Technology Park, Norcross, GA 30092, ATTN: Susan Sheridan. Reproduction in whole or part without permission is discouraged. Opinions expressed in by-line articles are not necessarily those of the editors or of ISC. Hardware/software items are checked to the best of our abilities, but are NOT guaranteed.

## Editors' Notes

The editorial duties of Colorcue have been assumed by a new editor. In fact, there are two of us. Since we plan to be around for a while, let us introduce ourselves.

David B. Suits obtained his PhD. in philosophy from the University of Waterloo in Ontario, Canada, and teaches philosophy at the Rochester Institute of Technology, He has published articles in both philosophy and microcomputer journals and has written a book about Intecolor/Compucolor graphics. Dave became interested in computers when he discovered that his programmable calculator was too small to handle a chess-playing program. He's been the librarian of the Rochester Compucolor Users Group for $21 / 2$ years and is interested in artificial intelligence, music, and hardware.

Ben Barlow is a computer hacker. A systems programer for many years, he is now involved with computers only as a hobby. He works for the Kerox Corporation in Rochester, and for the past 2 $1 / 2$ years he has edited and published Datachip, the Rochester Compucolor Users Group newsletter. His interests lie in the area of hardware and external world interfaces.

A word about policies is in order, especially in light of the recent shifts of Colorcue direction. Colorcue will continue to be a bi-monthly (every 2 months) publication directed toward both new and experienced users of Intecolor and Compucolor compu-
ters. We are committed to its timely publication and to a goal of providing top-quality, accurate and useful information to our readers. At the same time, however, we're dependent upon your participation. We need your articles, programs and ideas to make this magazine into what it should be and what we're sure you want it to be. Let us and your fellow users know what you've been doing with your computer.

A11 current subscriptions are good now for one year (six issues), as are the applications of those of you who responded to the query about an Intecolor publication. No renewal worries for anybody until next year! So, sit back, read, enjoy. (And, we trust, learn.)

Dave and Ben

## Compucolor Policy

Because of a ruling by the Federal Communications Commission, Intelligent Systems Corporation is no longer permitted to sell the Compucolor II personal computer in the United States. The ruling stems from the FCC's determination that most personal computers generate substantial interference--resulting in impaired television reception. You may have noticed that when you work with the Compucolor II, a television viewer in the next room can have troubles getting a good picture. This problem is especially troublesome to apartment resi-
dents, who have no control over the activities of neighbors in adjoining rooms. Since televisions and personal computers are communication media, the FCC stepped in. Vacuum cleaners, hair dryers, and other appliances are not affected by the ruling, even though they are often much more annoying to the television viewer, because they are not communications equipment.

But personal computer owners need not be worried that they are operating machines illegally. The ruling applies only to machines manufactured after the FCC deadline, which, in the case of Intelligent Systems, was April, 1981. Al1 Compucolor IIs already in service are perfectly accept-able--no modifications are necessary.

Virtually all personal computer manufacturers were affected by the FCC's decision. Companies such as Apple and Radio Shack were faced with modifying their products in order to meet the new standards. Intelligent Systems started to reseach the changes in design, components, and production that would have to be made in order to meet the new standards. Our financial planners and marketing people assessed the overall costs of such changes and weighed them against current profits and projected profits. Then we investigated all possible alternatives to come up with a final decision about how we would comply with the ruling.

In the course of studying the problem, it became evident that re-designing the Compucolor II in order to comply with the FCC standards was simply too expensive. It would mean further channeling of resources away from
the company's mainstay line or business and scientific computers, and that was not feasible. Consequently, Intelligent Systems Corporation had no choice but to discontinue manufacture of the Compucolor II. However, since the FCC ruling does not affect overseas clients, the company still manufactures Compucolor IIs for sale in foreign markets.

Once the decision was made to discontinue the Compucolor II for domestic markets, company leaders turned their attention to Compucolor II owners in order to determine how they might best be served. A two-point plan was developed that serves past customers well beyond any legal or moral responsibility. Intelligent Systems Corporation provides:
(1) A source of information -- Colorcue. Despite increasing costs of postage, editorial time, and materials, Colorcue will continue to be published. This newsletter supplies users with valuable information such as names of people who welcome contact from other users, and user group locations for those who wish to meet with other owners. Intelligent Systems decided to sponsor publication of the newsletter so that all users will have a place to turn for information.
(2) A source of parts and repair. Intelligent Systems continues to maintain its repair facility at the factory. Replacement parts and trained service personnel are available to give satisfactory turn-around time at a reasonable cost. This ensures that all owners will be able to maintain their units properly for best performance. E

## The Serial Port

by Ben Barlow

Stop! You apply the brakes as the traffic light changes first to yellow and then to red. You are obeying the rules associated with a standard established back in the 1930's just as your computer's RS232 port obeys the RS232 standard as it sends and receives data. The traffic light's purpose is to control traffic flow; the RS232's purpose is to control data flow. Just as the traffic light does not care what kind of traffic is stopping below it, the RS232 standard specifies rules for sending and receiving data but does not restrict the data being sent. The only requirement is that data characters be transmitted in a serial (bit by bit) manner, rather than in parallel (all bits simultaneously). RS232 is not concerned with the type of code, parity, or the number of stop bits, synchronous or asynchronous transmission - only with the signalling levels to be used to transmit data and the protocol used to interchange it.

Why have a standard, anyway? The answer is obvious, just as it is at the intersection of two streets. Some commonly understood method of control is required to make sure control is orderly, and to ensure compatible products. A good standard serves as the base for manufacturers to develop products, and promotes healthy competition (read lower price). Look at the host of different terminal and computer devices on the market manufactured by a myriad of manufacturers whose products are compatible because of the RS232 standard.

The RS232 standard (current revision level C adopted in 1969) governs the connection of terminals and computing equipment to modems. Although it's not intended to describe direct terminal to computer interfacing, it's been used that way often, especially since the onset of personal computing. The standard defines three aspects of the interface - 1)the electrical standard (which is what most advertising refers to) 2) the mechanical standard, and 3) the interchange, or protocol standard. The latter is frequently determined by the software in the respective computers and terminals. Let's look at these three aspects separately.

The easiest one to understand is the mechanical interface. The ubiquitous DB25 connector has become almost synonymous with RS232C. This familiar 25 pin plug (male) and socket (female) combination appears on the end of every cable and is built into most equipment that meets the
standard. Some of the ISC computers use an edge connector to provide the RS232 interface, which must have an added cable to give a DB25-type connection. By convention, terminals and modems usually have female connectors, which leads to male/male cables, but this is not universally true.


## DB 25 Connectors

Fig. 1

Electrical characteristics are also easy to understand. RS232 signals are based on specified voltage levels. Your computer is constructed largely of TTL family devices (integrated circuits). This family also uses voltage levels for signal determination. TTL levels range between 0 and 5 volts; logic 0 is recognized as being a voltage between 0 and .4 volts, and logic 1 is recognized as being between 2.4 and 5.0 volts. A voltage between .4 V and 2.4 V is undefined. This is a relatively small difference in voltage, and can easily be affected by electrical surges caused by motors turning on (refrigerators, copiers, vacuum cleaners), or even by line capacitance if the TTL signal must be passed over a long (measured in feet) distance. Thus, it's unacceptable as a data communication method over any useful distance.

The RS232C standard, on the other hand, specifies voltages ranging from -25 V to +25 V and provides much greater immunity to noise. A positive signal is a voltage between 3 V and 25 V , and a negative signal is a voltage between -25 V and -3 V . The area between -3 V and +3 V is undefined. ISC computers use plus 9 V and minus 9 V as the signalling levels.

The interchange standards are the most difficult to describe because of the many options which are possible. The simplest is the send-only interface, sometimes known as "send and pray". A simple send/receive interface can be constructed relatively easily, and a send/receive (or simply send) interface with appropriate handshake can also be constructed without much trouble. The standard includes many features such as
signalling on a secondary channel，ring detection（for dialup connections）， and data rate selection．These are fairly uncommon，though，so we＇ll not discuss their use．Let＇s look instead at the signals used in common interfaces that you may apply to your computer，and then actually put them to use．

Pin 非1（FG）is a common Frame Ground electrically connecting the frames of the two machines being connected．It can sometimes be omitted，especially if，as in ISC machines，frame ground and signal ground（see later）are connected．When connecting to another piece of equipment（e．g．modem or printer），check the outlets into which each is plugged to make sure the polarity is identical．If the electrician made a mistake，it could be a costly one for your machine．Sears sells an inexpensive tester for this purpose．（Cat．非 49662）

Pin 非2（TD）is for Transmitted Data，meaning data going out of the interface（or，in the case of a modem，into the interface）．The transmitted data circuit at one end must be connected to the received data circuit at the other．This reversal，sometimes accomplished by the modem and sometimes by the cable，is frequently the cause of initial operational difficulties when connecting gear．

Pin $⿰ ⿰ 三 丨 ⿰ 丨 三 3$（RD）is Received Data，data coming into the interface．
Pin 非7（SG）is Signal Ground and is the reference point for positive and negative signal determination．

Some other signals used for handshaking occur in pairs，such as：
DataSet Ready and Data Terminal Ready（pins 6 （DSR）and 20 （DTR） respectively）．It＇s not difficult to imagine the use of these．They tell each connected machine that the other is up．

Request To Send（pin 4，RTS）and Clear To Send（pin 5，CTS）are a pair of signals used to request and grant permission to transmit data．In the conventional terminal／modem arrangement，CTS is the modem＇s response which enables data flow from the terminal．

These eight signals are the most commonly used signals of the 25 （3
unassigned) defined by the standard.
When a sender can send data faster than the receiver can process it, there needs to be some interchange protocol established so that the sender does not swamp the receiver with data and cause data loss (overrun). Several different conventions have been established over the last decade and one of these is implemented using the RS232 signals. It is used by most of the popular personal computer printers, such as the IDS 440-450-460 lines, the Epson, Heath, Base-2 and others. The method is to lower the CTS (or sometimes DTR) line when the receiver does not want to receive any more data. (The ability to sense this signal was omitted from early Compucolors, and is the subject of an engineering change known as ECNO02137, the "handshake mod").

Now let's take a look at what is required to connect an ISC computer to an Epson MX80 printer and a Cat Novation modem, two of the most popular and useful peripherals to be found.

The first step is the construction of a cable that will connect the RS232 port to the devices. Figure 2 shows the connections needed for each. Notice that the signals for each are different. It would seem at first that two cables will be needed, one for the printer, and one for the modem. With a little cleverness, though, we can make one cable in the form of a Y , with your computer on the bottom and the printer and modem each on its own arm. It'll be necessary to plug and unplug devices, since both should not be powered on and connected to the Y-cable simultaneously. Fancy switching arrangements are an alternative, as are separate cables, but the Y-cable works, and it's cheap. If you follow the construction diagram (Figure 2) you should have no trouble, but some hints are in order:

If you use an IDC-type (insulation displacement connector) at the PC board end, press it solidly and evenly together in a vise using even pressure and care. Don't plan to take it apart again!

If you use a solder-type connector at the PC board end, insulate the pin by slipping a length of heatshrink tubing over the wires before soldering. Any 20-24 gauge 4 conductor stranded wire can be used. Belden 8723 is excellent, but can be obtained only in 500 ft. spools.

Solder-type DB25 connectors are the easiest to use.

Use lengths of heatshrink tubing on the $Y$ connections and on the Y itself.

Label the connectors on the legs as "printer" and "modem".


## "Y"Cable

Fig. 2

The popular and versatile Epson MX80 printer has recently received rave reviews, and its connection to your computer is not difficult, and will serve as an illustration for the use of other printers. The Epson must be ordered with the RS232/current loop interface (Cat. no. 8141). Follow the instructions for installing the interface, and when the case is open, make the following jumper and switch settings on the interface board:

JNOR - cut and remove.
JREV - solder (carefully) a jumper of 28 or 30 gauge wire. SW 1-8 - all off.

The printer is now set up to receive data at 9600 BPS, its maximum speed, and respond to potential overflow conditions by raising the clear to send line to your computer.

If you have a Compucolor with a Rev. 3 or earlier computer/logic board, you'l1 need to add ECN002137. Your dealer can do this for you, or by following the logic in Figure 3, you can make the modification yourself. (If sufficient interest develops among older Compucolorists,


Fig. 3
the editorial staff will develop a step-by-step installation sequence and kit. - ed.) Connect the "printer" connector of the $Y$-cable to the MX80 and turn on the power. There are several ways to direct BASIC output to the RS232 port, to the printer (or modem, when you connect it):

1. Enter BASIC with ESC $M$ rather than ESC E.
2. Execute a PLOT 27,13 statement.
3. Execute a POKE $\mathrm{x}, 14$ where $\mathrm{x}=33265$ (for V6.78 and V8.79) or x $=40886$ (for V5.79).

Return BASIC output to the screen by executing a POKE x,0 statement, where x is as above.

FCS output may also be directed to the RS232 port (handy for printing directories) by entering FCS by the ESC G sequence rather than ESC D. Output can be returned to the screen by typing either:

1. CPU RESET, or
2. executing POKE $y, 0$ where $y=33249$ (V6.78, V8.79) or $y=40884$ (V5.79).

With the MX80 set as described above, no other controls or options need be set (the computer defaults to 9600 BPS). Other printers may require speed and stop bit selection. Since the modem connection requires this, let's take a look now. As with directing output, the speed, stop bits, and half/full duplex modes can be set either with keystroke sequences or by executing BASIC statements.

## SPEED

Keystrokes:


BASIC:
PLOT $27,18, s$ where $s$ is as described above.

## STOP BITS

Keystroke:
A7 ON (one stop bit)
A7 OFF (two stop bits).

BASIC:
PLOT 14 (one stop bit)
PLOT 15 (two stop bits).

## HALF/FULL DUPLEX

Keystrokes:
ESC H (half duplex - echo each character locally
as well transmitting it)
ESC F (full duplex - just send it).

BASIC:
PLOT 27,8 (half)
PLOT 27,6 (ful1).

Connecting the Cat Novation modem is simple. Disconnect the printer and plug the "modem" connector of the Y-cable into the modem. Select speed and stop bits with appropriate commands, direct the output to the RS232 port and dial your favorite timesharing sytem, information utility, or friend. The command sequence to use is:

1. CPU RESET
2. A7 ON or A7 OFF
3. ESC R speed
4. ESC F or H
5. ESC M.

Two of the more popular information utilities to which you may subscribe are:

## MICRONET

Personal Computing Division Compuserve Inc.
5000 Arlington Centre Blvd.
Columbus, Ohio 43220
614-457-8600


## The SOURCE

Telecomputing Corporation of America 1616 Anderson Rd.
McLean, Virginia 22102
703-821-6660

Unfortunately, there may be more to using the port than is indicated in this article. Your printer may cause lower case y's or slashes to print on the screen or printer, and when connected to a remote timesharing system, the distant computer may send control sequences that will cause your disk to spin, your screen to clear, or strange characters to be displayed. These problems can be very annoying, but don't despair--there are cures. The lower case y problem can be easily cured by software or hardware. (See the following article). To solve the remote data/local interpretation problem, a program such as ISC's (or other vendor's) is needed to act as a data filter and translator for the incoming data. (A past issue of Colorcue contained a simple data filter program. For a reprint, send a SASE to our editorial offices requesting it. -ed.) e

## Lower Case y's

by Tom Devlin<br>3809 Airport Rd.<br>waterford, MI 48095

The following fix for the lower case $y$ problem is the simplest, cheapest one I've found to date. It requires the use of the handshake mod (described in the previous article - ed.) and the addition of a single 2 cent diode to the keyboard printed circuit board.

Basically, the fix takes advantage of the fact that the output of UCl which you connect to pin 10 of UEl for the handshake mod also goes to R6 in the keyboard via pin 11 of the keyboard cable. This gives us the CTS signal at TTL levels. Adding a diode from this point to any of the column outputs will cause the CCII to think that several keys are being pressed whenever CTS is low, causing the keyboard scan routine to give up and return without a character for that scan.

Following the placement of the diode as shown in the accompanying photograph, you should have no trouble adding the fix to your machine. Use a 1 N914 diode, and make sure the band, i.e. the cathode, is toward the right, as shown in the photo.
(Note: Although this fix works with no apparent side effects, it has not been sanctioned by ISC, and may affect your warranty.- ed.)

Photo 1.
Trace side of keyboard PCB. Solder a signal diode as shown, cathode end to right.


## Screen Dump to the MX80 Printer

by Mark D. Fairbrother<br>Carriage House East, A5<br>Rt. ll<br>Kirkwood, NY 13795

This program will produce a hardcopy of the 128 by 128 resolution ISC graphics screen on an Epson MX-80 printer. Only those character positions that have the plot bit set in their color byte will be printed; all others will be treated as blanks. The reason $I$ chose not to print other characters is due to the difference in size of an ISC plot block (4 by 2) versus an $M X-80$ plot block (3 by 2). In order to faithfully reproduce the ISC graphics, the screen is scanned three $Y$ positions and two $X$ positions at a time to form an MX-80 graphic character. Since a character such as an "A" occupies what looks to the printer like one and a third lines, to avoid gaps it is not printed.

The program consists of three sections: initialization, dot scan/character build/print, and point test. The following is a breakdown of the program, by section and line.

## InITIALIZATION

Set up the array MS, which is used by the point test
subroutine. These lines must be executed before the dump

DOT SCAF/CHARACTER BUILD/PRINT
9000 Disable keyboard.
9010 Set up for stop bit (PLOT 14 . . 15), 1200 Baud (PLOT $27,18,4$ ), and send all output to the RS-232 port (PLOT $27,13)$. Set print width to 132 characters per line. The PLOT 14 at line 9030 sets double width characters, which change the line width to 66 . The PLOT 14 must be reissued for every line output.
9020-9040 Scan the CRT Y positions in increments of three and $X$ in increments of two. This corresponds with the number of dots per MX-80 block.
9050-9200 CH is the character to be sent to the printer and is built
one dot at a time. The characters are located from CHR\$(160) to CHR\$(223), or all dots off to all dots on. On return from the point test routine, PT will be set to -1 to skip two X points, 0 if that $X, Y$ point is off, or $>0$ if that $X, Y$ point is on. The printing process will take 20 to 25 seconds per printer line (around 15 minutes for a whole screen) depending on the number of blanks. The code at line 9060 is to handle the scanning at the next to last lines.
9210 Move to top of the next sheet of paper.
Return output to screen.
9230-9240 Enable keyboard interrupt and return.

## POINT TEST - TEST POINT AT SX, SY

9510
9520

9530
9540

9550-9560 Return $P T=-1$ if two $X$ tests can be skipped as being unprintable (each character is two points wide).
9570
9580
Initialize PT.
Compute the address within the screen buffer of the character the dot specified by $S X, S Y$ is within.
Get the ASCII code and color byte of the character.
If it is not a blank and is a graphic character, continue testing. Convert SX and SY to an index into MS. See if point is on. MS and DO are used to isolate a particular bit in the byte stored for the graphic character. (See your Programming Manual for the relationship between a graphic point and the byte stored in memory.) PT will be zero if the point is off and non-zero if the point is on.

As mentioned above, it will take around 15 minutes for a complete picture to print. Translation of the algorithm into machine language will undoubtedly speed it up. (Any takers? - ed.)

To use the subroutine, simply bury it in a program that draws a picture, and GOSUB 9000 to start it. If you do not have the y-eliminator, put some sort of delay in before calling the subroutine to allow time to set up the printer and turn it on.

Listing 1. The screen dump subroutine

| 0 | DIf 朋(7) | 9150 | NEXT J |  |
| :---: | :---: | :---: | :---: | :---: |
| 1 | FOR I= 0 T0 7 | 9150 | NEKT I |  |
|  | :MC(I) $=$ INT $\{2 \times \mathrm{I})$ | 9170 | PRINT [HE\$ (CH): |  |
|  | :NEXT I | 9180 | NES |  |
| 9000 | OUT 8,4 | 9170 | FRINT CHF (13) |  |
| 9010 | PLOT 14,27, 18, 4, 15, 27, 13, 15, | 4200 | NEXT Y |  |
| 9020 | FOR $\gamma=0$ TO 126 ctep 3 | 9210 | PRINT CHR ${ }^{\text {(12) }}$ |  |
| 9030 | FLOT 14 |  | :REM OUTPIT A FGEMFEED |  |
| 9040 | FOR $\mathrm{L}=0$ T0 126 STEF 2 | 9220 | FOKE 33265,0 |  |
| 9050 | $\mathrm{CH}=160$ | 9230 | DUT 8,255 |  |
|  | : $M=1$ | 9240 | RETUEN |  |
|  | $: I M=2$ | 4500 | REM SUBROUTINE TO SEE If FIINT SX, cy is on. | PT=0 If MOT. |
| 9060 | IF $Y=126$ THEN $I M=1$ | 9510 | FT $=0$ |  |
| 9070 | FOR $I=0 \mathrm{TGI}$ IM | 8520 |  |  |
| 9080 | FOR $\quad 1=0$ TO 1 | 7530 | $D A=P E E K, ~(A D) ~$ |  |
| 9090 | $5 x=8+1$ |  | :CL $=$ PEE $(14 B+1)$ |  |
| 9100 | $5 Y=Y+1$ | 9540 | If DRC > 32AND CL ${ }^{\text {c }} 12760 T 09570$ |  |
| 9110 | 6U5UE 9500 | 950 | FT=-1 |  |
| 9120 | IF $\mathrm{PT}=-1$ THEN M $=4$ | 9560 | RETUPN |  |
|  | : 80 T0 9160 | 9570 |  |  |
| 9170 | IF FTY OTHEN CH= CH+M | 9580 | PT= MSIOD) AND DA |  |
| 9140 | $\mathrm{M}=\mathrm{M}$ \% 2 | 9590 | GETURN |  |

LATE NEWS: Compuworld Inc. has announced business software packages for Intecolor computers. Programs include ColorCalc, ColorGraph, General Ledger, and others. (See the August, 1981 issue of Mini-Micro Systems.) We' 11 have more on these next issue. In the meantime, for more information contact Fred Calev, Compuworld, 125 White Spruce Blvd., Rochester, NY 14623. (716) 424-6260.


## User Group Bulletin Board

The following list represents our most current information on Intecolor/Compucolor User Groups. In the interest of keeping our readers informed and in touch, we plan to run this list periodically. If you have additions (because your group is not listed) or changes (because names or numbers are different), drop a line to our editors so we can tell the world. We'd like to include news of your group here, too. Have your secretary write a piece about your group, its activities, interests, and members. We also welcome pictures. ( 5 x 7 black and white glossies, with good contrast.)

User groups are a very important medium for information transfer, and we at Colorcue want to provide a platform for group information interchange. Get your secretaries on the ball!

## ALABAMA

Compucolor Users Group
Eike Mueller
12117 Comanche Trail, S.E.
Huntsville, Alabama 35803
(205) 883-7614

```
San Jose User Group
Vicki Oliver
1358 Branham Lane 非4
San Jose, California 95118
(408) 267-5250
The SOURCE: CL0691
```


## CALIPORNLA

Compucolor/Intecolor User Group
Stan Pro
S.P. Electronics Systems

5250 Van Nuys Boulevard
Van Nuys, California 91401
(213) 788-8850

E1 Cerrito User Group
Frasier Hewitt
c/o P.C. Computers
10166 San Pablo Avenue
E1 Cerrito, California 94530
(415) 428-0468 (F. Hewitt)

GOTO Group
Tommy W. Schenck
c/o Tom \& Bobbie of Newberrys
1136 Fulton Mall
Fresno, California 93721

San Diego User Group
Hal Brehe
4671 Mt. Arnet Drive
San Diego, California 92117

## FLORIDA

JACKS - Jacksonville Area Compucolor Knowledge Seekers
Gary Haney
1723 Debbie Lane
Orange Park, Florida 32073
(904) 264-6785

Robinson High Computer Club
Mrs. Byman
6311 S. Lois Avenue
Tampa, Florida 33616
(813) 835-1211

GEORGLA
Compucolor Users Group Irv Mullins
2194 Briarcliff Road, N.E.
Atlanta, Georgia 30329
Day: (404) 586-5156
Night: (404) 634-3919

## MASSACHUSETTS

Compucolor Users Group
Richard Manazir
13 Grandview Street
Southwick，Mass． 01077
Day：（203）688－1911
Night：（413）569－6621

## NEM JERSEY

Compucolor Users Group
Peter J．Miller
125 Buena Vista Drive
Ringwood，New Jersey 07456
Night：（201）839－7251

## NET YORK

Compucolor Users＇Group of Rochester Gene Bailey
28 Dogwood Glen
Rochester，New York 14625
（716）381－4027

## OREGON

Compucolor II Users＇Group
Bruce Vanderzanden
2006 ＂C＂Street
Forest Grove，Oregon 97116
（503）357－2772

## PENNSYLVANIA

Philadelphia CCII Users Group
Howard Rosen
P．O．Box 434
Huntington Valley，Pa． 19006
（215）464－7145

## VIRGINIA

Compucolor Users Group
Rick Vick
702 W．Holly Avenue
Sterling，Virginia 22170
Day：（703）827－3894
Night：（703）430－3843

## INTERNATIONAL

Ham Radio Users Group
Bill Shanks，W2GTX
7 Lake Circle Drive
Vicksburg，Mississippi 39180

AJSTRALIA
Compucolor User Group／Syd－ ney Area
Andrew licIntosh
91 Regent Street
Chippendale，IJ．S．W．， 2008
Australia
Compucolor User Group／ilel－ bourne Area
iveil Brandie
212 High Street
Windsor，Victoria， 3181
Australia

## CANADA

Canadian Users Group
Glen Davis
Bsmt－59 Kendal Avenue
Toronto，Ontario
Canada M5R 1 L8

Compucolor Users Group
Mark Herzog
House of Computers
368 Eglinton Avenue，West
Toronto，Ontario
Canada M5N 1A2
（416）482－4336

## Cueties

10 RFM COLORED SQUARES
20 REM BY STEVE REUBART
30 PLOT $27,24,15,29,6,0,12$
$40 \mathrm{IN}=1: \mathrm{C}=1$
50 FOR Y＝Ø TO 31 STEP 4
60 FOR X＝Ø TO 63 STEP 8
$70 \mathrm{C}=\mathrm{C}+\mathrm{IN}:$ PLOT 6，C
$30 \mathrm{FOR} \mathrm{I}=\mathrm{Y}$ TO Y＋3：PLOT $3, \mathrm{X}, \mathrm{I}$
90 PRINT＂いいいいい＂：REM 8 NULLS
100 NENT $I, X, Y$
110 FOR W＝1 TO 600：NEXT
120 IF $I N=1$ THEN IN $=-1: C=64$ ： GOTO 50
$130 \mathrm{PLOT} 27,11,6,2,8$

# Assembly Language Programming 

by David B. Suits

The Intecolor and Compucolor computers have a great deal to offer in the way of color graphics. Using BASIC to access this power is relatively easy, but its main disadvantage is speed. Assembly language programming provides the ultimate in speed, but its main disadvantage is coding time. Another disadvantage is instruction: what can you do with assembly language if you don't know anything about it? Where can you go to learn about it?

Answer: here! I sympathize with all those of you who want to learn assembly language but are afraid to begin, or don't know where to begin. This series of articles will be devoted to you, the rank beginner. I will assume you know a little bit about programming your Intecolor or Compucolor in BASIC. Other than that, all you need is a little curiosity, and some perseverance. (And a subscription to Colorcue.)

Inside your computer's plastic shell there is an 8080 A microprocessor (call it " 8080 " for short), connected to some Read Only Memory (ROM) and some read/write memory, also called Random Access Memory (RAM). In addition, the 8080 has a tiny bit of read/write memory inside itself, referred to as the 8080 's registers. Each of these registers is eight binary digits (one byte)--except for the Stack Pointer (SP) and the Program Counter (PC), which are 16 bits (two bytes) each. Certain of the eight bit registers may be grouped together into pairs: registers $D$ and $E$ can be paired and referred to as register pair $D$ (some people prefer to be explicit and say DE). H and l can be referred to as register pair $H$ (or HL ), and so on. For the most part you will be concerned to read from or write to registers $A, B, C, D, E, H$ and $l$, and, often, the register pairs $B C, D E$ and HL. PC, SP and FLAGS are used automatically
REGISTER

| PC |  |
| :---: | :---: |
| SP |  |
| FLAGS | A |
| B | C |
| D | E |
| H | L |

ONE BYTE
ONE BYte
by the 8080 for its own purposes, and we won't bother with them for quite a while, except for the FLAGS, which we'll make use of frequently.

The 8080 microprocessor is capable of understanding and carrying out a varied set of instructions, each of which is simply a number (eight bit binary) held somewhere in the computer's memory. The 8080 will look at memory, copy the number it finds there into a secret internal register of its own, and execute the instruction. This means that the 8080 must know where in memory to find that instruction number in the first place. That's the function of the Program Counter (PC), which points to a specified memory location. When the information (always an eight bit binary number) is read into the 8080 , the PC is automatically incremented to point to the next instruction in memory: that is, it is automatically set up to get the next instruction as soon as the present instruction has been carried out. More or less as in BASIC, the 8080 will execute an instruction and then move on to the next one, execute it, move on to the next one, and so on. And, more or less as in BASIC, you can use GOTOs and GOSUBs to veer off to other parts of the program, except that the 8080 will jump hither and yon according to JMP and CALL instructions. More accurately, since the 8080 can understand only binary numbers, and not letters, it will GOTO (JMP) somewhere or other in response to the binary number 11000011 ( $=195$ decimal), and GOSUB (CALL) some subroutine or other in response to binary 11001101 ( $=$ decimal 205). Since we humans have difficulty in talking in long strings of 1 s and 0 s , we write programs using mneumonic names like JMP and CALL, which are then translated into binary numbers by a program called an assembler. So there are two steps to assembly language programming: (1) You write the program in assembly language, perhaps using ISC's Screen Editor (or, Heaven forbid, the old Text Editor), and save it on disk. (2) You run the assembler (or the macro-assembler), which translates your assembly language program into a series of binary numbers and stores the result (we can now call it machine language, since it is in a form which the 8080 can understand) on disk. You will need, then, to learn how to write an assembly language program, and then you will have to learn how to use the assembler program to translate what you've written. We'll cover both of those topics (and more) in this series.

A crucial difference between BASIC (and other high level languages) and assembly language is that in assembly language you do everything one tiny, painstaking step at a time. Most single instructions in BASIC would require a large number of assembly language instructions. Whereas in BASIC you could write $\mathbf{A}=\mathbf{B}+\mathbf{C}$, in assembly language you would have to
break it down into a series of minimal steps:
(1) Find the location in memory called "B".
(2) Copy the contents of location "B" into an internal register of the 8080 .
(3) Find the location in memory called "C".
(4) Add the contents of "C" to that internal 8080 register which has been holding the contents of "B".
(5) Find the location in memory called "A".
(6) Copy the contents of that 8080 internal register (which now contains the sum of "B" and "C") into that memory location called "A".

Steps (1), (3) and (5) are themselves a series of operations which must be broken down into a number of discrete 8080 instructions. (To make matters more complex, BASIC uses not one, but four locations in memory to store a number.)

Before we get into a discussion of the assembly language instructions and what they mean, let's get accustomed to binary and hexadecimal arithmetic. If you don't already know something of binary and hex, then the following will serve as a brief introduction. The real key to learning is experience. That means practice. And practice. The more you fiddle around, the more you will feel at home. (Isn't this how many of us learned BASIC?)

## BINARY NUMBERS

When you add one to a (decimal) number, you get the next higher number. Take the largest number representable by a single decimal character, 9, and add one to it. You get 10. The digits shift just like the wheels in an odometer: once a character reaches 9, it shifts back to 0 and the character to the left moves up one. Each column in a decimal number is occupied by some decimal character. The value of the whole decimal number can be determined by multiplying the value of the character by its column's weight. A decimal number has a one's column, a ten's column, a hundred's column, and so on. These are merely the weights each column has.

| weight | 1000 | 100 | 10 | 1 |
| ---: | ---: | ---: | ---: | ---: |
| example number | 0 | 2 | 3 | 8 |

The decimal number 238 is two one hundreds plus three tens plus eight ones: $(2 \times 100)+(3 \times 10)+(8 \times 1)=238$. These weights are simply powers of ten
(powers of ten, because this is the decimal number system). The powers start with $0: 10^{\circ}=1$. (In fact, any number to the zero power $=1$.)

| powers | 3 | 2 | 1 | 0 |
| ---: | :---: | :---: | ---: | ---: |

We now have all the principles necessary to develop any other number system we please. For the binary number system, just repeat the previous discussion, replacing "decimal" with "binary", the ten allowable decimal characters with the two allowable binary characters (1 and 0 ), and so on.
fa...:
A binary odometer would turn considerably slower than the decimal odometer (but there would be no savings in gas):

| 000 | 001 | 010 | 011 |  |
| :--- | :--- | :--- | :--- | :--- |
| $+\quad 1$ | $+\quad 1$ | +1 | +1 |  |
| 001 | +010 | +011 | 100 |  |

As soon as you reach 1 , you must start over again, whereas in the decimal system you don't have to start over until you reach the character "9". Don't misinterpret these numbers. The binary number 10 is not equal to the decimal number 10 . If you start at zero, you will have counted to two when you get to binary 10 , but you will have counted to ten when you get to decimal 10 .

Each column of a binary number has a weight, determined by the powers of two:

$$
\begin{array}{rrrrr}
\text { powers } & 3 & 2 & 1 & 0 \\
\text { binary weights } & 2^{3}=8 & 2^{2}=4 & 2^{1}=2 & 2^{0}=1
\end{array}
$$

Thus, abinary number is represented using the same underlying principles as a decimal number. If we wish to translate binary into decimal, we need only multiply the column weights by the column numbers and add them up. For example:

```
weight
```

(expressed
in decimal) $\quad 2^{3}=8 \quad 2^{2}=4 \quad 2^{1}=2 \quad 2^{0}=1$
$\begin{array}{llllll}\text { example number } & 0 & 1 & 1 & 0 & \text { (binary) }\end{array}$

In decimal this would be $(0 \times 8)+(1 \times 4)+(1 \times 2)+(0 x 1)=6$. Your computer stores numbers as eight binary digits (one byte), so we ought to become accustomed to this table:

## powers

(expressed
in decima1) $\quad 2^{7}=128 \quad 2^{6}=64 \quad 2^{5}=32 \quad 2^{4}=16 \quad 2^{3}=8 \quad 2^{2}=4 \quad 2^{1}=2 \quad 2^{0}=1$

This means, incidentally, that the largest number which can be stored in a given memory byte is 11111111 (binary) $=255$ (decimal). Some of the 8080 's internal registers are 16 bits wide; in addition, each memory location in the computer is referenced by its address (more on that later), which is 16 bits. So we ought to become familiar with this table:

$$
\begin{array}{cccccccl}
2^{15}=32768 & 2^{14}=16384 & 2^{13}=8192 & 2^{12}=4096 & 2^{11}=2048 & 2^{10}=1024 & 2^{9}=512 & 2^{8}=256 \\
2^{7}=128 & 2^{6}=64 & 2^{5}=32 & 2^{4}=16 & 2^{3}=8 & 2^{2}=4 & 2^{1}=2 & 2^{0}=1
\end{array}
$$

The largest number which two bytes can hold (and the highest address allowed) is therefore 1111111111111111 (binary) $=65535$ (decimal).

With some "hunting and pecking", these tables will allow you to translate from decimal to binary and also back again. To translate from decimal to binary, find the bit with the highest weight equal to or less than the decimal number to be converted. Make the corresponding bit in the binary number $=1$ and subtract its weight from the decimal number. Resume the process with the result of that subtraction until the result is zero. For example, we can translate 238 (decimal) into binary by noticing that the highest binary weight equal to or less than 238 is 128. Make that bit $=1$.

$$
\begin{array}{rrrrrrrrrr}
\ldots & 512 & 256 & 128 & 64 & 32 & 16 & 8 & 4 & 2 \\
0 & 0 & 1 & & & & & & &
\end{array}
$$

Subtract 128 from 238 to get 110 . The highest binary weight equal to or less than 110 is 64 . Set that bit $=1$.

$$
\begin{array}{rrrrrrrrrr}
512 & 256 & 128 & 64 & 32 & 16 & 8 & 4 & 2 & 1 \\
0 & 0 & 1 & 1 & & & & & &
\end{array}
$$

110-64=46. The bit under 32 becomes 1 .

| $\ldots$ | 512 | 256 | 128 | 64 | 32 | 16 | 8 | 4 | 2 |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| 0 | 0 | 1 | 1 | 1 |  |  |  |  | 1 |

$46-21=14$, so the bit under 8 becomes 1.

$$
\begin{array}{rrrrrrrrrr}
\ldots & 512 & 256 & 128 & 64 & 32 & 16 & 8 & 4 & 2 \\
0 & 0 & 1 & 1 & 1 & 0 & 1 & & &
\end{array}
$$

14-8=6. Etc. The final result is 11101110 (binary).

| bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
|  | 1 | 1 | 1 | 0 | 1 | 1 | 1 | 0 | (binary) |

An alternative method is to divide the decimal number by 2 and place the remainder into the right-most bit of the binary number. Now divide the quotient (i.e. the integer result) of the previous division by 2 , put the remainder in the next bit position to the left, and so on until the quotient $=0$.

```
    119 remainder = 0 (bit 0, i.e., right-most bit)
2\longdiv{238}
    59 remainder = 1 (bit 1)
2\longdiv{119}
    29 remainder = 1 (bit 2)
2\longdiv{59}
    14 remainder = 1 (bit 3)
2\longdiv{29}
14
    remainder = 0 (bit 4)
2\longdiv{14}
    remainder = 1 (bit 5)
2\longdiv{7}
2\longdiv{1}
    remainder = 1 (bit 6)
2\longdiv{0}
```


## HEXADECIMAL NUMBERS

The problem with binary numbers is that they are so confusing to work with: all those zeros and ones create a fertile environment for errors. Decimal numbers are easier, because there are fewer characters per number and because we're already familiar with them. Unfortunately, the computer is too stupid to deal with ten different characters. We might use decimal numbers for our own figuring and then translate between decimal and binary when required. But that is usually a tedious operation. What we need, then, is a way of expressing binary numbers without being restricted to two digits, and yet without the burden of tedious translation. Enter the hexadecimal system.

The binary system is built upon powers of two (that's 10 in binary). The decimal system is built upon powers of ten (that's 10 in decimal). So the hexadecimal system is built upon powers of 6 (hex) $+10(\mathrm{dec})=$ 16 (that's---you guessed it!--10 in hexadecimal). Hexadecimal is convenient because 16 is a power of two (the fourth power). This means
that we can represent four binary digits with one hexadecimal digit. Or, that is, two hex digits will represent one byte, and four hex digits will represent two bytes. That convenient symmetry does not exist between binary and decimal. (Ten is not an integer power of two.) This does two things for us: (1) it decreases the number of characters necessary to represent a number, and (2) it allows a human to convert quickly between binary and hex as the occasion demands.

The conversion is simple: start with the four least significant bits (i.e., the right most four bits, or half a byte--sometimes called a nibble, or nybble) and translate them into its hex equivalent. Move left and convert the next nibble, and so on. To convert from hex to binary, just replace each hex digit with its binary equivalent.

Well, in order to make these conversions, we'll have to define the hexadecimal characters. We need sixteen of them, but we are accustomed to using only ten ( $0-9$ ). So we invent six more. It really doesn't matter what they are, except that tradition calls for using the letters A-F. (Perhaps because these symbols were easily available on typewriter keyboards.) The relationships between the binary, decimal and hexadecimal schemes are shown below. After working with hex and binary numbers for a while, you will have this table memorized.

| HEX | BINARY | DECLMAL |
| :---: | :---: | :---: |
| 0 | 0000 | 0 |
| 1 | 0001 | 1 |
| 2 | 0010 | 2 |
| 3 | 0011 | 3 |
| 4 | 0100 | 4 |
| 5 | 0101 | 5 |
| 6 | 0110 | 6 |
| 7 | 0111 | 7 |
| 8 | 1000 | 8 |
| 9 | 1001 | 9 |
| A | 1010 | 10 |
| B | 1011 | 11 |
| C | 1100 | 12 |
| D | 1101 | 13 |
| E | 1110 | 14 |
| F | 1111 | 15 |
| 10 | 00010000 | 16 |
| 11 | 00010001 | 17 |
| - | - | - |
| - | - | - |
|  |  |  |

If you recall the principles which structure all number systems, you should have no trcuble dealing with hex. A hex odometer would turn slower than a decimal one:

$$
\begin{array}{cccccccccc}
00 & & 09 & 0 A & 0 B & O C & O D & 0 E & 0 F & 10 \\
+1 & \cdots & +1 & \frac{+1}{O A} & \frac{+1}{O B} & \frac{+1}{O C} & \frac{+1}{O D} & \frac{+1}{O E} & \frac{+1}{O F} & \frac{+1}{10} \\
\hline 01 & \frac{+1}{11} & \cdots
\end{array}
$$

I said before that one reason for not using decimal numbers was because the computer was too primitive to handle ten different characters; two is its limit. (After all, the computer is merely a clever arrangement of switches, each of which can be either on, $=1$, or off, $=0$.) And now I offer the hexadecimal system with its sixteen characters. So? I lied.

No, no. It is true that the computer can't deal directly with anything other than binary numbers, but hexadecimal numbers at least provid a fast and efficient manner in which to represent binary numbers to humais;, which is what we are. We can always wrice a simple routine to translate between hexadecimal and decimal if we need it. Besides, hex is firmly rooted in tradition (short as the computer tradition is), and any attempt to rip it up will probably cause the destruction of the universe.

Your computer's FCS understands only hex. (This makes it difficult for BASIC, which understands only decimal, to talk to its own FCS! Computer manufacturers work in mysterious ways.) Furthermore, ISC's Machine Language Debug Package (a spectacular piece of software, by the way) is easier to deal with in hex (although you can get by with decimal), and we will be using it in this series. (Do you have yours yet?)

Next time: the 8080 instruction set. In the meantime, read anything and everything you can get your hands on dealing with 8080 assembly language. There are articles in magazines, books, and even assembly language program listings--cryptic as they might seem at first--can heip you become acquainted with this new language. Every human is partly a self-teacher. You'd be surprised at how much information you can assimilate simply by looking over some 8080 assembly language listings. At least you will come to have an unconscious "feel" for the language. And that's important.

## HONARD ROSEN, INC.

```
    Futt the firishirig touches to your Compucolor II or ISC computer.
    Come up to the worla of word processirig.
    Evterid the utilizatiori of wour computer to the other members of your family.
        * Letters
        * School reports
        * Busiriess reports
        * If you now type-write it, Camp-U-Write it for a better product.
* Essic requiremerits for CCII or 3651/9651:
        16F! FiAM.
        1!7key keyboard.
        Fririter.
        Comp-U-Writer software arid iristruction mariual.
    For maximum capabilits!
        FBll 32r, FiAM.
        Lower case characters.
    Talk to other computers: Ads a MODEM to your system.
    We carry the entire CCII & ISC line of naroware/software, incluring spares.
    Send for our 4-page order form for hardware/software+ Request separately
    ly item your spare parts rieeds.
```

    Sens your order now. We pay the shippirig.
    Allow 5 weeks for delivery.
    

HOWAFD FOSEN. IMC.
FO EOK 434
Hiritirigdori Valles, Fas, 1900
Sigristure (please sign orger)

# Colorcue <br> Editorial Offices <br> 161 Brookside Dr. Rochester, NY 14618 

## Colorcue



Oct/Nov 1981

## Colorcue

## A bi-monthly publication by and for Intecolor and Compucolor Users

## Editors:

October/November 1981
Ben Barlow
Volume 4 , Number 2
3 Editors' Notes
4 ISC Product Review
6 Assembly Language Programming, by David B. Suits Part II: Using the Machine Language Debug Package
15 Sphere Program, by Mark D. FairbrotherGenerate a 3-D sphere on your screen
17 Incremental Plot Table, by Bob V. Smith
Find those incremental plot numbers the easy way
18 Lissajous Figures, by Trevor Taylor himicking oscilloscope graphics
20 Ask Mr. C. O. Lorcue
21 CALLable Sort Routine, by Alan D. MatzgerUltra fast sorting for your BASIC programs

Advertisers: You will find our advertising policies attractive. Urite for details.

Authors: This is a user-oriented and supported publication. Your articles/tips/hints are required to make it go. Send your articles or write for information.

Colorcue is published bi-monthly by Intelligent Systems Corporation, with editorial offices in Rochester, New York. Editorial and subscription correspondence should be addressed to the Editors, Colorcue, 161 Erookside Dr., Rochester, NY 14618. Froduct related correspondence should be addressed to ISC, 225 Technology Park, Norcross, GA 30092, ATTN: Susan Sheridan. Opinions expressed in by-line articles are not necessarily those of the editors or of ISC. Hardware/software items are checked to the best of our abilities but are ${ }^{\text {HOT }}$ guaranteed.

## Editors' Notes

Tine last issue of Colorcue carried an article by Toin Devlin about an elegant, inexpensive solution to the lower case $y$ problen that has bugged Compucolor owners for so long. As it frequently happens, inventors seem to come up with nearly the same ideas at nearly the same time. Ken Orford, a member of Foran, the Canadian users group, hit upon the same fix for the problem several months before Tom and published it in the July/August issue of the Forman newsletter. Since by that time we had already received Ton's article, we ran it in our August/September issue. But we neglected to mention Ken's article, so we want to take this opportunity to acknowledge Ken's efforts and Form magazine. It is a fine newsletter published six tines per year. Subscription is \$15. By the way, our User Group Bulletin Board in the last issue of Colorcue was out of date for Canada. Forum is the only user group. For information, write to Forum, 21 Dersinghan Crescent, Thornini11, Ontario, Canada L3T 4P5.

## DEBUG BUG DEBUGGED - Dave Suits

There is a fatal bug in at least some versions of the llachine Language Debug Package prograrn. In order to see if your version has a bug, and to correct it if it does, read on.

The MLDP disk comes with two versions of the Debug program. On the disk they are naned MLDP.PRG;1 and MLDP. PRG;2. Only version one has the bug. If you've a 32 K machine, you probaíly won't notice it. (Besides, you.'11 usually be
using version two.) But if you have a 16 F machine, you must use version one (since version two loads into very high memory--which would be non-existent on a 16 K machine). Enter FCS with ESC D and then run ILDP;1. You should be greeted with a heading and the DBG> prompt. Type in this: @A03A and hit RETURN. There should be a yellow A03A on the left. (If it's not there, try again.) Off to the right, and in white, you will see H,0E000H. If that's not there (for example, if you have H,OA000H), then your liLDP program is OK. But if it is there, then so is the bug. To correct it, type in: LXI H,A000 and hit RETURN. You should now be shown, on the left, a yellow A03D. (If that's not there, go back to the first step by typing @A03A and try again.)

You have now altered the MLDP program, and you should save this corrected version on disk. To do that, first hit CPU RESET and ESC D. If you have single density disks (i.e. Compucolors) then you'11 have to mount a fresh disk: there's no roon on the original disk to save the corrected progran. Now type: SAVE MDP A000-BF00. This new version of MLDP will now run correctly on both 16 F and 32 x machines.

The bug, in case you're interested, was the program's initializing the pseudo stack pointer to E000H. That will work witi 32F machines, but on 16 K machines, EOOOH is a non-existent RAM address. The alternative to the permanent correction is to correct the error each time the ILDDP progran is restarted (either run fron FCS or else entered with (ESC) (USER)). In this case, you will have to re-initialize the pseudo stack pointer by typing \%SP=A000.

## EPSON NOTES - Ben Barlow

Those of you with Epson printers have no doubt seen the advertisements for GRAFTRAX 80 , the bit graphics package for the $M X-80$. I got mine in the mail the other day, and what $I$ read sent me scurrying for the Aug./Sept. issue of Colorcue in which $I$ had written: "The Epson must be ordered with the RS/232 interface (Cat. no. 8141)". Then back to the front page of the GRAFIRAX 80 documentation: "Note: GRAFTRAX 80 will not function with the Epson 8141 serial interface". Curses! Not only can $I$ not try GRAFTRAX 80 , since I've got an 8141 interface, but I've possibly misled some readers into buying into the same dilemma.

If so, I offer apologies. If not, I can think of two solutions. One, purchase the 8150 version serial interface port (it costs about twice the 8141), or, two, build a serial to parallel interface and connect it to the $M X-80$ 's parallel connector. Always looking for the cheap solution, I'm opting for the second method. I hope to
have an article about the parallel interface and GRAFTRAX 80 in the next Colorcue.

Having heard rumors of printer nalfunctions when using GRAFTPAX30, and seeing a prominent warranty disclaimer in the documentation, I called Epson America to inquire. I was told that there were no problems using GRAFTRAX-80. It can cause no problems with the hardware because it is simply a replacement program for the printer. The warranty disclaimer is there to protect against claims due to customer mis-installation or claims arising from applications of the printer (for example, a bug in the program using a printer connected to an EKG machine, causing the death of three patients). I was advised of a problem with a kit for a friction feed roller to add to the MX-80. The kit apparently causes the stepper motor, which advances the paper, to burn out. Its installation will void the Epson warranty. (Epson has no modification or plans for one to add friction feed. They advise the use of the MX-80FT printer.) E

## INTELLIGENT SYSTEMS CORPORATION PRODUCT LINE REVIEW

Intelligent Systems designs, manufactures, markets and supports several different types of color graphic microcomputers. These products are used in varied applications such as process control, energy management, and network monitoring. There are tens of thousands of units in use in over 40 countries.

Currently, the product 1 ine consists of three different computer types, each of which can be ordered with various options such as different screen sizes, different memory amounts, printers, software, etc.

The "workhorse" of the Intecolor product line is the 8000 series. This series includes terminals, desktops, and complete computer systems. These
computers are in industrial cabinets and have card cages for adding features as required. All prices quoted below are for single quantity.

8001 G termina1 $\$ 2560$
Includes refresh RAM, terminal control software, $72-$ key keyboard, 19" CRT, RS-232C, 8 foreground 8 background colors, graphics software and editing features. 160 x 192 resolution.

## 8001I terminal $\$ 4460$

Same as 8001 G , but with graphics resolution of 480 x 384 . Each dot on the screen is addressable to any of eight colors. Includes dot, arc, vector, and incremental vector generators. 4 by 4 dot super-pixel yields 4913 individual lines.

8051 desktop $\$ 4265$
Complete desktop computer with color graphics, $19^{\prime \prime}$ CRT, Disk BASIC, lower case ASCII, operating system, 48 lines by 80 charactexs, 80 K mini-disk and 8 K of user RAM.

When the setting requires a contemporary cabinet, Intelligent systems has the 8300 and 8900 product series, with $13^{\prime \prime}$ and $19^{\prime \prime}$ CRTs respectively. These fit well in business environments, where they are used for MIS, conmunications, and word processing. Each of the following pairs has features comparable to the corresponding industrial units described above.

| $8301 G$ | $\$ 3560$ |
| :--- | :--- |
| $8901 G$ terminal | $\$ 2760$ |
|  |  |
| 8301 I terminal | $\$ 5460$ |
| 8901 I terminal | $\$ 4660$ |
|  |  |
| 8351 desktop | $\$ 5265$ |
| 8951 desktop | $\$ 4465$ |

Many users do not require the high density of information that can be displayed in the 80 character by 48 line format of the 8000 series. For these users, Intelligent Systems manufactures a line of single-board computers whose less dense (and less expensive) displays show 64 characters by 32 lines. These computers have fewer available options than the 8000 series products, but their low price makes them very attractive, They are
(Contimued on page 13.)

# Assembly Language Programming 

PART II: USING THE MACHINE LANGUAGE DEBUG PACKAGE
by David B. Suits

I mentioned last time that instructions 2 the 8080 are always eight binary digits (one byte). This is misleading. With one byte there are 256 possible numbers (0-255), but the 8080 does not understand 256 different instructions. The number 237, for example, is undefined: what the 8080 will do, if anything, with that number is unpredictable.

There is another sense in which it is misleading to say that the 8080 instructions are one byte numbers. Sometimes ( $I$ should say "often") the 8080 will know what to do, in response to an instruction code, but it will have to be given certain data to do it with, or else it will have to be told where to get that data. Let me explain.

In response to the number 71 ( $=01000111$ binary) the 8080 will copy the contents of the $A$ register (the Accumulator) into the $B$ register. That's simple and straightforward. The 8080 doesn't care what is in either register. It just blindly copies $A$ into $B$. (Consequently, whatever used to be in register B will no longer exist.) Suppose A contains 0001100 and B contains 11010110 .

```
Before executing 01000ll1 After executing 0l0001ll
    A 00011100 A A 00011100
```

For each possible register copy instruction (not counting PC, SP and FLAGS registers) there is some unique number which tells the 8080 to do that copying. In assembly language you specify the copy instruction by using some mneumonic. The standard set of mneumonics was defined by Intel Corporation (the inventors of the 8080) some years ago. Unfortunately, instead of something obvious, like COPY A TO B, they chose the more cryptic MOV B,A. 'MOV' stands for 'move', but what it means is simply copy. After the MOV, the destination register is specified, and then the source register. The general format for a MOV instruction is therefore MOV <dst. reg.>, <src. reg.>. Don't forget the comma between the two registers' names. The assembler translates such MOV instructions into their appropriate binary codes.

Some other 8080 instructions, however, need more than one byte to
operaie. ior example, the number 00001110 ( $=0 \mathrm{E}$ hex) instructs the 8080 to put the following number into the $C$ register. But what is "the following number"? It is the byte which immediately follows the 00001110 itself. Fence, this is called an immediate type of instruction.

```
00001110 00000001 put l into the C register.
0000lll\emptyset 0lø0000l put 65 into the C register.
```

And so on. In assembly language, this is a move immediate instruction, whose meumonic is MVI. The general format is MVI <reg.>, <number>. Some examples:

$$
\begin{array}{ll}
\text { MVI A, } & \text { put } 0 \text { into the accumulator } \\
\text { MVI C,25 } & \text { put } 25 \text { into the C register } \\
\text { MVI L,15 } & \text { put } 15 \text { into the is register }
\end{array}
$$

Note that the number to be placed into a register must be between 0 and 255 -- i.e., one byte. The instruction MVI A, 1024 is an impossible instruction (and your assembler will give you an error message), since the number 1024 is 0000010000000000 binary and occupies two bytes, not one. (By the way, negative numbers are allowed, but that's a topic for a future installment.)

Register pairs may also be loaded with immediate data (two bytes, this time), but we'll save discussion of that for anothex time.

## Using the MLDP

ISC's Machine Language Debug Package is a harvelous tool for debugging machine language programs. It is not an assembler, although it can do some of the same things that an assembler will do. And it can be used by us to learn about onoo instructions and what they can be inade to do.

Before we jump in for some hands-on experience, please read the section in this issue's Editors' Notes dealing with a bug in the MLDP program. When you're ready to go, enter FCS with ESC D and run MLDP; (if you have a 16 k machine) or MLDP; 2 (if you have a 32 k machine). The program will Fint a heading and then give you a DBG prompt. There are a number of comands you can now cive it. For example, it will translate a number from inexadecimal into decimal. or vice-versa if we give it the "=" command. (The computer's output is in bold face in the following examples; my input is in regular type and is followed by a REIURN.)

The number typed in after the＂$=$＂is assumed to be a hexadecimal number， but if it is preceded by＂非＂，then it is taken to be deciraal．So DBG understood my command to be：＂Translate the hexadecimal number 10 into decimal＂．This is did；it printed out hexadecimal 10 and then the decimal equivalent， 16 （preceded by＂挷＂to indicate decimal）．Similarly，

```
DBG>=42-#13
=0635 $53
```

tells me that subtracting decinal 13 from hex 42 yields hex 35 （＝decimal 53）．You now have a convenient hex／dec calculator－converter．You don＇t really need binary output as well，since you can easily read off the binary equivalent from the hex number．（See the previous article in this series．）

The MLDP can also set up a series of machine language instructions and then execute them．（Furthermore，we can choose two modes of execution， as I＇ll explain below．）The command＂＠＂followed by a number tells the program to point to the address in the computer＇s memory specified by the number．For example，

DBG＞＠82b0
820日 E5＇E＇POSH H

Just above the MEM there is a line of information consisting of four fields．The first（at the left of the screen）tells you the memory location （in hex，always）．In our case it is 8200．To the right of that is a two digit hex number，or else two two digit hex numbers，or else three two digit hex numbers．Whether there is one，two，or three numbers depends upon what information is presently found at memory location 8200H．We＇ll come back to this later．To the right of the numbers is one（or two，or three）characters within single quotes．Let＇s ignore that for the moment． And finally，there are two more columns；together they comprise the assembly language mneumonics corresponding to the information stored at 8200 H ．More on that later．

Notice that the DBG＞prompt has been changed to MEM＞．This means we are in the so－called＂memory mode＂rather than the debug mode．In the MEM mode we can put a byte of data－－for example， $3 E H-$－into the present memory location by entering

```
MEM>=3E
```

After pressing RETURN, the $3 E H$ ( $=62$ decimal) was stored in 8200 H and now anotner line of display tells us what is at the next location, 8201H. Let's enter 41 H there:

MEM $>=41$
(Renember that whatever is typed in is assumed to be in hex unless preceded by "非" to indicate decimal.) Now enter these three bytes: CD, 33 and 00 .

MEM $>=C D$
MEM $>=33$
MEM $>=00$
Now let's go back where we started and see what has been accomplished. Go back to location 8200 H by entering:

MEM> 98200
You should get this result:
829. 3E 41 ">A" HVI A, S1日

This says that we are at memory location 8200 H and that the bytes 3 EH and 41 H are found there. Well.... There can be only one byte per memory location. That's the 3EH. The 41 H is evidently in the next memory location, namely, 8201 H . Why are we shown two bytes at once? The answer is more easily seen from the information at the right: MVI A, 4IR. This means, in assembly language mneumonics, "Put the number 41 H into the Accumulator." The "MVI A..." is the assembly language equivalent of 3 EH ( $=00111110$ ), which is what is stored at 8200 H . That is, the MLDP program has translated the numbers at 8200 H and 8201 into assembly language mneumonics for us. It has done just the reverse of what an assembler does, and so we say that it has disassembled some machine language instructions. Instead of telling us that "MVI A" was stored at 8200 H , it tells us something more useful, namely, "MVI A,41H".

Now press the " + " key and RETURN. You will be shown the contents of the next location in memory, namely,

```
82g1 41 'A' MOV B,C
```

Oops! What is that "MOV $B_{3} \mathbf{C}$ " doing there? That is an instruction to copy the contents of register $C$ into register $B$. Eut the MOV $\mathbb{B}_{\mathbf{s}} \mathbf{C}$ instruction apparently overlaps the MVI $\mathbf{A}, 41 \mathbb{1}$ instruction. What's going on?

The 8080 understands numbers held in memory locations in two ways: (1) as a number, such as 41 H , or (2) as an instruction, such as ... er ... 41H! How does the 8080 know that 41 H is on one occasion merely a number to do something with and on another occasion an instruction? Well, that
 assembly language, $\mathbb{N} I \mathrm{~A}$ ). But the instruction is incomplete: move what into the accumulator? Why, the next number, of course: 41II. So in this case 41 H is understood as a mere byte of data to be used somehow. On the other hand, if execution were to begin at 8201 H , then 41 h would be interpreted as MOV B,C. So if we ask the MLDP to tell us what is at location 8201 H , it will oblige by displaying 41 H and then telling us that 41 H is equivalent to $\mathbf{M O V} \mathbf{B}, \mathbf{C}$. After all, the MLDP doesn't know where we plan to begin execution.

The "+" key moved us forward one byte in memory. The "-" key will move us backward one byte. Try it. You should be back at 8200 H .

Now press RETURN without entering anything. This will move you ahead not just one byte, but rather to the next instruction, even if, as in the case of $3 E H 41 H$, the instruction takes up more than one byte. Now you should see:

## 8202 CD 33 日6. 'M3@' CALL 0033H

The CDH (that's a hex number, by the way, and not some kind of mneumonic) is a CALL (like GOSUB in BASIC) instruction. But CALL what (or where)? In BASIC, a GOSUB must have a line number to GOSUB to. When you're talking directly to the 8080 , you will have to give it an address. Remember that memory addresses are two bytes long, and so an address can be anywhere from 0 decimal to 65535 decimal. So 33 H 00 H following the CDH represent (in hex) the address of the subroutine. But notice that out to the right it says CALL 0033H, whereas to the left it says CD $\mathbf{3 3} \mathbf{0 0}$. Why have the "33" and the "00" been reversed? The 8080 always assumes that the low byte (i.e. the right-most, or least significant eight bits) of the address comes first, then the high byte. Some people feel that this is backwards, but $I$ have found it quite natural. Just remember low byte first in memory, eventhough we write it with the high byte first.

When the 8080 performs a CDH (=CALL) instruction, it goes off to execute some subroutine until it encounters C9ll ( $=201$ decimal, or RET in assembly language), in which event it will return to where it was before and continue. How does it know where it was before? When CDH (=CALL) is encountered, the 8080 will store the address of the next instruction after the CALL instruction. When it returns fromithe subroutine it can retrieve that address and continue. Since our CALL instruction is located at addresses $8202 \mathrm{H}, 8203 \mathrm{H}$ and 8204 H , the address of the instruction after the CALL is address 8205 H . Where is this address saved during a CALL? It is saved in a reserved area of memory called the STACK. We' 11 talk more
about that in a later installment.

The display is still in MEM mode. Let's go back to the DBG mode. We do that by entering " /" and RETURN. (Be careful not to enter / when in the $D B G$ mode, because that is a command to exit the MLDP altogether and return control to the machine's CRT mode.)

Now we're going to set a BREAKPOINT. A breakpoint is a location in memory that will halt execution of a machine language program and return control to DBG, sort of like pucting an END statement into a BASIC program to halt execution and return control to BASIC. A breakpoint is set by entering $\mathbf{A T}$ <address>. Thus,

DBG>AT 8265
will set a breakpoint at 8205 . You can put in a number of breakpoints throughout a program for debugging purposes. To list all the breakpoints, rype in L. In our case, there is only one breakpoint, so we'll see:
nBG>L
8285: C3
This tells us that a breakpoint is set at 8205 H , and that the contents of 8205 H is C 3 H. (Your memory might have something else there.) We don't care what's there; we only want to make sure that the program will stop there and not continue on trying to execute whatever random assortment of numbers might be in menory. If we wanted to clear a breakpoint, we would type in DBG>C 8205
But don't do that, because we want the breakpoint there.
Believe it or not, we now have a program to execute. There are two
 will transfer control to the 8080 , which will begin executing at the address specified. It will run at top speed until it bumps up agairst a breakpoint, in which case control will return to DBG. The I command will execute much slower and in this case execution may also be interrupted by pressing the ATTN/BREAK key. Sometimes the I conmand is helpful because you will want to see things happen in slow motion.

Let's execute the machine language instructions beginning at 8200 H . UBG>R 8260 or else DBG>I 3200

In this case it won't matter whether you ase $\mathbb{R}^{2}$ or $\mathbf{I}$, because the program is so short.

What will be the result?

DBG>R 8200
A
BREARPOINT AT 8265

(Some of the values on your display will differ from mine.) The first thing to notice is the $A$ on the second line. Our program was a simple one: it put 41 H ( $=65 \mathrm{dec}$., which is the ASCII code for " $A$ ") into the accumulator and then CALLed a subroutine in ROM (located at 33H). That subroutine takes whatever is in the accumulator and puts it on the screen. The subroutine was RETurned back to our program, where it immediately encountered a breakpoint. The rest of the display shows the resulting contents of the 8080 's registers ( $\mathrm{A}, \mathrm{B}, \mathrm{C}, \mathrm{D}, \mathrm{E}, \mathrm{H}, \mathrm{L}$ ), the contents of the Status Word (i.e. the FLAGS), the Program Counter, the Stack Pointer, and the contents of part of the stack. On the last line of the display, we are shown the contents of memory location 8205 H , where the breakpoint was encountered.

For now, nevermind all that information. The interesting thing is that we have written a program to put something onto the screen. It works just like the PLOT statement in BASIC: PLOT 65 would put an "A" on the screen at the present cursor location. The rule is, each time you want to put a character onto the screen (or change colors, or draw vectors, or do anything else you do with BASIC'S PLOT statement), just put the appropriate number into the accumulator and then CALL $33 H$. Try this:

| DBG>C 8205 |  | (clear the breakpoint) |  |  |
| :---: | :---: | :---: | :---: | :---: |
| DBG>@8200 |  |  |  |  |
| 8200 | 3 E 41 | ${ }^{\prime}>\mathbf{A}^{\prime \prime}$ | MVI | A, 411 |
| MEM> |  |  |  |  |

Now type in each of the following 40 two digit hex numbers, preceding each with "=" and following each with the RETURN key.

BASIC equivalent


You should now be pointing to location 8228 F . (If not, go back to 8200 H and check your entries.) Return to the DBG mode with "/". Set a breakpoint at 8228 H and then execute the program beginning at 8200 H . (Use either $\mathbf{R}$ 8200 or I 8200.) Notice how fast the 8080 executes when allowed full speed. The $I$ command is interesting, since the machine language program will then be interpreted, and the results will be much slower--slower, sometimes, than even BASIC. For example, the PLOT 12 instruction evidently first homes the cursor and then erases each line of the screen from the bottom up. Did you know that?

NEXT TIME we'll learn more about the 8080 instruction set, and we'll explore some much easier methods of putting stuff onto the screen. Until then, fiddle around with the present method. Have fun!
(Continued from page 5.)
especially well-suited to small businesses, fersonal use, and MIS environments.

## 3651 desktop \$2945

Complete desktop computer with color graphics, $13^{\prime \prime}$ CRT, 72-key keyboard, RS-232C, selectable baud rate, disk BASIC, internal mini-disk drive with 92 K bytes storage, 16 K user RAM. Graphics resolution: 128 x 128.

```
9651 desktop $3345
```

As above, but with $19^{\prime \prime}$ CRT.

3654 desktop $\$ 5445$
Same as 3651 but with 1182 K dual $8^{1 \prime}$ double-headed floppy disk drive.

9654 desktop $\$ 5854$

Same as above but with $19^{\prime \prime}$ CRT.


# Sphere Program 

By Mark D. Fairbrother

Carriage House East, A5
Rt. 11
Rirkwoou, NY 13795

This program will generate a 3-D line drawing of a sphere on a $128 \times 128$ screen using routines from "Mathematical Elements of Computer Graphics." Note that lines 2000 and up are for use with the $M \mathbb{X}-80$ screen dump routine if you wish to add it. (See Aug./Sept. Colorcue.) And in that case you ought to add line 10 DIM MS (7) and line 20 FOR $I=0$ to $7: \operatorname{MS}(I)=\operatorname{INT}\left(2^{\wedge} I\right): N E X T$.

100 REM GENERAIE A SPHERE USING THE SPHERICAL COORDINATE SYSTEM
110 DIM X(100), Y(100), Z(100)
$120 \mathrm{RO}=30: 1=1$
130 FOR TH=ø TO 2.51327 STEP . 628317
140 FOR $P H=\varnothing$ TO 6.28319 STEP . 339694
$150 \mathrm{X}(\mathrm{I})=\mathrm{RO}^{*} \operatorname{COS}(\mathrm{TH}) * \mathrm{SIN}(\mathrm{PH})$
$160 \mathrm{Y}(\mathrm{I})=\mathrm{RO} * \operatorname{SIN}(\mathrm{TH}) * \operatorname{SIN}(\mathrm{PH})$
$179 \mathrm{Z}(\mathrm{I})=\mathrm{PD} \mathrm{D}^{*} \operatorname{Cos}(\mathrm{PH})$
$180 \mathrm{I}=\mathrm{I}+1$
190 NEXT PH
200 NEXP TH
$210 \mathrm{P}=100$
1089 REM DIMEIRIC PROJECTION ROUTINE
1690 REM FROM 'MATHENATICAL ELENENTS FOR COMPUTER GRAPHICS'
1100 REM $P=$ NUMBER OF X,Y, $Z$ TRIPLETS
1110 REM $\quad \mathrm{X}()=$ ARRAY CONTAINING X-COORDINATES
1126 REM $Y()=$ ARRAY CONTALNING Y-COORDINATES
1130 REM $Z()$ = ARRAY CONTAINING Z-COORDINATES
1140 DIM U(100,4), V(106,4),T(4,4)
1150 FOR I=1 TO P:FOR J=1 TO $4: U(I, J)=\emptyset: V(I, J)=0:$ NEXT J:NEXT I
1160 FOR $I=1$ TO $P$
$1170 \mathrm{U}(\mathrm{I}, \mathrm{I})=\mathrm{X}(\mathrm{I}): \mathrm{U}(\mathrm{I}, 2)=Y(\mathrm{I}): \mathrm{U}(\mathrm{I}, 3)=\mathrm{Z}(\mathrm{I}): \mathrm{U}(\mathrm{I}, 4)=1$
1180 NEXT I
1190 FOR $I=1$ TO $4: F O R \quad J=1$ TO $4: T(I, J)=0:$ NEXT $J: N E X T$ I
$1200 \mathrm{~T}(1,1)=.92582: T(1,2)=.133631: T(1,3)=-.353553$
$1219 T(2,2)=.935414: T(2,3)=.353553$
$1220 T(3,1)=.377964: T(3,2)=-.327327: T(3,3)=.866825$
$1230 \mathrm{~T}(4,4)=1.0$
1240 FOR $I=1$ TO $P: F O R$ J $=1$ TO $4: F O R \quad K=1$ TO 4
$1250 \mathrm{~V}(\mathrm{I}, \mathrm{J})=\mathrm{U}(\mathrm{I}, \mathrm{K}) * T(\mathrm{~K}, \mathrm{~J})+\mathrm{V}(\mathrm{I}, \mathrm{J})$
1260 NEXT K:NEXT J:NEXT I
1279 FOR I=1 TO P
$1280 \mathrm{X}(\mathrm{I})=\mathrm{V}(\mathrm{I}, \mathrm{I}): Y(\mathrm{I})=\mathrm{V}(\mathrm{I}, 2): \mathrm{Z}(\mathrm{I})=\mathrm{V}(\mathrm{I}, 3)$
1290 NEXT I
$1306 \mathrm{~N}=3$
1310 REM AXONOMEIRIC PROUECTION ROUTINE
1320 REM FROM 'MATHEMATICAL ELEMENTS FOR COMPUTER GRAPHICS'
1330 PEM $P=$ NUMBER OF X,Y,Z IRIPLETS
1340 REM $X()=$ ARRAY CONIAINING X-COORDINATES
1350 REM $Y()=A R R A Y$ CONTAINING Y-CDORDINATES
1360 REM $Z()=$ ARRAY CONTAINING Z-COORDINATES
1379 REM $\quad \mathrm{N}=\mathrm{NUMBER}$ INDICATING THE PERPENDICULAR AXIS
1380 REM $N=1$ X-AXIS, $N=2$ Y-AXIS, $N=3$ Z-AXIS
1390 FOR $I=1$ TO $P: F O R$ J=1 TO $4: U(I, J)=\varnothing: V(I, J)=\varnothing: N E X T$ J:NEXT I
1400 FOR $I=1$ TO $P$
$1410 \mathrm{U}(\mathrm{I}, \mathrm{I})=\mathrm{X}(\mathrm{I}): \mathrm{U}(\mathrm{I}, 2)=\mathrm{Y}(\mathrm{I}): \mathrm{U}(\mathrm{I}, 3)=\mathrm{Z}(\mathrm{I}): \mathrm{U}(\mathrm{I}, 4)=1$
1420 NEXT

1430 FOR $I=1$ TO $4: F O R \quad J=1$ TO $4: T(I, J)=\varnothing: N E X T$ J:NEXT I
$1440 T(1,1)=1: T(2,2)=1: T(3,3)=1: T(4,4)=1$
1450 IF $\mathrm{N}=3$ ITHEN 1490
1460 IF $\mathrm{N}=2$ THEN 1480
$1470 \mathrm{~T}(1,1)=9:$ GOIO 1500
$1480 \mathrm{~T}(2,2)=0: \mathrm{GOIO} 1500$
$1490 \mathrm{~T}(3,3)=0$
1500 FOR I=1 TO P:FOR J=1 TO 4:FOR K=1 TO 4
$1510 \mathrm{~V}(\mathrm{I}, \mathrm{J})=\mathrm{U}(\mathrm{I}, \mathrm{K}) * \mathrm{~T}(\mathrm{~K}, \mathrm{~J})+\mathrm{V}(\mathrm{I}, \mathrm{J})$
1520 NEXT K:NEXT J:NEXT I
1530 FOR I=1 TO P
$1540 \mathrm{HX}=\mathrm{X}(1): L X=\mathrm{X}(1): H Y=Y(1): L Y=Y(1)$
1550 FOR $I=2$ TO $P$
1560 IF $X(I)>H X$ THEN $H X=X$ (I): GOTO 1580
1570 IF $X(I)<L X$ THEN $L X=X(I)$
1580 IF $Y(I)>H Y$ THEN $H Y=Y(I):$ GOTO 1600
1590 IF $Y(I)<L Y$ THEN $L Y=Y(I)$
1600 NEXT I
$1610 \mathrm{DX}=\mathrm{HX}-\mathrm{LX}: \mathrm{DY}=\mathrm{HY}-\mathrm{LY}: \mathrm{NX}=127 / \mathrm{DX}: M Y=127 / \mathrm{DY}$
1620 MOOT $12,6,6,2, \operatorname{INT}\left(0.5+M X^{*}(X(1)-I X)\right), \operatorname{INT}\left(0.5+M Y^{\star}(Y(1)-L Y)\right)$
1630 PLOT 242
1640 FOR $I=1$ TO P
$1650 \mathrm{X}(\mathrm{I})=\operatorname{INT}(0.5+\mathrm{MX} *(\mathrm{X}(\mathrm{I})-I X))$
$1660 \mathrm{Y}(\mathrm{I})=\operatorname{INT}(0.5+\mathrm{MY} *(\mathrm{Y}(\mathrm{I})-\mathrm{LY}))$
1670 PLOT X(I),Y(I)
1680 NEXT I
1690 PLOT 255
2000 REM $2 X X X$ LINES ADDED TO RUN SCREEN-PRINT IF WANTED.
2010 PLOT 15: INPUT "HARDCOPY (Y/N)";AS
2020 IF AS="N" THEN END
2040 Gasub 9000
2950 END

Latge is 65\%
of original size.


## Incremental Plot Table

by Bob V. Smith<br>498 Brown Street<br>Napa, CA 94558

This table allows you to discover the proper plot number for incremental point and bar graph plotting. In the general plot mode, PLOT 251 enters the incremental point plot subriode. The next number specifies the placement of two nore plot blocks out fromi the first, and this number can be found in the table. For example, if you want to move nor theast and tinen southeast for the two incremented blocks, then PLOT 169 is the correct number. If the oriéinal block is at, say, $X=63, Y=63$, then the instruction:

PLOT 2,63,63,251,169,255
will plot the orisinal block and the two incremented blocks.
The table also applies to incremental bar graph plotting, where the increment number will move the end point of the bar graph in the specified direction(s).

First Block

|  |  | N | NE | E | SE | S | SW | W | NV1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | N | 34 | 162 | 130 | 146 | 18 | 82 | 66 | 98 |
|  | NE | 42 | 170 | 138 | 154 | 26 | 90 | 74 | 106 |
|  | E | 43 | 168 | 136 | 152 | 24 | 88 | 72 | 104 |
| Second | SE | 41 | 169 | 137 | 153 | 25 | 89 | 73 | 195 |
| Block | S | 33 | 161 | 129 | 145 | 17 | 31 | 65 | 97 |
|  | SW | 37 | 165 | 133 | 149 | 21 | 85 | 69 | 101 |
|  | W | 36 | 164 | 132 | 148 | 20 | 84 | 68 | 100 |
|  | WW | 38 | 165 | 134 | 159 | 22 | 36 | 70 | 1¢2 |

# Lissajous Figures 

by Trevor Taylor<br>36 Tarm St.<br>Wavell Heights<br>Brisbane 4øl2<br>Australia

Lissajous figures are drawn by applying sine waves to the horizontal ( $x$ ) and vertical (y) deflection plates of an oscilloscope. The two variables that determine the shape of the figure are the ratio of the frequencies of the two sine waves and the phase angle between them.

When the frequency ratio is an integer, it will be the number of lobes (loops in the figure) if the phase angle is 90 degrees. If you consider the figure to be three dimensional, the phase angle (which is how much one sine wave is shifted with respect to the other) determines how much it is rotated. A frequency ratio of one and a phase angle of zero will draw a straight line, i.e. a circle turned on edge!

In the program, line 30 determines the foreground and background colors. The step size (degrees) used in the calculations is set in line 70. This is a compromise between speed and accuracy. Note that line 1050 and 2010 convert from degrees to radians as required for the SIn function. In lines 2040 and 2050 the sine waves are scaled for the screen and the aspect ratio is taken care of so that a circle will appear round.

Lines 2085 and 2087 take care of the case when you want only one figure drawn, which you specify by giving an increment of zero. Removing the 12 in line 1060 stops the program from exasing the screen between figures, and can give some pretty patterns. You could also modify the program to run through a range of frequency ratios.

A good example is: Frequency ratio $=3$, Starting Phase angle $=90$, Ending Phase angle $=90$, and Increment $=30$.

Try to visualize the figure rotating as the program draws successive pictures. Taking the reciprocal of the frequency ratio (for example, 0.33333 ) will turn the figure sideways, i.e., interchange $x$ and $y$ on the screen. Experiment and have fun! Try to understand the math if you can.

Note: The Australian Broadcasting Commission used this particular Lissajous figure as its symbol for a number of years. Of course, on TV it was done using an oscilloscope and rotated much faster than is possible on your computer.

```
    10 2E: LISSAJOUS FIGUPES
    11 REM
    12 REM BY GREVOR TAYLOR
    13 REI
    20 DI=3.14159
    30 PLOT 6,2
    40 PLOT 12
    50 Z=360
    60. 50=1
    70 SS=10
    79 REM IIDPOIFT X,Y OF SCRERJ
    39 XI=64:REN FOR ITTECOLOR 80G1 USE NM=30
    90 YM=64:RE1 FOR INTECOLOR 800l USE Y:I=96
    99 REH SCALING EACTORS
    100 XF=.75* (NT-1)
    110 YF=Y:-1
    500 IIPUP "EREOUENCY RATIO (0.1-10): ";R
    510 IIPUT "STAPTITG PHASE ANGLE (9-360): ";PG
    520 INPUT "EmDIIG PIASE ANGLE (0-360): ";PE
    530 INPUT "INCRE&ENT: ";IC
    540 PLOT 12
    550 IC=ABS(IC)
    560 P17=PS
1000 IF E>10 THE\ E=10
1010 IF E<0.1 MHEm: E=0.1
1920 IF E<l THEN Z=360/E
1030 IF RH<O THIET PH=-PH
104n IE Pit>36E TMET PH=PF-360
1050 A=PHI*PI/180
1060 PIOT 12:RET RHIS IANY BE OILTMED
1070 SR=YM+YF*STJ (N)
1080 RLON 2,253,年,ST,242
2000 FOR I=SS TO Z STEP SS
2010 T=T*PI/180
2020 X=SIIT(T)
2030 Y=SII (E*T+A)
2040 X=X11+XF*X
2050 Y=Y:1+YF*Y
2060 PLOT X,Y
2070 14XT I
2080 PLOT 255
2085 IF IC=0 THEM\ 2500
2037 IF GO<>0 THEN 2100
2090 IE Pif=pE THENJ 2500
2109 MI=PH+IC
2110 FOR I=1 TO 50:NEXT
2120 GO=0
2130 COTO 1000
2500 RD
```


## Ask Mr. C. O. Lorcue . . .

## Dear Mr. Lorcue,

Strange things happen when I press the INSERT LINE key on my extended keyboard when I'm in BASIC. Can you explain what is going on? Or is my machine defective? Signed: Keyed Up

## Dear Keyed Up,

The explanation is not hard to find. Each key on the keyboard is translated into a number from 0-255. The 'A' key, for example, translates as 65 (the ASCII code for ' $A$ '); the ERASE PAGE key translates as 12 ; and so on. These numbers are PLOTted by BASIC. So PLOT 65 produces an ' $A$ ' on the screen, and PLOT 12 erases the screen. The INSERT LINE key, just like CONTROL C, translates as 3. Now, what is PLOT 3? Cursor control! The next two keys you press would ordinarily determine the $X, Y$ coordinates of the cursor on the basis of those keys' ASCII values. Uinfortunately, BASIC doesn't know enough to stay in control, and it will go away after sucn key presses. You can get it back with ESC E. Unlike many keyboard controls, the INSERT LINE cannot be embedded within quotes in BASIC statements. So it looks like the only obvious use of this is for one method of cursor control in the CRT mode. For example, go into CRT mode and try these:

## Keystrokes

INSERT LINE (or CONTROL C) $=3$
SPACE $=32=\operatorname{PLOT} 3,32,0$
CONTROL @

INSERT LINE (or CONTROL C) $=3$
A7 ON $=14=$ PLOT 3,14,15
BL/A7 OFF
$=15$
(If you have questions for Mr. Lorcue, send them in and we' 11 pass them on. -Eds.)

## Cueties


PLOT 255,6, X AND 7,2,251: NEXT: PLOR 255,6,2

# CALLable Sort Routine 

by Alan D. Matzger<br>960 Guerrero St.<br>San Francisco, CA 94llo

A few issues of Colorcue ago, Myron Steffy had an article introducing the CALL statement. I came upon its utility first when $I$ was writing a text editor for Assembler source code. Like Mr. Steffy, I wanted to move large blocks of memory from one place to another and BASIC's PEEKing and POKEing was unconscionably slow. Having solved the memory moving with an assembly language routine, $I$ went on to use the CALLed routine to list text, edit individual lines, search for particular strings, and insert, replace, delete, move, and copy one or more lines.

Which function was performed depended only on the particular $X$ in $A=C A L L(X)$. I have since used the CALL statement in the screen editor, in a "stolen" HP-85 Cribbage game to draw cards and evaluate hands and, as illustrated below, to sort any one of BASIC's string arrays.

The January 1981 Colorcue listed the labels for scratchpad addresses in alphabetical order. I wanted the listing in ascending order of the hex addresses. First, I copied each entry, interchanging the label with the hex location, and saved them in a .RND file with 137 records, each 50 bytes long. Using a simple program which read the records into a string array and then CALLed the sort routine, I was able to sort the 137 records in about 3 seconds.

The algorithm for the routine is the Shell-Metzner sort, modified to leave the array in place and exchange the entries' index numbers in a "pointer array" if the lower entry is "higher" than the higher one in the ASCII collating sequence. An explanation of the Shell-Metzner algorithm may be found in any of several computer science texts, and in articles in computer magazines, such as BYTE. My intent here is to demonstrate another way of using the CALL statement. A short explanation of the assembly language routine is in order, however.

When RUN through FCS, the routine is entered at the label START, where it sets up the CALL vector and adjusts BASIC's end of memory pointer to just in front of START itself, to avoid having the routine overwritten. Doing this within the routine itself cleans up the BASIC program, leaving only a mandatory CLEAR statement to be done there.

When CALLed, the routine is entered at the label MAIN with the value of $X$ in $A=C A L L(X)$ converted into two bytes and received in the DE register
 and it must be restored to HL upon RETurning. Tie two bytes in de upon RETurn are converted to a decinal number which is the value of A. It ay or may not be useful to the CALLing program; in tinis case it is used to pass back the address of the pointer array. Strictly speaking then, only one parameter may be passed and returned by the CALL statement. But as demonstrated by the BASIC demonstration program delow, the CALLing program can have access to any amount of data calculated by the assembly language coded routine.
(At the start of the routine, the pointer array INDX0 is set up in $1,2,3, \ldots$ order. At the end, the index number of the "least-valued" string will be in the first slot of this pointer array. The demo program reacis this array sequentially, converting the two bytes to a decimal number which is the index of the next higher string and prints it. Similarly, more than one value can be passed to the routine by POKEing into appropriate locations within the assembled code. One must be careful to note the addresses of wanted data areas from the listing by the assembler, convert them into decimal numbers, and then insert these numbers into the BASIC program.)

FNDAR finds the array name passed to it in BASIC's array table, and saves the number of entries and the address of the first string for the sort routine.

Finally, the CALLed routine need not have its own stack; I have consistently used whichever one BASIC uses and have never run out of PUSHing roon. The CALL statement, thus, is not nearly as restrictive as suggested in the Programing Manual. BASIC is fast enough for most input-output operations, and the use of assembly language routines inserted via CALLs greatly enhances its usefulness by adding speed when needed. CALLing machine language routines has made my machine even more enjoyable, fast and powerful.

## 22

## PNGE

 50| $;$ | ASSEMBLY LANGUAGE SORT |
| :--- | :--- |
| $;$ | Alan D. Matzger |
| $;$ | 950 Guerrero St. |
| $;$ | San Francisco, CA 94110 |

Establish run-time linkages START is entered when the PRG file is RUN by FCS ENTRY START
START:

## PUSH PSN

;save regs at entry
PUSH
LKI H,MAIN ; set up CALL vector with address,
SHLD 33293
MVI A, (JMP) ; and JMP op-code
STA 33282
LXI H,START-
SHLD 32940 ; set top limit of BASIC to below us
POP $H \quad$;restore regs
MVI $B, 0$; (except $B$, a 0 in which indicates $O K$ to $F C S$ )
POP PSS
RET
Shell sort of leave.
MAIN is entered when BASIC program executes a CALL
;
MAIN:

| PUSH | H | ; Save BASIC address |
| :--- | :--- | :--- |
| CALL | FNDAR | ;Subrtn to find list to sort |
| CALL | MFNDX | ; Initialize pointer array |
| CALL | SHSRT | ;The actual sort |
| LXI | D, INDXO | :Return the address of ptr array |
| POP | $H$ | :Retrieve return address |
| RET |  | ;Back to BASIC |


| V678 | FQU | 1 | : Select version by setting aprop. $V$ to 1 |
| :--- | :--- | :--- | :--- |
| $V 879$ | EQU | 0 | ; and others to 0 |
| V990 | EQU | 0 |  |

CALLable Shell-Metzner sort

|  | IF | V979 OR | V980 |
| :---: | :---: | :---: | :---: |
| $\begin{aligned} & \text { OSTR } \\ & \text { NEGH } \end{aligned}$ | EQU | 182AH | ; Puts out string to screen ; Negates HL |
|  | EQU | 195A |  |
|  | ENDIF |  |  |
|  | IF | V678 |  |
| OSTR | EQU | 33 FdH |  |
| NEGH | EQU | 3524H |  |
|  | ENDIF |  |  |
|  | PAGE |  |  |
| ; |  |  |  |  |
| ; | DE contain array name as 2.55 first ch + second ch +128 |  |  |
| ; | e.g. if array name is $A L S$, then $D=x^{\prime} A^{\prime}{ }^{\prime}, C=x^{\prime} C C^{\prime}$ |  |  |
| ; |  |  |  |
| ; | on exit, NAMEU --> ARRAY |  |  |
| ; | NELEM has \# elements |  |  |
| FNDAR: |  |  |  |
|  | LHYLD | BGARR | ; HL points to first array |
|  | MOV | B, D | ; Keep name in BC |
|  | MOV | C.E |  |
|  | PUSH | H |  |
| FNDR1: | POP | H |  |
|  | MOV | A, M | : look at first byte |
|  | INX | H | ;point to second |
|  | CMP | C | ; lobytes the same? |
|  | JNZ | FNDR2 | ; IF NOT, look at next entry |
|  | MOV | A, M | ; look at second |
|  | CMP | B | :hibytes the same? |
|  | JZ | FOUND | iyes. |
| FNDR 2 : | INX | H | ;point to next byte |
|  | MOV | E,M | ; these contain offset to next entry |
|  | INX | H |  |
|  | MOV | D, M |  |
|  | JNX | H |  |
|  | DAD | D | ;add offset |
|  | PUSH | H |  |
|  | XCHG |  | ;but maybe we're |
|  | DCX | D |  |
|  | LHLD | NDARR | ;beyond last entry |
|  | MOV | A, H |  |
|  | CMP | D |  |
|  | TM | FNDER | ; we are - show error |
|  | JNZ | FNTRIL | ; we're not - look at next |
|  | MOV | A, L | ;compare lobytes |
|  | CMP | E |  |
|  | JNC | FNDR1 | ;we're not $=100 k$ again |


| FNDER: | LXI | H, FNEMG |  |  | RAR |  | ; we are dividing by two |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | CALL | OSTR | ; display error msg |  | MOV | H, A | ; to get the partition factor |
|  | POP | H |  |  | MOV | A, L |  |
|  | POP | H |  |  | RAR |  |  |
|  | JMP | EXIT | ; nothing more to do |  | MOV | L, A |  |
| FOUND: | LXI | B,4 | ; in a one dimension list, |  | ORA | H | ;if it is zero |
|  | DAD |  | ; \# elements is 4 bytes away |  | RZ |  | ; we're all done |
|  | MOV | E,M |  |  | SHLD | PARTN | ; but we aren't |
|  | INX | H | ; get that number |  | CALL | NEGH | ; this \# is used |
|  | MOV | D, M |  |  | SHLD | NGPTN | ; in many subtractions |
|  | XCHG |  |  |  | XCHG |  |  |
|  | SHLD | NELEM | ;and store it away |  | LHLD | NELFM | ;as here |
|  | XCHG |  |  |  | DAD | D |  |
|  | INX | H | ;next byte is first for ${ }^{4} 0$ |  | SHLD | LPLIM | ; this is loop limit |
|  | SHLD | NAME0 | ; save that away |  | LXI | D, 0 | ; start with first string |
|  | RET |  |  | SSLP2: | INX | D | :get next index |
|  | PAGE, |  |  |  | XCHG |  |  |
| ; | MKNDX | - Make index array subroutine |  |  | SHLD | INNOI | ;store it |
| ; |  | During | the sort, the strings themselves are not |  | XCHG |  |  |
| ; |  | changed | ; their indices in a pointer array (INDX0) |  | LHLD | LPLIM | :is it > limit? |
| ; |  | are swi | tched. This subrtn initializes that array. |  | CALL | NEGH |  |
|  |  |  |  |  | DAD | D | ;compare them |
| MKNDX: | LXI | H, INDXO ; addr of 0th element |  |  | LHLD | PARTM | ; but load this before the test |
|  |  |  |  |  | JC | SSLPI | ; it is greater, goto loopl |
|  | LXI | D, 0 | ;index and value start the same | SSLP3: | DD | D | ;it's not, add partition factor |
|  | LXI | B,1280 | ; there are 640 entries MAX |  | SHLD | INNO2 | ; to get second index |
| MKNX1: | MOV | M, E | ; 2 bytes each |  | CALL | DTSTR | ;sr to get strimg indexes |
|  | INX | H |  |  | CALL | GT2ST | ; this gets their len and address |
|  | MOV | M, D | ; value stashed |  | CALL | CCMPR | :sr to compare two strings |
|  | INX | H | ;prepare for next |  | LILD | INNOI | ; load first index in case |
|  | INX | D | ; value is one more |  | XCIG |  |  |
|  | DCX | B | ;are we all done? |  | CPI | DOSNT | ; the two must be switched |
|  | MOV | A, C | ; we'll see |  | JNZ | SSLP2 | ; they don't, go back for more |
|  | ORA | B | ;not if result is <>0 |  | CALL | SWTCH | ; they do |
|  | TNZ | MKNXI | ; nope |  | LHLD | INNOl |  |
|  | RET |  | ;yup |  | XCIGG |  | ; ${ }^{\text {dNNOL }}$ now in DF. |
|  | PAGE |  |  |  | LHLD | NCPTN | ; lst str of next comparison is |
| ; | SHSRT - The actual sort routine |  |  |  | DAD | D | ; [NNO - PARTN, if that's |
| ; |  | INNO1 | INNO2 are the pointer array indices |  | MOV | $\mathrm{A}, \mathrm{H}$ | ; not zero or less |
| ; |  | NSTR1 \& NSTR2 are the values in the index |  |  | ANA | A |  |
| ; |  | and are themselves the indices to the two |  |  | , M | SSLP? |  |
| ; |  | strings in their own array. |  |  | JZ | MRTST |  |
|  |  |  |  |  | JP | RTTO 3 |  |
| SHELT: |  |  |  | MRTST: | ORA | L | ; is L zero too? |
|  | LHLD | NELFM | ; the number of strings |  | JN7. | RTTO3 | ;if not, goto LP3 |
| SSLPI: | MNA | A | :clear array |  | LhLD | PARTN | ; if so, goto LP? |
|  | MOV | A, H |  |  | JMP | SSLP2 |  |


| RTT03: | SHLD | INNO1 | ; stash new strl index |  | INX |  | ; stash this addr in DE |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | XCHG |  | ;put in DE |  | MOV | D, M |  |
|  | LHLD | PARTN |  |  | XCHG |  | ;put it in HL |
|  | JMP | SSLP3 |  |  | POP | D | ;retrieve NAMEO |
|  |  |  |  |  | RET |  |  |
| íSTR: | DTSTR - Obtain NSTR's from INNO's |  |  |  |  |  |  |
|  |  |  |  | ; | COMPR - The comparison routine |  |  |
|  | LXI | D, INDXO | ;start of index | ; COMPR: |  | If lst <= second, A returns FFH , else 0) 0 H |  |
|  | LHLD | INNO1 |  |  |  |  |  |
|  | LXI | B,NSTR1 | ;address, not the value |  | LHLD | ASTR2 | ; Point to its first byte |
|  | CALL | NFRIN | ; this gets and inserts the value |  | XCHG |  | ; In DE |
|  | LHLD | INNO2 | ;repeat for second |  | LHLD | ASTRI | ;addr of lst string |
|  | LXI | B,NSTR2 |  | COMP1: | LDAX |  | iget the byte |
|  | CALL | NFRIN |  |  | CMP | M | ;is str2 > strl? |
|  | RET |  |  | JM | OGT2 | ;no, it's less |
|  |  |  |  |  |  | JZ | OEC2 | ;no, it's equal |
| NFRTN: | DAD | ${ }^{1}$ | ; $\mathrm{HL}^{\star} 2$, each entry is 2 bytes | OLT2: | XRA | A | ;yes -- put 00 in A |
|  | DAD | D | ;points to value in index |  | RET |  |  |
|  | MOV | A, M |  | OGT2: | MVI | A, OEFH | ; tell caller to switch |
|  | STAX |  | ;store lobyte |  | RET |  |  |
|  | INX | H |  | OEQ2: | DCR | B | ;end of strl? |
|  | INX | B |  |  | JZ | OLT2. | ;yes, so 2nd > lst |
|  | MOV | A, M |  |  | DCR |  | ;end of str2? |
|  | STAX |  | ;store hibyte |  | JZ. | 0 OTP2 | ;yes, 1st > 2nd |
|  | RET |  |  |  | INX |  | ;point to next byte |
|  | PAGE |  |  |  | INX |  | ; ditto |
| ; <br> GT2ST: | GT2ST - Get length and starting addresses of the 2 strings |  |  |  | TMP | COHPl |  |
|  |  |  |  | PAGE |  |  |
|  | LHLD | NAMEO | ;addr of ptr to string $\# 0$ |  | $\begin{aligned} & \vdots \\ & \text { SNTCH: } \end{aligned}$ | Swich - Switch values in index array |  |  |
|  | XCHG |  | ;put in DE |  |  |  |  |  |  |
|  | LHLD | NSTR1 |  |  |  |  |  |
|  | CALL | GLNAD | ; this gets 'em | LXI |  | D, INDXO |  |
|  | MOV | B, C | ;len1 now in B, len2 will be in C | LHLD |  | NSTR2 | ;this value will go |
|  | SHLD | ASTR1 | ;addr returned in HL | PUSE |  | H | ; where NSTRI was |
|  | LHLD | NSTR2 |  | LHLD |  | INNOl | ; but we have to find |
|  | CALL | GLNAD |  | DAD |  | H | ; original address |
|  | SHLD | ASTR? | ; do same for 2 nd str | DAD |  | D | ;here it is |
|  | RET |  |  | POP |  |  | ; now the value is in BC |
|  |  |  |  | MOV |  | M, C |  |
| GLNAD: | DAD | H | ; each entry is 4 bytes |  | INX | H |  |
|  | DAD | H | ; so mult nstr by 1 |  | MOV | M, B | ;all moved, now for other |
|  | DAD | D | ;add to NAME0 |  | LHLD | NSTR1 |  |
|  | PUSH | D | ; this byte is len of str |  | PUSH | H |  |
|  | MOV | C, M | ; this one is filler |  | LHLD | INNO2 |  |
|  | INX | H | ; lobyte of addr in string space |  | DAD | H |  |
|  | INX | ${ }^{\mathrm{H}}$ |  |  | DAD | D | ;here's the address |
|  | MOV | E, M |  |  | POP | B |  |

CALLable Shell-Metzner Sort PAGE 7

|  | MOV | M, C | ;switch completed |
| :---: | :---: | :---: | :---: |
|  | INX | H |  |
|  | MOV | M, B |  |
|  | RET |  |  |
| ; | VARIABLES |  |  |
| ; |  |  |  |
| NELEM: | DS | 2 |  |
| INNOL: | DS | 2 |  |
| INNU2: | DS | 2. |  |
| NSTR1: | DS | 2 |  |
| NSTR2: | DS | 2 |  |
| ASTR1: | DS | 2 |  |
| ASTR2: | DS | 2 |  |
| PARTN: | DS | 2 |  |
| NGPTN: | DS | 2 |  |
| LPLIM: | DS | 2 |  |
| NAME0: | DS | 2 |  |
| DOSNT | EOU | UFFH |  |
| FNEMG: | DB | 5,1.3 | 20,5,237,50 |
|  | DB | ${ }^{\prime} \mathrm{LIS}$ | NOT FOUND', 11, 239,239 |
| INDX1): | DS | 1280 | ;THERE ARE 520 2-BYTF |
|  | END | STAR | EvTRTES |

CALL IN MSSEMBLY LANGUAGE SORT ROUTINE

10 PLOT 27, 4: PRINT "RUN SMSORT": PLOT 27, 27
20 CLEAR 20000
25 REM
SET UP TEXT ARRAY TO BE SORTED

30 DIM AS(100)
40 FOR $X=1 T 0100$
$45 \mathrm{SS}=$ ""
50 FOR L= 1 TO RND (3)* 20:REM LENGTH OF AN ITEM
50) $55=5 \$+$ CHRS (55+ 25* (RND (2)))
70) NEXT

75 AS $(X)=5 \$$
80 NEXT
85 REA
CALL SORT ROUTINE TO SORT AS
$90 \mathrm{X}=128+255 *$ ASC ("A")
95 PRINT "GOING TO SORT"
$100 \mathrm{~B}=$ CALL (X)
105 REM
PRINT SORTED ARRAY, USING POINTERS NT LXC. B

110 DEF FN $I(Z)=\operatorname{PEEK}(Z)+255 * \operatorname{PEFK}(Z+1)$
125 REM STEP THRU POINTER ARRAY, FACH 2 BYTFS LONG
130 FOR $\mathrm{I}=2 \mathrm{TO}$ 2* $^{\text {* }}$ lUOSTEP 2
140 $I X=F N I(I+B)$
150 PRINT AS (IX):NEXT


HOWARD ROSEN, INC.

```
    Put the finishing touches to sour Compucalor II or ISC computer.
Come up to the world of word processing.
> Extend the utilization of your computer to the other members of your family.
            x Letters
            x School reports
            * Busimess reports
            * If you now type-write it, Comp-U-Write it for a better product.
Basic requiremerits for CCII or 3651/9651:
            16K RAM.
            117key kesbosrd.
            Frinter.
            Comp-U-Writer software and instruction manual.
> For maximum capability:
            Full 32k RAM.
            Lower case characters.
> Talk to other computers: Add a MODEM to your system.
    We carry the entire CCII & ISC line of hardware/software, includirig spares.
    Send for our 4-page order form for hardware/software. Request separately
    by item your spare parts reeds.
            Send your order now. We pay the ghipping.
            Allow }5\mathrm{ weeks for delivery.
```



## Colorcue

## Editorial Offices

161 Brookside Dr.
Rochester, NY 14618

An Pr Publication

# Coloreue 



Dec/Jan 1982

## Colorcue

A bi-monthly publication by and for
Intecolor and Compucolor Users

Editors:
Ben Barlow David B. Suits

December 1981/January 1982
Volume 4, Number 3

## 3 Editors' Notes

## 5 Protected Fields, by Bernie Raffee <br> Control keyboard input with a machine language patch

13 Serial to Parallel Interface, by Ben Barlow A simple but useful project for your peripherals

19 Assembly Language Programming, by David B. Suits Part III: More on the 8080 instruction set

25 'The' BASIC Editor, reviewed by David B. Suits A powerful new utility for BASIC programmers

Advertisers: You will find our advertising policies attractive. Write for details.

Authors: This is a user-oriented and supported publication. Your articles/tips/hints are required to make it go. Send your articles or write for information.

Colorcue is published bimmonthly by Intelligent Systems Corporation, with editorial offices in Rochester, New York. Editorial and subscription correspondence should be addressed to the Editors, Colorcue, 161 Brookside Dr., Rochester, NY 14618. Product related correspondence should be addressed to ISC, 225 Technology Park, Norcross, GA 30092, ATTN: Susan Sheridan. Opinions expressed in by-line articles are not necessarily those of the editors or of ISC. Hardware/software items are checked to the best of our abilities but are NOT guaranteed.

## Editors' Notes

## Notes on Assemblers

The Macro Assembler is a very powerful tool when working with assembly language programs. It is, however, different from and a little more difficult to use than the old original Assembler program. The differences caused some trouble to a few of our readers who keyed in the Assembly Language Sort subroutine in the last issue, and then tried to "Assemble" it (rather than "Macro Assemble" it). We appologize to those readers we misled. A1though the Macro Assembler rates its own set of tutorial articles, the following brief notes will tell you what is needed to "translate" between the two versions of assembly language. Colorcue will often publish programs in the Macro version, since one of its editors threw away his old Assembler. [Tsk, tsk. --Other Ed.]

The most powerful feature of the Macro Assembler (hereafter known affectionately as Mac) is its support of relocatable, modular code. Instead of one mammoth source program, the Mac user can develop small modules, assemble them individually (faster editing, faster assembly, faster debug) and then "link" them together using the L80 linker. Mac doesn't know (generally) where the final PRG program will reside; the linker is told that. Mac tells the linker which labels must be known outside the current module by making available ENTRY and EXTRN commands. An ENTRY point is a name in a module that may be referenced by another module; an EXTRN is a name in another module that the current module wishes to reference. In our subrouines, replace an ENIRY START line with an appropriate ORG statement.

The Sort routine needed an ORG 0 FOOOH .

Another one of Mac's features is that of "conditional assembly". It's possible to define conditions within a source file, and then select one of several options with assembly-time "switches". This feature can be seen in the Sort routine where different EQUs are generated based on the setting of the version variables, V879, V980 or V678. To use with the regular Assembler, simply delete the statements for the version in which you're not interested, as well as any LF, FLSE, and ENDIF commands you may find.

Mac's third major feature is the ability to define and use macros. Since these are just notes, and macros haven't been used in any Colorcue programs to date, we won't discuss them now, but save them for the later article they deserve.

## Busses

We know that several hardwareoriented readers have developed peripheral devices and busses for the 50 pin bus. It would be advantageous to all of us--developers, users, would-be consumers--if we had a standard for an external bus interface. Colorcue (and some of its corresponding electrical engineers) will be happy to coordinate a stan-dard-setting activity. Just send a definition of your bus to our editors, and we'll chair a remote committee of respondents. We'1l also be happy to publish a dialog on the matter. Send us your comments, notes, articles.

## Back Issues

Please note that Colorcue back issues are for sale. See inside of this issue's back cover.

# LOOKING FOR CUSTOM KEYCAPS 

 for yourCOMPUCOLOR II Or INTECOLOR COMPUTER?


We stock blank keycaps that match the original keycaps supplied on ISC's keyboards

$$
\begin{aligned}
& 1 \times 1 \text { size in } 16 \text { colors } \\
& 1 \times 2 \text { size in } 14 \text { colors }
\end{aligned}
$$

Single quantity prices for one line engraved keycaps start at:
$\$ 1.17$ for $1 \times 1$ size with 5 characters
$\$ 1.95$ for $1 \times 2$ size with 10 characters

Other sizes and specific colors are available on special order.
Custom molded keytops available for high volume users.

> For order form call or write:
> Arkay Engravers

2073 Newbridge Road
Bellmore, NY 11710
516-781-9859

# Protected Fields 

by BERNIE RAFFEE

31 Malvern Avenue
South Harrow, Middlesex
England HAU 9EU

Although the common method of input to a BASIC program (prompted INPUT) is adequate for most applications, there are times when it is not so suitable. Consider, for example, a data entry application where unskilled users must enter names, addresses, and numbers from forms through a screen for a program to write to a disk file. It would be nice if INPUT would allow the programmer to exclude all the undesirable keys (the menu, escape, blinks, erases, home, or down arrow) that can cause such harm if pressed. Programmers themselves can press those keys by mistake, and trying to get new users to realize the importance of NOT striking the wrong key is sometimes a tough job. (And understandable, from their point of view. If the computer doesn't want to see a menu key in the middle of typing in an address, why doesn't it just ignore it?) Another drawback to INPUT is that the user can type in 26 characters for a 25 character field. Nothing stops them. The program can generate an error message after the fact, but it would be much neater to inhibit input after the 25 th character.

This article describes a CALLable assembly language subroutine that does just that. In addition, it can disallow entry of alpha data in fields defined as numeric, and has editing features such as the use of the insert/delete character and delete field (line) keys. There are two programs presented in this article. One of them, the assembly language subroutine, will not be described. If you're interested, it makes interesting reading; if not, just type it in and assemble it. The second program is a BASIC program that illustrates the use of the CALLable "FIELDS" assembly language routine. The assembly language program can be assembled by either assembler (see the Editors' Notes in this issue about the differences between the macro and regular assemblers), and that when RUN from FCS, it establishes all necessary linkages, and adjusts for different software versions.

The major divisions in the BASIC program should be evident. Lines 550-870 display the initial screen. Each input field is shown as a blue rectangle of appropriate size. For example, the first field - NAME - is displayed on line 10 . The title of the field, NAME, identifying the data item, is displayed followed by a blue rectangle 12 bytes long. (Lines 630, 680.) Lines $880-1080$ are the main body of the code, where each field is "read" from the screen. Lines $880-920$ read the first field, NAME. Variables $X$ and $Y$ are set to the first position of the blue rectangle for NAME, ML is set to 12 , the maximum field length, NAS is set to " $A$ " indicating an
alphanumeric field, and the CCI set to white on blue. The control subroutine at line 100 (through 195) is called to interface to the machine language subroutine. The 100 subroutine POKEs parameter values for the assembly language routine into a known common location (right after the end of BASIC RAM), and CALLs the routine. On return, FI\$ is constructed from the field, byte by byte, and control returned to line 910 . The variable $K B$ is used to return the ending character of the field, and if a TAB key was pressed, the program "backs up" a field. Lines 1210-1220 are where the main field processing program would be. In this demo program, the entered field is simply written back on the same line along with its length.

When the assembly language program gets control after the CALL, it takes over all keyboard input and disallows any control characters and any "undesired" characters. It processes editing keys, such as cursor controls and character insert keys. The delete line key causes the field to be cleared from the cursor to the end of the rectangle. When the return key signifies the end of field entry, what is shown in the field will be what the program gets, unlike INPUT, which inputs only to the cursor. (In the assembly language program, changing the value of ALPHMAX to 126 will permit lower case letters.)

Although the routines here may be a little long, I'm sure you'll find them to be quite useful. They can greatly extend the use of your computer by making the "feeding" of it more friendly. ©
(Program listings begin on next page.)

## CALL FOR ARTICLES

Colorcue gets its material from those who write it. It don't grow on no trees. Nevermind your bad sppeling or badly grammar: that's the job of the Editors. You come up with the ideas, splash them onto paper (or disk), and we'll wrestle them into an intelligible form for publication. We can't pay you for your time--you won't become rich. But, since Colorcue is read all over the world (well, not all over, exactly), then maybe you'll become world famous.... We're looking for ideas, programs, reviews of books and/or software and/or hardware applicable to ISC machines, hardware/software modifications, user group news, and.... Well, you get the idea.

FOR SALE. 32 K Compucolor II with 117 key keyboard and lower case. Version 8.79. Centronics 779-2 dot matrix printer with lower case and interface/cable. Novation CAT modem with cable and Com-Tronics TERMII communications software. Disks: Sampler, Basic Editing, Assembler, Text Editor, Star Trek, Personal Data Base, Blackjack, Cubic Tic-tac-toe, plus 25 additional disks. Compucolor Programming Manual. \$2200.00. Contact: Don Miller, 112 Marble Drive, Rochester, NY 14615 (716) 663-1175.

;THIS CHARACTER INPUT INITIALISATION
;ROUTINE SETS UP THE PARAMETERS
;NECESSARY FOR THE 'CHRINT' AND
;'CI' ROUTINES.
CIINIT: PUSH PSH ;SAVE REGISTERS REQ'STD BY BASIC
PUSH H
LDA KBDFL SAVE BASIC'S JUMP VECTOR
HVI A,31 ;SETUP NEL JUMP VECTOR
STA KEDFL
MUI A,OC3H
STA PLACE 'JMP' AT VECTOR LOCATION
IMPCRT
LXI H,CHRINT ;GET ADDRESS OF 'CHRINT' ROUTN
SHLD INPCRT+1 ;FLACE ADDRESS AFTER 'JHP'
HOU B,A
STA KBRDY
STA CHARIM
STA COUNT ;RESET CHARACTER COUNT
JMP GETNXT

- this character input routine is ; VECTORED TO FROM THE KEYBOARD INPUT ;ROUTINE through the jump vector ( (31). ; THE CHARACTER FROM THE KEYBOARD ;ROUTINE IS IN REGISTER 'g'.

CHRINT: LKI H,CHARIN
GEET ADDRESS OF TEHP
;CHARACTER STORAGE.
*RA A ;CLEAR ACCU犃LATOR
CMP M ;TEST FOR 'CHARIN' FOR IERI
JHZ CFIN ;IF NOT ZERD THEN IGMORE INPUT
MOV $A, E$ :GET CHAR FROM 'E'
AMI 127 ;STRIP UPPER BIT FOR ASCII
MOV M, ;PUT CHAR IN 'CHARIN'

| CFIN: | $\begin{aligned} & \text { EI } \\ & \text { RET } \end{aligned}$ | EENABLE INTERRUPTS <br> ; RETURN FROM INTERRUPT |
| :---: | :---: | :---: |
|  | ; Cl - | this character infut routine gets a character |
|  | ; | FROH THE TEMPORARY STORAEE LDCATION 'CHARIN' |
|  | ; | Clears the keyboard reany flas' |
|  | ; | AND RETURNS WITH THE CHARRCIER IM 'A': |
|  | ; | IF THERE IS NO CHARACTER : 'CHARIN'; |
|  | ; | THEN 'CI' hILL HANG ANU HAIT FOR ONE. |
| CI: | EI | ; ENABLE INTERRUPTS |
| C110: | LDA | Charin ; GEt character |
|  | CPI | 0 ¿HAVE Á CHAR? |
|  | 32 | CIIO ; IF NOT, HANG F-T ONE |
|  | PUSH |  |
|  | XRA | A |
|  | STA | kbrdy ; CLEAR KEyplard ready flag |
|  | STA | CHARIN ;CLEAR TEMP STARAEE FOR NEYT CHAR |
|  | POP | PSY \# ; RESTORE CHAR |
|  | RET |  |
| GETNXT: | CALL | CI ;GET NEXT KEYBOARD CHARACTER |
|  | MOV | D,A :STORE IN 'D' |
|  | CPI | LEFT ;BACKSPACE? |
|  | d2 | LEFTKEY : JuMP TO BACKSPACE ROUTIME |
|  | CPI | HT ;TAB? |
|  | J2 | MAIWFIN ; YES-JUMF TOE END ROUTIME |
|  | CPI | CR PRETUR ? |
|  | 31 | MAINFIN :YES - JUMP TO EMD ROUTT RE |
|  | CPI | RIGHT ;CURSOR RIGMT |
|  | J2 | RGhtkey ;Put It on Screen |
|  | CPI | INSERT ; INSERT A NEL CHARACTER |
|  | Jl | INS10 |
|  | CPI | DELETE ;DELETE A CHARACTER |



```
STA COUNT
MOU E,D :PUT LAST CHAR IN 'E'
MRA A
MOV D,A
POP PSA ;RSTORE BASIC JUMP VECTOR 
STA KBDFL
POP H {RESTORE USED REGISTERS
POP PSH
RET
LEFTKEY: MOV A,B ;BACKSPACE SUEROUTINE
CPI 0 ;IF ON FIRST CHAR
JZ GETNXT ;THEN IGMORE
DCR B ;DECREMENT CHAR COUNT
mVI A,LEFT ;BACKSPACE CURSOR
CALL LO
JMp GETNXT
RGHTKEY: MOV A,g ;RIGHT CURSOR FUNCTION
    CMP E ;CHECK IF FIELD FULL
    MOV A,D :RESTORE CHARACTER
    JNZ KEYOK ;NO SO O.K.
    JMP GETNKT ; IGNORE
    ;INSERT A CHARACTER ROUTINE....
    ;REGISTER USAGE:-
    ; A=NO OF CHARS ALREADY IN FIELD
    ; B=POSITION LITHIN FIELD
    ; C=NO OF CHARS TO SHUFFLE
    ; E=MAX ND OF CHARS IN FIELD
IHS10: CALL CNTIO ;DETERMINE SIIE OF FIELD RTN
    CPI 0 ;IGNORE IF ND CHARS IN FIELD
    JI GETNXT
    PUSH D
```

|  | LHLD | FLDADR | ; SET UP HL |
| :---: | :---: | :---: | :---: |
|  | DCX | H |  |
|  | DCX | H |  |
|  | SUB | B | ;CALCULATE NO OF CHARS |
|  | JM | IMS55 | ;RETURN IF PAST LAST CHAR (NAUGHTY:) |
|  | JI | IN555 |  |
|  | W0V | C, A | ;PUT IN C |
|  | ADD | B | ;RESTORE A |
|  | CMP | E | ; IS FIELD FULL |
|  | JNZ | INS20 | ;NO |
|  | DCX | H | ;YES - SO ARRANGE FOR LAST CHAR |
|  | DCX | H | ; TO DISAPPEAR FROM FACE OF EARTH |
|  | DCR | C |  |
|  | 32 | IMS55 | ;IGNORE IF SITting on last char |
| ING20: | MVI | D,0 | ;POSITION HL TO LAST CHAR |
|  | Moy | E, A |  |
|  | DAD | D |  |
|  | DAD | D |  |
|  | LXI | D, OFFFCH | : TO Subtract 4 later on |
| INS30: | MOV | $A_{\text {, }}, 4$ | ;GET CHAR FROM FIELD |
|  | INX | H |  |
|  | IN\% | H |  |
|  | Hov | H, A | AND SHIFT UP |
| INS40: | ICR | $\bigcirc$ | :ANY MORE? |
|  | J2 | INS50 | ; NO |
|  | DAD | 0 | ;YES POINT TO PREVIOUS CHAR |
|  | JMP | IN530 |  |
| IN550: | DCX | H | ;PUT A Space at cursor position |
|  | DCX | H |  |
|  | mVI | M, SPACE |  |
| INS55: | POP | D |  |
|  | JMP | GETNXT | ;RETURN |



|  | DAD |  | ;yEs - point to next one |
| :---: | :---: | :---: | :---: |
|  | JMP | DEL30 |  |
| DEL50: | INX | H | ;PUT A SPACE AT LAST |
|  | INX | H | ;CHARACTER POSITION |
| DEL53: | HVI | M, SPACE |  |
| DEL55: | P9P | D |  |
|  | JMP | GETNXT | ;RETURN |
|  | ;ROUT <br> ;CHAF | $\begin{aligned} & \text { E TO DELI } \\ & \text { TERS IN } \end{aligned}$ | ete all the rembining THE FIELD |
| DLINIO: | CALL | CNT10 | ;GET NO OF CHARS IN FIELD |
|  | PUSH | 0 |  |
|  | LHLD | FLDADR | ; START OF FIELD |
|  | SUB | B | ;CALCULATE MO. OF Chars to delete |
|  | JH | DLIN40 | ;igNore If at end |
|  | J2 | DLIN40 |  |
|  | MOV | C, A |  |
| DLIN20: | NVI | D, 0 | ;SET HLL TO IST CHAR TO DELETE |
|  | MOV | E, B |  |
|  | DAD | D |  |
|  | DAD | 0 |  |
|  | MVI | A, SPACE |  |
| DLIN30: | MOV | $M, A$ | ; DELETE CHARS |
|  | DCR | C | ;ANY MORE? |
|  | 31 | DLIN40 | ; NO |
|  | INX | H |  |
|  | INX | H |  |
|  | JMP | DLIN30 |  |

$\begin{array}{rll}\text { DLIN40: } & \text { POP } & 0 \\ \text { JMP } & \text { GETNXT }\end{array}$
; ROutine to determine the exact length ;OF A FIELD ON THE SCREEN.
:FINAL COUNT IS PUT IN 'A'.
CNTIO: PUSH 8 ;SAVE IT
LHLD FLDADR ;START OF FIELD
DCX H
DCX H
WI $\quad 8,0$
CNT20: MUI C,O :C HILL CONTAIN NO OF SPACES ;PAST THE END OF THE FIELD
CNT30: HOU A,B ;B CONTAINS THE FIELD COUNT
GMP E ;ARE HE AT MAXIMUM
JI CNT40 ;YES - FINISHED
INK H EGET TO NEXT CHAR
INX H
INR B
MOV A, H
CPI SPACE IIS IT A SPACE
JNI CNT20 : NDPE
INR $\quad$ IINCREMENT SPACE COUNT
JMP CNTJO ;BACK : SHUFFLE NEXT OME
CNT40: MOV A,B ;CALCULATE EXACT COUNT
SuB C
POP B ;RESTORE USED REGISTER
RET

END START

Listing 2. A BASIC program which demonstrates the fields machine language patch. Note: numbers in brackets are control codes entered from the keyboard. Thus, [16] is black; [17] is red; [29] is foreground on; etc.

[^0]660 PLOR 3,1,19:PRINT
679 IP AS ${ }^{*}{ }^{\prime \prime}$ GOID 730
680 PLOT 3,14,10:FRNNT " [30] [20] ..... [16] $^{\text {[16] }}$
699 RLOF 3,14,13: FRDNT "[2[29]
16]
16]
710 RLOT 3.14.19:PRINT - [2 710 RLOT 3.14
30 불FLOT 3.14.10:FRINT "[29][22][30][20]BERNIE RAFFE[16]*
4 ELOT 3.14:13:FRTNT -[29]HARRGN, ENGLAND [16]"
750 ELOT 3,14,16ะPRINT - [29]534.43-[16]"
750 RLOT 3,14,16:PRINT [29] 534.43-[16]"
769 RLOR 3,14,19:PRNNT [29]AVERAGE[16] [29]*
770 LLOT 3.0 .22
770 ELOT 3: 0.22 .
780 PRDNT " [22]SUNARY OR FACILITIES:-[18]"

810 PRINT [19] 1) CONTROL CHARACIERS \& CURSCR FOSITIONING KEXS (EXCEPT[18]-
810 PRINT "[19] LEFTT \& RIGHR ORSOR) ARE DISABLED. [18]"
830 PRNNT - [19] 3) FTNAL FIELD VALIE DEPENDS CN CNVIENT OF FIETD ON THE[18]-

840 PRINT [19] SCREEN \& NOT ON THE ACIUAL KETY DEPRESSIONS[18]*
850 PRINT' [19]4) 'DELEIE/DNSERTS CHAR' \& 'DETENE LINE' WORR WILHIN AFIELD[18
860 PRINT "[19] 5) ALPHABETHC CHARACIFRS ARE DISABLED IN NUMERIC FIELDS[18]"
870 ERNTI [19] 6) THE 'TAB' KEY REIURWS ORSOR TO THE FREVIOUS FIELD[18]"
875 REM [16][10]GET EACH FIELD[10][10]
$889 \mathrm{X}=14: Y=10: \mathrm{ML}=12: \mathrm{NAS}={ }^{-1} \mathrm{~A}^{*}$
890 PLOT 6,38
900 gasub 100
910 IP KB=9COIO 900
920 Gasus 1290

940 LLAT 6,38
950 GOsus 19A
960 IF KB=9GOID 88
978 GOSUB 1200
$980 \mathrm{X}=14: \mathrm{Y}=16: \mathrm{ML}=7: \mathrm{NASO}^{-N *}$
990 FLOT 6,38
1000 casub 100
1010 IF KB=9GOIO 939
1020 GOSUB 1298
$1030 \mathrm{X}=14: Y=19: \mathrm{ML}=7: \mathrm{NAS} \mathrm{man}^{-1}$
1045 FLOT 6,38
1050 GOSJB 190
1060 IF KB=96010 980
1676 GOSUB 1209
1089 HLOT 3,15,21:INPUT " [17]HIT REIURN FOR AROIBER GO [18]";AS

$1100601058 \%$
1290 REM [10][10]FIEND PROCESSING ROUITNE. THIS JUST PRINIS VALUE; [10]REAL PRO
GRAM WOULD PROCESS FIELD EERE
1210 FLOT 6,2,3,33,Y:FRTNT"[29][17]":PI\$;" ([29][22]";LEN(FI\$);"[29][17]) [29][
18] " ${ }^{\circ}$
1220 REIURN

## Serial to Parallel Interface

by Ben Barlow

The Epson line of printers is certainly one of the most popular ones in the personal computer industry. Their combination of features (including high resolution bit-level graphics), high reliability, and low price are hard to beat. Have you noticed the low-end printers being offered by IBM and HP on their lines of personal machines? OEMed Epsons. Many ISC and Compucolor owners have added Epson printers to their computers, and several of them, as I did, were probably surprised to find that the Graphics package (GRAFTRAX-80) does not work with the low cost serial interface board option in the printer; the $\$ 150$ buffered serial interface is required for GRAFTRAX. Having chosen the Epson at least partly because of price, and badly wanting the graphics option, I was distressed to think of $\$ 150$ fleeing after the $\$ 75$ (inexpensive interface) after the $\$ 90$ (graphics). At that rate, the options would soon exceed the basic printer's cost. This article is about a cheaper solution--one that you might want to consider as an alternative to either of the serial boards: a serial to parallel converter that allows the computer (serial) to talk to the Epson in its own tongue (parallel). (After writing this article, $I$ came across an article in the Sept. 1981 issue of BYTE magazine by Steve Ciarcia on the construction of an unlimited-vocabulary speech synthesizer that connects to any computer through a parallel interface. So there may be value in this article even if you don't have a printer.)


Photo 1.

The converter is a low cost device based on a 5 V UART which accepts data at 9600 BPS (Bits Per Second)from the computer, turns them into a parallel byte, and strobes them into the Epson's parallel interface. It accepts the BUSY signal from the printer and passes that back to the computer as not clear-to-send to control data flow. It's constructed with only four IC chips and a handful of other parts, and with judicious shopping and no junk box to pull from, it can be built for about $\$ 45$. With any kind of parts box to paw through, the cost can go down significantly. Construction is relatively simple, and the converter can be a rewarding project for the winter months.

Looking in more detail at the operation of the converter, it may help to look at the schematic. The AY-3-1014A UART is one of several chips which are adequate; choosing one that operates on a single 5 volt supply simplifies the process, though those chips are more expensive than the +5 and +12 ones. The 1014 is set up with its selectable options determined by connecting option selecting pins $34,35,36,37$ to +5 V . Pin 21 must be grounded, and pin 4 must be low to place received data on the output lines. Bits of a byte a clocked into the 1014 under the control of the Receiver Clock ( $\mathrm{RCP}=$ pin 17). When a serial byte has been completely received, the 1014 puts it onto the output lines RD1 through RD8, and raises the DAV signal on pin 19 to indicate that a byte is ready. DAV is used to generate a short duration pulse which strobes the data on the output lines into the printer's data buffer. The 74123 chip (a 74121 would have been more appropriate, but I didn't have one) provides this pulse, which is about 500 nanoseconds long.

The printer returns the ACKNLG* ("*" indicates active low) signal which is routed to the 1014 on pin 18, RDAV*. The 555 timer chip provides the clock for the 1014, which must be at 16 times the baud rate, or 153.6 K ( $9600 \times 16$ ). Although this is outside the 555's specification (100K is its published limit), it seems to produce an adequate (though not quite square) pulse train. (I experimented with CMOS oscillators with no luck. A baud rate generator chip, which would be the professonal way to get a clock, would have added about another $\$ 15-\$ 20$ to the low-cost project.) An adjustable potentiometer adjacent to the 555 provides the necessary degree of adjustment of clock frequency to compensate for temperature or burn-in. (My clock seemed to drift for a couple of days until the capacitor apparently stabilized. Since then it's been fine, and unadjusted.) The 1489 chip accepts RS232C signal levels, converting them to TTL levels and inverting them. Serial data comes from the computer at an RS232C level, and is routed through the 1489 before feeding to the 1014. The BUSY signal from the printer also is routed through the 1489 , but only for its inverter function. (The BUSY* signal thus supplied works with ISC computers and Compucolors containing the ECN 002137 published in the Aug/Sept Colorcue.)

The converter consists of three separately constructable pieces: the interface itself, a 5 V power supply, and a printer cable. The interface pictured (photo 1) was built on a small piece of vector board using wire wrap methods. A 16-pin socket was wired in as the parallel cable connector, allowing the cable to be removed from the interface itself, a simplifying factor during construction. I've gotten into the habit on wire-wrap projects of providing labelled test pins for +5 V and GND, so that I can easily attach logic probes or meters without crowding the components needlessly. The potentiometer for the 555 should be a multi-turn one, possibly with a series resistor as shown in the schematic, depending on the values you can locate. Having the adjustment screw available from the outside (notice the hole in the case in photo 2) simplifies adjustment somewhat. The power required is low, but I chose not to pull it from the computer. It's not readily accessible there, nor did it seem esthetically correct. My junk box had some non-working calculator power supplies that came from a surplus store as fantastic bargains a year or two back, and modification of one of those proved to be easy. For someone without that type of junk, a 6 or 9 V transformer with regulation to 5 V on the converter board would be an inexpensive solution. There's room on the board for the required regulator and capacitor.


Photo 2.
Cable construction is straightforward. The special connector for the parallel interface is soldered on one end, a $16-$ pin DIP header on the other. The only trick is to carefully map pin to pin. Although I used a hunk of 12 conductor wire from under the workbench, ribbon cable or separate wires should work as well. Remember to connect pins 35 and 31 on the parallel connector.

Once built, you'll want to check the wiring (and, if you're like me, recheck it). An easy way to do this is to check (with no chips in the sockets) each pin, each wire, with a voltmeter set to measure resistance-make a series of continuity checks. If all the wiring checks out, check the power supply wiring, and if it looks all right, plug it in and hope for 5 V . Check power and ground pins on each socket on the board, check the cable, and with the power off, plug in the chips. If it seems to be working (as opposed to smoking), check the 555's output with an oscilloscope. (If you don't have one, don't worry. The following procedure works as well.) Set the pot so that the $555^{\prime}$ s pulse width is 6.5 microseconds. Or, write a small BASIC program, such as:

```
10 FLOT 27,13
20 PRINT "ABCDEFGHIJKLMNOPQRSTUVWXYZ//////TEST....TEST.
*.TEST----TEST----TEST"
30 GOTO 20
```

Connect the computer's RS232 cable to the interface, connect the printer, power everything on (all the connections were made with power off, right?), and RUN the program. Adjust the pot until the printer prints text, not garbage. Find the middle of the "print" range, in between garbage settings, and leave it there. You may have to adjust it again, but if it acts like mine, it will settle in like a rock after a day or two.

In summary, if you are willing to trade a little time for money, and enjoy (or want to experiment with) the hardware side of the business, you'l1 find this project appealing. It's a simple, low cost, safe interface.

## Parts List

```
    Small box
    2.5" x 5" Vector board
1 8 pin WW socket
1 14 pin WW socket
2 16 pin WW socket
1 40 pin WW socket
l LM555 IC
1 MC1489 IC
1 74123
1 AY-3-1014-A
2 50 pF capacitor
2.1 mF capacitor
l 10 mF capacitor
1 1K Ohm resistor
1 20K Ohm resistor
1 56K Ohm resistor
```

150 K multi-turn pot.
1 DB-25 female soldertail connector misc. wire (stranded 22 ga. \& WW) .025" square post terminals

1 Amphenol 36 pin connector for Centronics-type printers
1 length 12 conductor cable (or substitute)
116 pin DIP header

5 V power supply or: 6 or 9V DC transformer 78 M 05 regulator 1000 mF 35 V capacitor


Serial to Parallei Converter

## RENAISSANCE MARKETING ANNOUNCES

Low Cost Business Programs<br>for your Compucolor model 4 , model 5 and Intecolor 3600 series Computers

GENERAL LEDGER* (16K, 32K) ....... $\$ 59.95$
(1) ACCOUNT PROGRAM: Display account data, list all accounts, add account, delete account, change account data.
(2) JOURNAL PROGRAM: Enter journal data, display journal data by entry number, charge account data.
(3) PROOF PROGRAM: Display or print journal entry proof sheets with titles.
(4) POST PROGRAM: Applies journal entries to accounts.
(5) REPORT PROGRAM: Print balance sheet, print income statement with titles.

INVENTORY CONTROL* (16K, 32K) ........\$ 34.95
Gives the following reports:
(1) Display item data by quantity.
(2) Print or display all items on file.
(3) Display item data by class code.
(4) Print or display item data by vendor.
(5) Updating section:
a. add new item
b. update item quantity
c. change item data d. delete item
(6) Display item data by item number.

Provides quantities and data for 750 items or models.
MAXELL MINI DISK for Compucolor. Box of 10 ..... $\$ 34.95$
$30 \%$ OFF all Compucolor Corp. software in stock. CALL

RENAISSANCE MKT.
7 So. Pierson Rd.
Maplewood, NJ 07040
201-762-0585
[ ] General Ledger \$59.95
[ ] Inventory Control \$34.95
[ ] Maxell Disk Box Qty. $\qquad$
Terms: payment with order Available for Intecolor Feb., Mar. 1982 Taking orders now.

# Assembly Language Programming 

PART III: MORE ON THE 8689 INSTRUCIION SEI by DAVD B. SUIHS

Last time we used ICS's Machine Language Debug Package in order to learn about a few of the 8080 machine language instructions (or "operation codes", or just "op codes", as they are often called). We saw how to put an eight bit number into the accumulator and then CALL a certain subroutine contained in your computer's ROM in order to put something onto the screen. This was just like BASIC's PLOT statement. In BASIC, you can specify a whole string of PLOT numbers just by separating them with commas. For example, PLOT $2,0,0,242,127,0,127,127,0,127,0,0,255$ will draw a line around the screen. Being able to separate the numbers with commas makes things a lot easier than having to say PLOT for each number: PLOT 2: PLOT 0: PLOT 0 : PLOT $242 \ldots$... etc. Can we do something like that in assembly language so that we don't have to load the accumulator each time with the next by te and CALL the subroutine? Yes. We will investigate a method which works something like BASIC's DATA and READ statements.

As we learned last time, the MOV instruction tells the 8080 to move (or, really, copy) the contents of one register into another. The destination register is specified first, and then the source register. For example, MOV A, E will copy the contents of the E register into the $A$ register (the Accumulator). The diagram shows the registers. But there is

a register (well, sort of a register) which is not shown: register M. 'M' stands for Memory. Although memory locations are surely not registers in the 8080 chip itself, the 8080 is cleverly designed so that we can sometimes treat them as if they were. But certain conditions must be fulfilled. Specifically, we can call that memory location whose address is contained in the $H L$ register pair a register. Thus, MOV $A, M$ is an instruction to copy the contents of the memory location into the accumulator. Which memory location? Why, the one whose address is contained in the HL register pair.

For example, if the HL pair contained the 16 bit address 0011111100001110 (which is $3 F 0 E$ in hex), then MOV $A, M$ would copy the contents of memory location $3 F O E H$ into the accumulator. The instruction MOV M, E would copy the contents of register $E$ into memory location 3 FOEH . And so on. There are other ways of getting data to and from memory, but this is a handy one which will suit our present purposes.

How do we set up the HL pair so that it contains the address of the memory location we want? We can use the MVI instruction which we dealt with last time. Remember, MVI <reg>, <num> will put <num> into 〈reg>. If we want 3 FOEH in the HL pair, then we could write:

| MVI | $\mathrm{H}, 3 \mathrm{FH}$ |
| :--- | :--- | :--- | :--- | :--- |
| MVI | $\mathrm{L}, \emptyset \mathrm{EH}$ | or $\quad$ MVI $\quad \mathrm{L}$, QEH

(Note that $I$ have explicitly specifed that each of the numbers is hex by using the capital 'H'. Although the MLDP assumes that all numbers are in hex, your assembler doesn't: it assumes that everything is in decimal unless followed by 'H' for hex. Since we will be dealing with the assembler eventually, we might as well get used to this way of doing things. Besides, specifying 'H' when hex is meant will help avoid confusion. Of course, for numbers less than 0 AH , it won't matter: 1 decimal $=1$ hex, 2 decimal $=2$ hex, and so on up through 9.) But there is a more convenient method of loading a register pair with two bytes, and that is: LXI <reg pair>, <num>. The 'LXI' stands for Load eXtended Immediate. 'Load' in this case is the same as 'move'. So this is just like the MVI (MoVe Immdiate) instruction, except that 'extended' refers to a register pair instead of just one register. The number which is to be loaded into the specified register pair will be interpreted as two bytes. Thus, LXI H, 3 FOEH is an instruction to load the register pair $H L$ with the two bytes of immediate data, $3 F 0 E H$. And LXI H,2 will load 0002 H into the HL pair. (The 'H' in the LXI H,... instruction will not be confused with the single register $H$, because the LXI instruction requires reference to a register pair. 'H' in this case stands for ' $H L^{\prime}$. ' $\mathrm{B}^{\prime}$ would stand for ' BC ', and ' D ' would stand for 'DE'. 'SP' would stand for 'Stack Pointer'. The Accumulator and FLAGS together constitute a register pair only in certain circumstances, and this is not one of them.)

That's fine, but we need some more tools in order to do what $I$ want us to be able to do. Enter the increment command. Any register, or register pair, may be incremented by one by the use of a single command. To increment a single register, INR <reg> will do. For example, if register B contains 4, then after INR B it will contain 5. If the Accumulator contains lllll111B ( $=0 \mathrm{FFH},=255$ decimal), then what will be in the Accumulator after INR A? $0 \mathrm{FFH}+\mathrm{l}=100 \mathrm{H}$, which is 100000000 binary. But that's 9 bits long, and the

Accumulator, like all the other single registers, can hold only 8 bits. As a result, the left-most bit is lost; the Accumulator will now have zero in it.

If we wish to increment a register pair, the INR instruction just won $^{8} t$ do. Suppose the HL pair contains $00 F F H$, that is, $H$ holds 00 and $L$ contains 0FFH. Then INR $L$ will result in $H=00$ and $L=00$. But if we want to jacrement the $H L$ register pair, then we want $00 \mathrm{FFH}+1$ to yield 0100 H ; that is, $H$ should contain 01 and $L$ should contain 00 . In this case we use the INX instruction: INcrement eXtended. Thus, INX H will increment the register pair HL; INX $B$ will increment the register pair $B C$; and so on. (Corresponding to the increment and increment extended instructions, there are the decrement and decrement extended instructions. DCR A will subtract one from the Accumulator. DCX $D$ will subtract one from the register pair DE. And so on.)

Where does all this get us? Well, now we can store all our PLOT numbers in a section of memory. Then we can load the $H$ register pair with the address of the start of that section of memory. That is, the HL pair will act as a pointer. Each time we wish to read a number from memory into the Accumulator, we use the MOV A,M instruction. Remember, that instruction gets the contents of the memory location which is pointed to by (i.e., whose address is in) the HL register pair. And then to output what is in the accumulator to the screen, we do what we did last time: CALL 33H. To get the next number, we simply make $H L$ point to the next memory location by using the INX H instruction. Then we loop back to the MOV A, M. And so on. Something like this:

```
MON A,M ;Get the number from memory.
CALL 33H ;Put it onto the screen.
INX H ;Point to next location in memory.
; (Now loop back to the MOV A,M instruction.)
```

(Notice, by the way, that your assembler will treat a semicolon just as BASIC treats "REM".)

Of course, we need something more. We need a LOOP. We want to execute those three instructions a certain number of times. How do we make a loop? We need a counter to keep track of the number of times we have gone through the loop. And we need a conditional branch instruction to go back to the start of the loop in case the counter hasn't yet finished counting. In BASIC, the loop is handled in the familiar way using a FOR-NEXT structure. Unfortunately, the 8080 doesn't understand either 'FOR' or 'NEXT'. How could we make a loop in BASIC if we didn't have the FOR-NEXT structure? Perhaps like this:

| 10 COUNI=1 | :REM INITIALIZE COUNIER. |
| :--- | :--- |
| $2 \emptyset$ blah | :REM HERE'RE THE STATEMENTS |
| 30 blah | :REM WHICH OCCUR INSIDE |
| 40 blah | :REM THE LOOP. |
| $5 \emptyset$ COUNI=COUNI +1 | :REM INCREMENT COUNIER. |
| $6 \emptyset$ IF COUNT<8 THEN $2 \emptyset$ |  |

In that example, we repeated a loop 7 times. Alternatively, we could count down instead of up:

10 COUNI=7
20 blah
30 blah
40 blah
50 COUNT=COUNT-1
$6 \varnothing$ IF COUNT>Ø THEN 20

As it happens, this is a perfectly natural and easy way for the 8080 to handle a loop. For the counter, we could use, say, register C. We initialize register C to 7 with MVI C,7. And we decrement the counter each time through the loop with DCR C. But how do we test for zero and go back to the beginning of the loop? In BASIC, you jump around with GOTO statements. If you want to GOTO some place unconditionally, you simply write GOTO xxxx. But if you want to GOTO some place only when a certain condition is satisfied, you have to put in a test: IF such-and-such THEN GOTO xxxx. The 8080 has its unconditional GOTO instruction, too, only it's called a JuMP instruction. JMP xxxx will cause the program to go to (jump) to address xxxx. (It jumps to an address in memory, rather than to a line number.) But in addition, the 8080 has quite a few conditional jump instructions, all of which take place automatically by testing the status of one or another of the bits held in the FLAGS register. One of those bits is called the zero bit, and that is set (=1) whenever the result of certain operations end up as zero. Some operations affect that flag bit, and some do not. (The DCR instruction does.) The conditional jump instruction we want is: JHZ <adr>, or Jump if Not Zero to address <adr>.

Now we are ready to put together a program. Just as an example, let's erase the page and then plot a line around the screen. That will require 13 plot numbers, just as in BASIC: $12,2,0,0,127,0,127,127,0,127$, $0,0,255$. Let's put these numbers (the data) at some convenient but out of the way location in memory; say, at 9000 H . Get out your MLDP program and type
and then enter those 13 numbers in this fashion:

MBD=\#12
MEPD=2
MEMD= $\emptyset$
$\cdot$
etc.

NOTE: If you have a 16 K machine, your MLDP program may have a bug in it. Please read last issue's Editors' Notes and make the necessary correction, otherwise your efforts here will be in vain.

SECOND NOTE: Remember that you place numbers into memory using the " =" command. Whatever number you type in will be assumed to be hexadecimal unless preceded by "非". Since the numbers 0-9 decimal equal the numbers 0-9 hex, the "非" won't be necessary in such cases. That's why I didn't bother with it in the example above.

And now for the program to transfer those 13 bytes onto the screen. We can put the program anywhere (but not, of course, where it would interfere with the data). Let's put it at 8200 H .

| Address Contents |  | Assembly Language |  | Comments |
| :---: | :---: | :---: | :---: | :---: |
| 8200 H | 21H | LXI | Н,90Ø0Н | ;Point to start of data. |
| 8201H | 90H |  |  |  |
| 8202H | 90H |  |  |  |
| 8203H | OEH | MVI | C, ODH | ;Initialize counter. |
| 8204H | 0 DH |  |  |  |
| 8205H | 7EH | MOV | A, M | ; Get next byte of data. |
| 8206H | $\square \mathrm{CDH}$ | CALL | 33H | ;'Plot' it. |
| 8207H | 33H |  |  |  |
| 8208 H | 00H |  |  |  |
| 8209 H | 23H | INX | H | ;Point to next byte of data. |
| 820AH | ØDH | DCR | C | ;Done with all bytes? |
| 820 BH | $\square \mathrm{C} 2 \mathrm{H}$ | JNZ | 8205H | ;Not yet. |
| 820 CH | 05H |  |  |  |
| 820 DH | 82H |  |  |  |
| 820EH |  |  |  | ;Yes. All done! |

Notice that two byte numbers are always stored low byte first, then high byte. That's why the contents of locations $8201 \mathrm{H}-8202 \mathrm{H}, 8207 \mathrm{H}-8208 \mathrm{H}$ and $820 \mathrm{CH}-820 \mathrm{DH}$ are the way they are.

There are two ways of entering this program into memory using the MLDP. You can load each memory location with the necessary hex values. For example:

MERD@8200
MEMD=21
MEDD $=0 \emptyset$
MEDD=90
MED $=0$ E
-
$\dot{-}$
Or you can just type in the assembly language mneumonics (but not the comments):

```
MEMD@8200
MEHDLXI H,9000
MEPDMVI C,#13
MEEDMOV A,M
-
etc.
```

That is, the MLDP will act as a mini-assembler and translate the assembly language code into the required hex values for you.

Don't forget to put a BREAKPOINT right after the program:
MEED/
DGB>AT 820E
And now you can execute the program, either full speed with

## DBG $>$ R $82 ø \emptyset$

or else in the interpreted mode with

## DBG>I 82øø

NEXT TIME: Some more 8080 instructions, a look at the Status Flags, and still more ways of putting data onto the screen. In the meantime, have you been reading anything and everything you can get your hands on regarding 8080 programming? That's the way to learn. Our sessions here cannot hope to cover everything; I expect you to be doing some homework. E

# 'THE' BASIC Editor 

Reviewed by DAVID B. SUIIS

The BASIC in ROM in your machine has some rudimentary editing capabilities: if you wish to make some changes to your BASIC program, you needn't write the entire program over again. But a more sophisticated editor would be a great boon to BASIC programmers. That's why I was so excited when ISC released 'FREDI', their BASIC editor, some time ago.

But now a newer, even more powerful editor is available, written by M. A. E. Linden and distributed by Quality Software Associates. It is called 'The' BASIC Editor. It does all that FREDI does and much, much more besides. Like FREDI, 'The' Editor lives in the top of your computer's memory; it takes up only slightly more room than FREDI. Un1ike FREDI, 'The' Editor is a screen editor: instead of editing one line (shown at the top of the screen) as FREDI does, 'The' Editor allows you to move a cursor up and down the screen and then move into a line in order to edit it. You can copy one line or a block of lines into a different part of the program. You can delete a line or a block of lines; scroll the screen up or down; insert one or more lines (with auto line numbering); and search for a specified string.

The good news does not stop there. 'The' Editor will also renumber your program. And it will restore a lost program. (Ever hit ESC $W$ and RETURN and then regret it? Not to worry!) It keeps you informed about the amount of memory you have left. It will LOAD a BASIC file for you. 'The' Editor will also politely disappear in case you don't want it around any longer.

But perhaps the most spectacular feature of this spectacular editor is its ability to append a program on disk to the one you have in memory! What's more, it will put the appended program starting at any line number you wish, resolving all the GOTOs and GOSUBs. This means that you can have a library of routines on disk and use 'The' Editor to append any one or more of them as and where you need them.

The version we received was an early one. Although the upgraded version was not available at press time, it should be available as you read this. The upgrade will have typomatic, it will list the program's variables (including arrays), and it will have a few other esoteric goodies which will bogg1e your mind.
'The' Editor is priced at $\$ 49.95$ (US), and, in my opinion, it is well worth every cent of that. It comes with an extensive manual, and if future changes are made, purchasers receive a free upgrade. If you are more than a very casual BASIC programmer, this is one software tool you will be overjoyed with. 'The' Editor works with V6.78, V8.79 and V9.80 systems, either 16 K or 32 K . (It will even be available very soon in a ROM version.) Contact:

Quality Software Associates 21 Dersingham Crescent Thornhil1, Ontario, Canada, L3T 4P5.

## HONARD ROSEN, INC.

```
    Fut the finishing touches to your Compucolor II or ISC computer.
\becauseCome lip to the world of word processing.
* Extend the utilizetion of sour computer to the other members of your famils.
    * Letters
    * School reports
```



```
            * If you now type-write it, COMFin-W-wite it for a better product.
` Essic requirements for CCII or 3651/9651%
            16% FIAM.
            1174es keyboard.
            F'ririter +
            COMF-U-writer software aris iristruction marigal.
* For maximum capability:
            Full 32t FAM.
            Lower c`ase characters.
    Talk to other compiters: Aisd a MODEM to your system.
    We carry the ertire CCIT & ISC lire of hardware/software, iricludirig spares+
    Serd for our 4-page order form for hardware/software. Request separately
    by item your spare parts reeds. Call us for computer servicing arid mpgraring
        Serid your order now. We pay the shippirig.
        Allow S weeks for delivery.
        We are ari authorized ISC ूealer ....
```



## Back Issues Sale

Back issues of Colorcue are an excellent source of information about Compucolor computers, ISC computers, and programming in general. Interviews, interesting articles, and programs are all there with a touch of history.

The list below includes every Colorcue ever published. If it's not on the list, then there wasn't one.

RETROVIEN: Vol. 3, 非1 (Dec 79/Jan 80) includes: an interview with Bill Greene; Compucolor-teletype interface; user group hotline; introduction to the Screen Editor; PEEKing at BASIC programs; talking to other computers; making programs compatible with $V 6.78$ and $V 7.89$ software; software modifications.

MULTI-ISSUES at $\$ 3.50$ eachOct, Nov, Dec 1978
_ Apr, May/June
1979
__ Jan, Feb, Mar 1979 $\qquad$ Aug, Sept/Oct
1979
IISDIVIDUAL ISSUES at $\$ 1.50$ each


INDIVIDUAL ISSUES at $\$ 2.50$ each
__ Dec 1980/Jan 1981 _ Aug/Sep 1981 __ Oct/Nov 1981

## POSTAGE

US and Canada -- First Class postage included.
Europe, S. America -- add $\$ 1.00$ per item for air, or \$. 40 per item for surface.
Asia, Africa, Middle East -- add $\$ 1.40$ per item for air, or \$. 60 per item for surface.

## DISCOUNT

For orders of 10 or more items, subtract $25 \%$ from total after postage.

ORDER FROM:
Colorcue Editorial Offices 161 Brookside Dr. Rochester, NY 14623


## Colorcue

# Editorial Offices 

161 Brookside Dr.
Rochester, NY 14618

## Address Correction Requested

## Coloreue



Business Software Reviews

Assembly Language
Programming

More


Feb/Mar 1982

## Colorcue

## A Bi-monthly Publication by and for Intecolor and Compucolor Users

## Editors:

Ben Barlow
David B. Suits

## February/March, 1982

Volume 4 , Number 4

## 3 Editors" Notes

3 Compuworld Business Software
5 Frepost Computers, Inc. ROM Board

## 5 Cueties

7 3-D Graphics, by Doug Van Putte Reflection, shear, rotation and scaling

15 Compuwriter Word Processor, by Howard Rosen
Software review
16 Classified Advertising
17 CRT Mode Plotting, by Bob V. Smith Keyboard tables to make graphics easier

## 19 Assembly Language Programming, by David B. Suits Part IV: The Status Flags and the Stack

Advertisers: You will find our advertising policies attractive. Write for details.

Authors: This is a user-oriented and supported publication. Your articles/tips/hints are required to make it go. Send your articles or write for information.

Colorcue is published bi-monthly by Intelligent Systems Corporation, with editorial offices in Rochester, New York. Editorial and subscription correspondence should be addressed to the Editors, Colorcue, 161 Brookside Dr., Rochester, NY 14618. Product related correspondence should be addressed to ISC, 225 Technology Park, Norcross, GA 30092, ATTN: Susan Sheridan. Opinions expressed in by-line articles are not necessarily those of the editors or of ISC. Hardware/software items are checked to the best of our abilities but are NOT guaranteed.

## Editors' Notes

## Dial Up Colorcue

Some of our subscribers have accounts with the Source, Compuserve (Micronet), or other system. We'd like to publish a directory of such persons so that Compucolor/Intecolor users can contact each other to exchange ideas, chew the fat, and otherwise pursue the hobby. If you would like your name included in the directory, send us your name, the name of the dial-up facility, and your ID number. Better yet, if you're already dealing with Micronet, leave the message for us there: Colorcue's editors (yours truly) are now on Micronet at account number 70045,1062.

By the way, it's easy to get onto Micronet. Drop by your local Radio Shack store and purchase the "Dumb Terminal" information package (Cat. No. 26-2224). This gives you an account number and password for Compuserve's Micronet information facility and one hour of free connect time. There is also an application for permanent status which you must fill out and send in. You will also need a simple program to make your machine into a "dumb terminal". Finally, you will need a modem. If you're already dealing with Micronet, the Source, etc., why not tell us all about it via an article for these pages?

## Corrections

It's just not true that there's always an error in Colorcue. Who said that?! Well, there's the occasional oversight.... Dave's Assembly Language Programming had a small omission last issue which is corrected in Part IV this month. And Ben's article on the serial to parallel interface ought to have indicated, on the schematic on page 17 of last issue, that pins 2 and 6 of the 555 should be connected. Please draw that in so you won't forget.

## Doings Down Under

Ralphe Neill of the Victorian (Australia, not B.C.) Users' Group writes a "Coripucolor Column" in Australian Personal Computer magazine. His column is an unexpected bright spot in these non-Compucolor days. It almost seems as though, now that the Compucolor II is out of production, support is starting to emerge. Witness also Tony Watson and John Newman's book Programing Colour Graphics for Compucolor/Intecolor Computers. The book takes you from initial power on through BASIC programming to color (and colour!) graphics, random files and even a driver program for the HIPAD digitizer. Bernie Muldowney, also of the Victorian Users' Group, is selling an assembly language tutorial series on disk. We have seen part of it, and it looks very good. $\$ 50$ (Australian) plus postage. Contact Bernie at 5 Dixon Street, Wangaratta, Victoria, 3677, Australia. And there are two active users' groups with growing program libraries. The Victorian Group can be contacted through its Secretary-Treasurer, Keith Ochiltree, P.O. Box 420, Camberwell, Victoria, 3124. The Western Australian group's librarian is Tony Lee, 52 Cowan Rd., St. Ives, NSW, 2075. ©

## COMPUWORLD BUSINESS SOFTWARE

Compuworld, Inc. has announced a line of business software for the Compucolor II/ISC 3600 series computers. All programs require 32 K user memory, one disk drive and a minimal (i.e. regular) keyboard.

## Colorcalc

(CCII/ISC 3600 series, $\$ 199.00$ )

## ColorGraph

(ISC 3600 series only, $\$ 150.00$ )
The so-called spread sheet is a traditional managerial tool which aids in the analysis of financial information. Data is presented in rows and columns and may represent sales informa-
tion, inventory, etc. As a company's budget changes, the data in the rows and columns are changed accordingly, and the entire array can be re-examined. ColorCalc is a computerized spread sheet which automatically totals the information you are after. The results can be handed over to another program, ColorGraph, which will then provide color graphic displays. Both programs are written in assembly language.

ColorCalc has 12 single-key commands, such as "L" to load data files or "D" for data entry. As many as 51 columns and 52 rows of data may be entered. A formula may be entered on the basis of which the data will then automatically be analyzed. Such a formula may operate on any column or row any number of times. And up to ten formulas may be saved. ( 24 formulas on 3600 machines.) ColorCalc operates with a triple-precision math package which is accurate to the penny up to five billion dollars. Hard copy print out is supported.

ColorGraph takes data from ColorCalc and draws graphs of any section of the spread sheet. You can re-scale the graphs, re-color them, re-label them, and print them -- in black and white on an Epson, for example, or in color using the PrintaColor inkjet printer.

## Inventory Control

(CCII/ISC 3600 series, $\$ 150.00$ )
Among the many features of this package is the ability to keep track of distributor's name, quantity on hand, quantity on order, re-order level, last re-order date, late received date, wholesale/retail pricing, and current order price. Reports may be generated by different parameters, including re-ordering, on order, and total retail and wholesale stock value at any time.

## Accounts Receivable

(CCII/ISC 3600 series, $\$ 150.00$ )
Designed with the average operator in mind, this package can keep track of credit ratings and credit limits. Also included are aged account status, his-
tory of total business done with any account, total monies owed at any time by all accounts, statements as often as required, and mailing labels by two parameters. On the 3600 version, this program has linkage to the General Ledger.

## Accounts Payable

(CCII/ISC 3600 series, $\$ 150.00$ )
This package supports invoice aging, with account status reporting, total monies owed, and a history of total business done with each vendor. Partial payments are supported, as well as manual and automatic check writing and check registers and mailing labels. On the 3600 series, this program has linkage to the General Ledger.

## General Ledger

(CCII/ISC 3600 series, \$299.00)
This package is menu driven with 10 levels of account totals, account balances for current month, quarter, previous three quarters, year to date and previous year. It generates trial balance and income statements, balance sheets, special reports and account status reports. A triple precision math package is incorporated and hard copy print out is supported. (This also creates data files for ColorCalc. And for 3600 series machines, there is linkage to Accounts Receivable, Accounts Payable, and, in the near future, Inventory Control.)

## Mailing List

(CCII/ISC 3600 series, $\$ 100.00$ )
The Mailing List system stores data for personal and business mailing: names, addresses, phone numbers and comments. The program provides for searching on one or more of eight fields, alphabetical storage of names, support of printing one, two or three labels across a page, sorted in alphabetical or zip code order.

Contact Compuworld, Inc., 125 White Spruce Blvd., Rochester, NY, 14623 (716) 424-6260.

## FREPOST COMPUTERS, INC. ROM BOARD

## Add-on ROM Board

Most ISC computers have a slot for additional ROM (Read Only Memory). In the Compucolor II, the area in memory from 4000 hex to 5 FFF hex is reserved as a user ROM area. Therefore, if you have software you wish to be available at any time, it can go into this area.

Frepost Computers, Inc. will supply a plug in ROM board for your ISC computer. The board is ready for you to plug in EPROMs of your own, or pre-programmed EPROMs available from several sources. Since the Compucolor V6.78 system allows only one user jump vector (ESC ^), your EPROMs can be accessed from a jump instruction poked at 33215. As an added enhancement, Frepost Computers has available a replacement system EPROM to replace one of the original V6.78 system chips. This allows four software or keyboard jumps to the PROM board and adds two more user programmable jumps (like having three ESC ${ }^{\wedge}$ ). And as an added feature, you can customize your system to your requirements. Frepost Computers will even burn your program into EPROM for you.

## Bank Selectable EPROM Controller

We also have a device that will plug into the Compucolor CPU board in the Add-on ROM connector just like the small board mentioned above. This device accepts commands from a program or from the keyboard and selects up to seven
more banks of ROM, all occupying the same memory addresses, 4000-5FFF. The EPROM bank selector module adds 56 K of additional ROM that is software selectable in 8 K byte segments. It uses 2532 type chips which are the same as your system ROMs. If you already have an add on ROM board in your machine, it can be pulled out of its sockets on the logic or CPU board. The EPROM bank select board plugs into these sockets and your present ROM board plugs into the bank selector board. There is a 50 pin bus connector that plugs into the back of the logic board. If you're already using the 50 pin bus connector for something else, we supply you with all instructions on how to solder the appropriate leads to your computer, eliminating the need for our 50 pin bus connector. (If you order it that way, you save \$10.)

The Add-on ROM board is available for $\$ 49.95$, assembled and tested, or in kit form for $\$ 39.95$, complete with sockets, capacitors and connectors. The EPROM bank selector board is $\$ 249.95$ assembled and tested. In kit form, with all sockets, capacitors and chip select parts, it is \$199.95. (Either version with 50 pin bus connector is $\$ 10$ more.) The Frepost V6.78 enhanced system PROM is available for $\$ 29.95$. Orders less than $\$ 100$, please include $\$ 5.00$ shipping and insurance. All orders include complete instructions and documentation. Write or call Frepost Computers, Inc., 431 East 20th Street 10D, New York, NY 10010. Phone: (212) 673-6476. The Source: TCI251. Micronet: 70210,374.

## Cueties

PLOT 12: $\mathrm{X}=64: \mathrm{Y}=\mathrm{X}: F O R \quad \mathrm{Z}=\emptyset \mathrm{TO} 50 \emptyset 0: \mathrm{X}=\mathrm{X}+2-\operatorname{INT}(\operatorname{RND}(3) * 5): Y=Y+2-$ INT (RND (2)*5): PLOT $2, X, Y, 255: N E X T$

HOWARD FOGEN, IMC.
Aththorixed ISC dealer

TALKEDTOYOUR COMFUTER LOTELY? DID IT TAKK EACK? NOW TTMCAN!!!
with the Amazing
TYFE… N-TALK

- Type \& hear the written word spoken - Simply commect to the FS232C port

To order semd check for $\$ 362,00$ whin your mame, adsress, city, state, \& zip. (FFa, resident please add $6 \%$ sales tam). We pay the shipping charges.

- We carry the word processor for the cctx, 3621, 3650. Your computer isn't the seme, once you've added the word processar.
*..... letters .... reports .... forms are easjup prepared with the aid of the word processor.
- Servicing ayalable for sour Com
- New hardware 8 software 1 ist with itemized bescription
- Send for new list and order form *
- Never a shippiris charge for software 8 hardware purchases.

HOWAFD FOSEN. TInc.
F. O. EOK 434

Humtingeron Valley, Fa, 19006
(215) 464-7145

* Xrquire about our free programs


# 3D Graphics 

by<br>Doug Van Putte<br>18 Cross Bow Dr.<br>Rochester, NY 14624

Graphic displays on the computer are enhanced by giving depth and motion to objects. The world of animinated graphics using matrix mathematics to present objects with simulated depth and motion will be introduced. I will cover the basic concepts and then describe the elementary motions that can be attained. These motions will be illustrated by a program that can be applied to any object by the reader. Those of you who find the technical treatment too superficial are referred to the references.

First, you will be given a review of the Cartesian coordinate system, followed by the object matrix which is used to define the coordinates of a point on the object. This will be followed by an introduction to the transform matrix which contains the elements of motion that can be applied to the object (matrix). The multiplication of the object matrix and the transformation matrix will then be shown to yield a new, transformed coordinate matrix from which the object can be drawn in its new position. I hope you can visualize now that a previously defined object can be moved through a series of motions by successively applying the operations above to each point on an object, followed by redrawing the object from its "new" points. The types of motion which include scaling, reflection, translation, rotation, and shearing will be shown in simple matrix notation.

CARTESIAR COORDINATE SYSTEM. The Cartesian coordinate system used exclusively here is the PLOT 2 coordinates familiar to ISC computer users. While the origin is usually in the lower left hand corner of the screen, the demo program uses an origin at the center. The X-axis is the horizontal axis, the $Y$-axis is the vertical axis, and the $Z$-axis, by convention, is perpendicular to the $X$ and $Y$-axes and runs toward the viewer and into the depth of the screen. Positive values are toward the right for $X$, up for $Y$, and out of the screen for $Z$.

OBJECT MATRIX [P]. The object in all cases will be a series of coordinate points which are connected by lines (PLOT 242). The individual $X, Y, Z$ coordinates for a point are stored in a column matrix $P$, which looks like this:

| (P is a matrix with <br> one column and <br> three rows) | $\|$$X$ <br> $Y$ |
| :--- | :--- |
| $Z$ |  |$|$| Points are stored in a |
| :--- |
| 2 dimensional array, i.e. |
| $X 1=P(1,1): Y 2=P(2,1): Z 1=P(3,1)$. |

If ten points are required to define the object, then ten column matrices, each with one column and three rows, must be created. E.G., $P(1,1), P(2,1)$, $P(3,1), \ldots P(1,10), P(2,10), P(3,10)$. The object is constructed by connecting the coordinate points $P$ by using PLOT 242 to draw lines between each successive point.

TRAMSFORTATION MATRIX [T]. This matrix is sometimes called the operator matrix because each non-zero element is a term which operates on or modifies one of the coordinates, point by point, to transform the object. For now, T will be defined as a $3 * 3$ matrix as follows:

| (I is a matrix with | A B C | Row 1 modifies X. |
| :---: | :---: | :---: |
| three columns and | D E $\bar{m}$ | Row 2 modifies Y. |
| three rows.) | G H I | Row 3 modifies Z . |

The transform matrix $T$ can be multiplied by the object matrix $P$ co yield a transformed object matrix $P^{\prime}$ as follows:

$$
P^{\prime}=T * P
$$

Upon expansion, the new coordinates of $P^{\prime}$ become:

$$
\begin{aligned}
& X^{\prime}=X * A+Y * B+Z * C, \\
& Y^{\prime}=X * D+Y * E+Z * F, \text { and } \\
& Z^{\prime}=X * G+Y * H+Z * I_{0}
\end{aligned}
$$

Rules of matrix order and shape have to be observed in the multiplication process. (See reference 3.) Using the ( $3 * 3$ ) transformation matrix $T$, operations of scaling, reflection, rotation, and shearing can be used to manipulate the coordinates of a point $P$. Other operations, called translation and overall scaling, will make use of an expanded (4*4) transformation matrix. All the operations will be illustrated with a simple cube in a BASIC program later on.

SCALING. For pure scaling, T becomes

$$
T=\left|\begin{array}{lll}
A & 0 & 0 \\
0 & E & 0 \\
0 & 0 & I
\end{array}\right|
$$

And when $T$ is multiplied by $P$, then

$$
\begin{aligned}
& X^{\prime}=A * X, \\
& Y^{\prime}=E * Y, \text { and } \\
& Z^{\prime}=I * Z .
\end{aligned}
$$

Independently then $X, Y$ \& $Z$ coordinates can be scaled to produce a reduction or a magnification of an object. The operation of reflection is similar.

REFLECTION. To obtain the reflection of an object in a given direction, all the signs of that coordinate direction must be reversed. For example, to obtain the mirror image of an object through the plane of the screen ( $\mathrm{X}-\mathrm{Y}$ ), T becomes:

$$
T=\left|\begin{array}{rrr}
1 & 0 & 0 \\
0 & 1 & 0 \\
0 & 0 & -1
\end{array}\right|
$$

And when $T * P$ then

$$
\begin{aligned}
& X^{\prime}=X, \\
& Y^{\prime}=Y, \text { and } \\
& Z^{\prime}=-Z .
\end{aligned}
$$

SHEAR. An object is sheared by modifying one or more coordinate values of a point in a special manner. A coordinate's value is modified from the original value by an amount computed from the point's position on one or both of the other two coordinate axes. An object can be sheared in its X-direction by:

$$
T=\left|\begin{array}{lll}
1 & B & 0 \\
0 & 1 & 0 \\
0 & 0 & 1
\end{array}\right|
$$

And when $T * P$, then

$$
\begin{aligned}
& \mathrm{X}^{\prime}=\mathrm{X}+\mathrm{B} * \mathrm{Y}, \\
& \mathrm{Y}^{\prime}=\mathrm{Y}, \text { and } \\
& \mathrm{Z}^{\prime}=\mathrm{Z}
\end{aligned}
$$

On evaluation, this operation tilts an object by displacing the horizontal ( $X$ ) values an amount $B * Y$. Then as a point's vertical ( $Y$ ) position increases, the displacement of $X$ increases and therefore the redrawn object will be tilted. Note that the $X$ value can be modified by both the $Y \& Z$ position of the point at the same time by making the ' $C$ ' value non-zero. All three coordinates of a point can be sheared simultaneously by the other coordinates by making B, C, D, F, G, \& H non-zero. This operation yields interesting object motions, but the operation of rotation can produce the most life-1ike movement.

ROTATION. An object can be rotated about any axis to produce a specific orientation in space. The transform examples, however, are restricted to the coordinate major axes. By specifying the rotation angle, an object can be rotated in one of three planes about its origin by using one of the following transformations:

X-Y PLANE (about the Z-axis):

$$
T=\left|\begin{array}{ccc}
\cos (a) & -\operatorname{Sin}(a) & 0 \\
\sin (a) & \cos (a) & 0 \\
0 & 0 & 1
\end{array}\right|
$$

And when $T * P$, then

$$
\begin{aligned}
& X^{\prime}=X * \operatorname{COS}(a)+Y * \operatorname{SIN}(a), \\
& Y^{\prime}=X * \operatorname{SIN}(a)+Y * \operatorname{Cos}(a), \text { and } \\
& Z^{\prime}=Z .
\end{aligned}
$$

X-Z PLANE (about the Y-axis):

$$
T=\left|\begin{array}{ccc}
\cos (a) & 0 & -\operatorname{SIN}(a) \\
0 & 1 & 0 \\
\operatorname{SIN}(a) & 0 & \cos (a)
\end{array}\right|
$$

And when $T * P$, then

$$
\begin{aligned}
& X^{\prime}=X * \cos (a)+Z *-\operatorname{SIN}(a), \\
& Y^{\prime}=Y, \text { and } \\
& Z^{\prime}=X * \sin (a)+Z * \cos (a) .
\end{aligned}
$$

Y-Z PLANE (about the X-axis):

$$
T=\left|\begin{array}{ccc}
1 & 0 & 0 \\
0 & \cos (a) & -\operatorname{SIN}(a) \\
0 & \operatorname{SIN}(a) & \cos (a)
\end{array}\right|
$$

And when $T * P$, then

$$
\begin{aligned}
& X^{\prime}=X \\
& Y^{\prime}=Y * \operatorname{COS}(a)+Z *-\operatorname{SIN}(a), \text { and } \\
& Z^{\prime}=Y * \operatorname{SIN}(a)+Z * \operatorname{COS}(a)
\end{aligned}
$$

Although not within the scope of this treatment of rotation, it can be stated that:

1) The final positon of an object rotated sequentially about different axes is a function of the order of applying the transformations.
2) An object may be rotated about any arbitrary point by translating that point to the origin prior to rotation and subsequently returning the point to its original position.
3) An object can be rotated about any arbitrary axis placed through the object.

Now for the theory and transformations for operations of Translation and Overall Scaling. This treatment will be followed by a demonstration program.

In order to present the last two operations, both the $P$ (object) matrix and the $T$ (Transform) matrix must be expanded. They become:

$$
P(X, Y, Z, I)=\left|\begin{array}{c}
X \\
Y \\
Z \\
I
\end{array}\right| \quad T=\left|\begin{array}{cccc}
A & B & C & U \\
D & E & F & V \\
G & H & I & W \\
0 & 0 & 0 & \text { S }
\end{array}\right|
$$

On multiplication, $T^{*} * P$ yields $P^{\prime}\left(X^{\prime \prime}, Y^{\prime \prime}, Z^{\prime \prime}, S\right)$. On expansion, the new coordinates become:

$$
\begin{aligned}
& X^{\prime \prime}=X * A+Y * B+Z * C+U, \\
& Y^{\prime \prime}=X * D+Y \div E+Z * F+V, \text { and } \\
& Z^{\prime \prime}=X * G+Y * H+Z * I+W,
\end{aligned}
$$

The transformed coordinates are now 'normalized' by dividing them by $S$ to yield $P^{\prime}\left(X^{\prime}, Y^{\prime}, Z^{\prime}, 1\right)$ as follows:

$$
\begin{aligned}
& X^{\prime}=X^{\prime \prime} / S, \\
& Y^{\prime}=Y^{\prime \prime} / S, \text { and } \\
& Z^{\prime}=Z^{\prime \prime} / S \text {. }
\end{aligned}
$$

In addition to all the operators described earlier, these new forms of the $P$ and $T$ matrices can be used to move and scale an object. For example:

TRAMSLATION. For translation only, T becomes

$$
T=\left|\begin{array}{llll}
1 & 0 & 0 & \mathrm{U} \\
0 & 1 & 0 & \mathrm{~V} \\
0 & 0 & 1 & \mathrm{~W} \\
0 & 0 & 0 & 1
\end{array}\right|
$$

And when T*P, then

$$
\begin{aligned}
& X^{\prime}=X+U, \\
& Y^{\prime}=Y+V, \text { and } \\
& Z^{\prime}=Z+W \text {. }
\end{aligned}
$$

Thus it is seen that the first three elements ( $U, V, W$ ) of the four th column are the elements of translation for each of the three coordinates ( $X, Y, Z$ ), respectively. This transformation, in conjunction with the others discussed above, provide us with the means of giving realistic movement to an object.

OVERALL SGALING. For Overall Scaling only, T becomes

$$
\mathrm{T}=\left|\begin{array}{llll}
1 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 \\
0 & 0 & 1 & 0 \\
0 & 0 & 0 & \mathrm{~S}
\end{array}\right|
$$

And when $T * P$ and the coordinates are normalized by dividing by $S$, then

$$
\begin{aligned}
& X^{\prime}=X / S, \\
& Y^{\prime}=Y / S, \text { and } \\
& Z^{\prime}=Z / S .
\end{aligned}
$$

Thus, the fourth element of the last column can be used to scale the three coordinates identically.

## DEHONSTRATION PROGRAM.

A few comments about the demo program are in order.

1) All prompted inputs are numeric, but the values need to be tailored for the type of operation. Values for each input which cause no operation are given with each prompt. The rotation angles are interpreted as degrees. Start at values close to the no-op values given to begin experimentation.
2) There is no internal check for exceeding the screen limits.
3) The program produces a composite transform matrix before the object is operated upon. To test the effect of each operation on the object, the user can change the program, or simply use the no-op values for the other operations.
4) The user can input his own object. Follow the program comments and don't forget to change the 'NR' value for the number of object points and the matrix dimensions.
5) The program is unfortunately slow. What is needed to make the program swifter is an 8080 subprogram which can be called to multiply matrices. Does anyone have one? If not, why don't one of you 8080 Wizards write one and share it.
6) All the lines, including the hidden lines, are drawn by the Plot subroutine. Can anyone add a hidden line subroutine?
7) Undoubtedly, there are the inevitable bugs. Please pass any you find along to me. E


## REFERENCES:

1. J. Posdamer, "The Mathematics of Computer Graphics", BYTE, Sept., 1978, pp. 22-39.
2. J. Hungerford, "Graphic Manipulations Using Matrices", BYTE, Sept., 1978, pp. 156-165.
3. D. Rogers and J. Adams, Mathematical Elements for Computer Graphics (McGraw Hil1, 1976).
```
100 REM ************* 3-D GRAPHICS DPMO PROGRAM *********************
110 REM ************* BY D. A. VAN FUITIE *****************
120 FLOT 12,15:DDM G(4,30),P(4,30),S1 (4,4),S2(4,4),T(4,4):NR=24
130 Il(1)=1:I2(1)=2:Il(2)=1:I2(2)=3:II (3)=2:I2(3)=3:DX=45:DY=45:DZ=45
140 REM * READ OBJECT POINTS FROM ARRAY (LINE 9000)
150 RESTORE :FOR C=1 TO NR:P(4,C)=1:FOR R=1 TO 3:READ P(R,C):NEXT :NEXT
160 GOSUB 530:GOSUB 600:GOSUB 500
170 REM * SCALE OPERATION-REOUIRES INPUT OF 3 NOS.
180 INPUT "SCALE X,Y,Z (1, 1,1)? ";A,E,I:Sl (1,1)=A:Sl(2,2)=E:Sl ( 3,3)=I:IF A=1
    AND E=1 AND I=1 THEN 210
190 GOSUB 450
200 REM * OVERALL SCALE OPERATION-REQUIRES INFUT OF 1 NO.
210 INPUT "OVERALL SCALE EACIOR S (1)? ";S:IF S=1 OR S=0 THEN 240
220 Sl (4,4)=S:GOSUB 450
23ø REM * SHEAR OFERATION-REQUIRES INPUT OF 6 NOS.
240 INPUT 'SHEAR X(BEC),Y(D&F),Z(G&H) (G,0,0,0,0,0)?";B,C,D,F,G,H:IF B=0 AND
                C=0 AND D=0 AND F=0 AND G=0 AND F=0 THEN 280
250 Sl(1,2)=B:Sl (1,3)=C:Sl (2,1) =D:Sl (2,3)=E:Sl (3,1)=G:Sl (3,2)=日
260 GOSUB 450
270 REM * ROTATION OPERATICQ-REQUIRES INRUT OF }3\mathrm{ NOS.
286 INPUT 'ROTATION ANGLE IN X-Y;X-Z,Y-E PLANES (0,0,0)? m;A(1),A(2),A(3):IF
    A(1)=0f AND A(2)=0 AND A (3) =0 THEN 310
290 GCSUB 550:GOSUB 450
```




```
        35@
320 SI (1,1)=A:S1 (2,2)=E:S1 (3,3)=1
330 cas!m 450
```




```
        AND W=01 THEN 38%
```



```
370 FOR C=1 TO RR:P(1,C)=P(L,C)*DD:P (2,C)=E (2,C)*DD:NEXI
380 GNSTB 418:PLOM 12:GOSTB 501/
390 INEUT MHIT RETURN TO RESTART :2Z:FLOT 12:GOTO 150
ADG RE% * SUB TO MULTIFLY T MATPIX & P MATRIX
410 ROR I=1 TO 4:FOR J=1 TO NR:SM=0:FO{ R=1 TO 4
420 SM=SM+T(I,R)HP(R,U):NEXT R:G(I,J)=SN:NEXX J:NEXT I
430 FOR R=1 TO 4:FOR C=1 TO NR:P(R,C)=TNT(G(R,C)+.5)/G:NEMT :NEXM :RETUR
    N
```



```
450 ROR I=1 TO 4:ECR Jm1 TO 4:SN={:ECR R=1 TO A
```



```
470 S2(I_J)=SM&MENT J:NEWT I:FOR RED TO &
4EO FOR C=1 TU 4:T(R,C)=S2(R,C):NENT C:NRET R
490 REM * SUS TO PIOT TRANGFONUCD OBJECT (P')
500 FLOT 2,F(1,1)+DX,P(2, 1)+NX,242:FOR C=2 TO 18:ROT P(1,C)+DX,P(2,C)+DY
:NEXT :PLOT }25
519 FLOT 2,P(1,19)+DX,P(2,19)+DY,242:FOR C=20 TO 24:PMOT P(1,C)+DX,P(2,C)
+DY:NEXT :FLOT 255:REIUTNT
52g RE* * SUB TO INITIALIZE T MATRIX
530 FOR R=1 TO 4:FOR C=1 TO 4:T(R,C)={:NEXT C:NEXXT R:T(1,1)=1:T(2,2)=1:T
(3,3)=1:T (4,4)=1:RETURN
540 REM * SUB TD CREATE ROTATICN TMNSEORN MRTRICES
550 FOR L=1 TO 3:A=A(L)*3.14I59/180:IF A(L)=0 THEN 580
560 Sl (II (L),Il(L))=COS (A):S1(II (L),I2 (IN) )=-SIN(A)
570 Sl(I2(L),I1 (L))=SIN(A):S1(I2(L),L2(L))=COS(A)
580 NENT L:REIURN
599 IENG STB TO INITIALIKE GENERAL TRARSFCRTH MATRTX (S1)
```



```
        1:SI (3,3):1:Sl (4,4)=1:REIURN
```



```
620 REM * EXAMPLE DATA IS A 3-D CUBE W/ LETWMCY "g8 CN FRONT EACE
```



```
640 DATA 0.0,30,0,30,30,30,30,30,0,30,30,0,30,0,0,0,0
650 DATA D,0,30,30,0,30,30,30,30,30,0,30,30,0,0
660 DATA 14,13,30,14,17,30,17,17,30,14,17,30,14,15,30,16,15,30
```


# Compuwriter Word Processor 

by<br>Howard Rosen<br>P.O. Box 434<br>Huntington Valley, PA<br>19006

Ten ways to use a word processor: write a letter; build a form; write a story; write a report; write a book; write a note to 15 people; write an article; write a research paper; write mailing labels; prepare a sheet of vertical and horizontal lines. The list is not complete, and I'm open to suggestions.

Do any of the above apply to you? Read on and learn the workings of the Comp-u-Writer word processor. How do you use it? Do you have to know how to program? No, you do not do any programming. Let's examine how the word processor is actually used on a day-to-day basis.

While a typewriter can be used to prepare a neat looking document, the word processor does it better and easier. Take, for example, centering the first line of a letter. Just press the key marked CHTER and type. The text magically starts at the very center of ?e lime and moves left, then right, while maintaining a perfect center. A typist has to consider page width, line ending, tab setting, and indentation. All of these are child's play to the word processor. Just press the PRIMT key and a menu with color squares appears. Each color square is associated with a setting: $\mathbb{E}$ 拃 for lines per page; YELLOW for characters per line; RED for margin settings; and more. After the page is typed in, if you wish to change the number of characters per line, return to the menu and change to a different setting. The text will immediately shift to acconodate your new choice.

Tabs are set by moving a little blue square right or left along a scale at the top of the screen and locking them in place by depressing the RETURN key. The TAB key is then exactly like your typewriter, except you can change tabs if you like right in the middle of a letter to get a different effect. Suppose you've finished typing and you've found mistakes. Return to the scene of the error, delete the word or character and make your correction. Perhaps you should have made the first paragraph the third one. Use the two keys labeled MARK BEG and MARK END to place marks on either end of the section in question-watch it turn from green to red to indicate that it's marked--and then move the cursor with the four arrow keys to the exact spot where you want the paragraph to be. Then simply press the MOVE BLOCK key and watch how fast all is forgiven. If the marked block is to be repeated, press the COPY BLOCK key for each copy desired. The ERASE BLOCK key will totally remove the marked block from your text.

To avoid typing long words used repeatedly, you might use $x$ or $y$ or $z$ in substitution, then with the DEFINE, SEARCH, and REPLace keys you can specifically or generally make your replacements. Anything can be underlined or made to be boldface by the press of a key. Any character on the keyboard will automatically repeat if held just a little longer.

Now you're ready to print. You can choose single sided, where each page is sequentially printed, or double sided, where the odd pages are printed, then
you reverse the paper for the even pages. If you wish to do two column printing then each consecutive pair of pages becomes the left and the right columns, respectively. This can be done for double sided printing as well, with the word processor keeping an eye on how things are going. If only a part of what was typed is to be printed, then use the marking technique discussed earlier and print only the marked part. Multiple copies can be printed by responding to the question "HOW MANY COPIES?" on the print menu. Documents can be saved on disk for future use. Either the whole document or just the marked portion can be saved.

The most recent version of Comp-uWriter is version 3.5. There are two types available: the Executive and the

Mail Merge. The Mail Merge, while capable of printing personalized letters, does not support the two sided printing. Mail Merge is useful when a specific letter is needed for a number of customers, each individually addressed.

Comp-u-Writer for the Compucolor II/ 3600 series computers requires the 117 key keyboard and a minimum of 16 K RAM. It is recommended that both 32 K RAM (to support long documents) and the lower case letters (to allow normal typing) be available.

As an authorized ISC dealer, we carry the Comp-u-Writer and the entire ISC hardware and software line. We may be reached at HOWARD ROSEN, INC., P.O. Box 434, Huntington Valley, PA 19006 or by phone at (215) 464-7145. ©

## CALL FOR ARTICLES

Colorcue gets its material from those who write it. It don't grow on no trees. Nevermind your bad sppeling or badly grammar: that's the job of the Editors. You come up with the ideas, splash them onto paper (or disk), and we'll wrestle them into an intelligible form for publication. We can't pay you for your time--you won't become rich. But, since Colorcue is read all over the world (well, not all over, exactly), then maybe you'll become world famous.... We're looking for ideas, programs, reviews of books and/or software and/or hardware applicable to ISC machines, hardware/software modifications, user group news, and.... Well, you get the idea.

## CLASSIFIED ADVERTISING

WANTED: Compucolor II system. Purchase or trade for TRS-80 Model 1. Mike Charlton, (502) 926-3021.

FOR SALE: Compucolor II, V6.78, 24 K , non-working. Analog board problems. Complete with manuals, programs, etc. Ben Moser, Rt. 2, Box 550, Stanley, VA, 22851. (703) 778-2861.

FOR SALE: Heathkit microprocessor learning system. M6800 based trainer, with documentation. Excellent package for introduction to hardware. Micheal Ezzo, 692 S. Drake Apt. 0-10, Kalamazoo, MI 49009 。

FOR SALE: Compucolor II, extended keyboard, 16 K , with diskettes, manual, graphics. Can add modem, sound, and memory. (716) 872-2322 (after 5).

FOR SALE: Intecolor 3651.32 K machine with $51 / 4$ inch disk, std keyboard. Manufactured in mid 1981. Asking \$2000. Jim Dantin, (502) 927-6921 x377 (days), or (502) 926-8333 (nights).

FOR SALE: Intecolor 3651. New in original carton. $32 \mathrm{~K}, 51 / 4$ inch disk. Asking \$2000. George Wilson, (404) 458-1431 (nights).

# CRT Mode Plotting 

by<br>Bob V. Smith<br>498 Brown Street Napa, CA 94559

Using the deluxe keyboard, graphic displays can be drawn in the CRT mode. Control B enters the plot mode (just like BASIC's PLOT 2.) From then on, the ASCII values of each key press will determine what is plotted. (See the keyboard table on the next page.) The special function keys F0-F15 are used to enter the various plot submodes. In order to simplify incremental plotting, the table below can be of help.

For example, draw a triangle this way: enter the CRT mode, Control B for the plot mode, plot a point with two appropriate key strokes, enter the vector plot submode with the F2 key, plot the other end of the vector with two
more appropriate keys, and finally enter the incremental vector plot submode with the F0 key:

ESC CRT
Control B
Control CRT
Shift ?
F2
CRT
Shift ?
F0
Now form a triangle by holding the left end of the vector stationary and moving the right end up (North). According to the table, that is the 2 key. ©

SHF = Shift
CTL $=$ Control
COM = Command (shift \& control)

Left side (A)

|  | N | NE | E | SE | S | SW | W | NW | HOLD |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| N | $\begin{gathered} \mathrm{SHF} \\ 2 \end{gathered}$ | $\begin{gathered} \mathrm{COM} \\ 2 \end{gathered}$ | $\begin{array}{\|c} \mathrm{COM} \\ \mathrm{~B} \end{array}$ | $\begin{gathered} \mathrm{COM} \\ \mathrm{R} \end{gathered}$ | $\begin{gathered} \mathrm{CTL} \\ \mathrm{R} \end{gathered}$ | R | B | $\begin{gathered} \mathrm{SHF} \\ \mathrm{~B} \end{gathered}$ | 2 |
| NE | $\underset{*}{\text { SHF }}$ | $\mathrm{COM}_{*}$ | $\begin{gathered} \mathrm{COM} \\ \mathrm{~J} \end{gathered}$ | $\begin{array}{\|c} \mathrm{COM} \\ \mathrm{Z} \\ \hline \end{array}$ | $\begin{array}{\|c} \mathrm{CTL} \\ \mathrm{Z} \\ \hline \end{array}$ | Z | J | $\underset{\mathrm{J}}{\mathrm{SHF}}$ | : |
| E | $\begin{gathered} \mathrm{SHF} \\ 8 \end{gathered}$ | $\begin{gathered} \mathrm{COM} \\ 8 \end{gathered}$ | $\begin{gathered} \mathrm{COM} \\ \mathrm{H} \end{gathered}$ | $\underset{\mathrm{X}}{\mathrm{COM}}$ | $\underset{\mathrm{X}}{\mathrm{CTL}}$ | X | H | $\underset{\mathrm{H}}{\mathrm{SHF}}$ | 8 |
| SE | $\begin{gathered} \mathrm{SHF} \\ 9 \end{gathered}$ | $\begin{gathered} \mathrm{COM} \\ 9 \end{gathered}$ | $\begin{array}{\|c} \mathrm{COM} \\ \mathrm{I} \end{array}$ | $\underset{\mathrm{Y}}{\mathrm{COM}}$ | $\underset{\mathrm{Y}}{\mathrm{CTL}}$ | Y | I | $\underset{\mathrm{I}}{\mathrm{SHF}}$ | 9 |
| S | $\begin{gathered} \mathrm{SHF} \\ 1 \end{gathered}$ | $\begin{gathered} \mathrm{COM} \\ 1 \end{gathered}$ | $\begin{gathered} \mathrm{COM} \\ \mathrm{~A} \\ \hline \end{gathered}$ | $\begin{gathered} \mathrm{COM} \\ \mathrm{Q} \end{gathered}$ | $\begin{array}{\|c} \mathrm{CTL} \\ \mathrm{Q} \end{array}$ | Q | A | $\underset{\mathrm{A}}{\mathrm{SHF}}$ | 1 |
| SW | $\begin{gathered} \mathrm{SHF} \\ 5 \end{gathered}$ | $\begin{gathered} \mathrm{COM} \\ 5 \end{gathered}$ | $\begin{gathered} \mathrm{COM} \\ \mathrm{E} \end{gathered}$ | $\begin{gathered} \mathrm{COM} \\ \mathrm{U} \end{gathered}$ | $\begin{array}{\|c} \mathrm{CTL} \\ \mathrm{U} \end{array}$ | U | E | $\underset{\mathrm{E}}{\mathrm{SHF}}$ | 5 |
| W | $\underset{4}{\mathrm{SHF}}$ | $\begin{gathered} \mathrm{COM} \\ 4 \end{gathered}$ | $\begin{gathered} \mathrm{COM} \\ \mathrm{D} \end{gathered}$ | $\underset{\mathrm{T}}{\mathrm{COM}}$ | $\underset{\mathrm{T}}{\mathrm{CTL}}$ | T | D | $\underset{\mathrm{D}}{\mathrm{SHF}}$ | 4 |
| NW | $\begin{gathered} \mathrm{SHF} \\ 6 \end{gathered}$ | $\begin{array}{\|c} \hline \mathrm{COM} \\ 6 \end{array}$ | $\underset{\mathrm{F}}{\mathrm{COM}}$ | $\mathrm{COM}$ | $\underset{\mathrm{V}}{\mathrm{CTL}}$ | V | F | $\underset{\mathrm{F}}{\mathrm{SHF}}$ | 6 |
| HOLD | 1 | CTL | $\begin{array}{\|c\|c} \mathrm{COM} \\ \mathrm{O} \end{array}$ | $\begin{aligned} & \text { COM } \\ & \text { CRT } \end{aligned}$ | $\begin{aligned} & \mathrm{CTL} \\ & \mathrm{CRT} \end{aligned}$ | CRT | 0 | $\begin{gathered} \text { SHF } \\ 0 \end{gathered}$ |  |

Keyboard Table．Each row lists ASCII values produced by the indicated keys（plus Control，Command or Shift，if necessary）．

|  | CONTROL |
| :---: | :---: |
|  |  |
|  <br>  |  |
|  | SHIFT <br> NOTHING |
| $\cdots \cdots \cdots$ ．． $0^{\circ}$ |  |
|  <br>  | NOTH ING |
|  |  |
|  <br>  |  |
|  | SHIFT |
|  <br>  |  |
|  | COMMAND |
|  <br>  |  |
|  | COMMAND CONTROL |
|  |  |
|  <br>  |  |
|  | CONTROL |
|  जヵいいい | SHIFT |
| NNNNNNNNNNNNNNNNNNNNNNNNNNNN <br>  <br>  |  |
|  | COMMAND |
|  <br>  | NOTH ING |

## BUSINESS SOFTWARE FOR SALE

－ 1982 INCOME TAX PREPARER（for Federal and New York State）．
－COMPLETE CASH REGISTER PROGRAM（Includes invoice print，order，inventory count， cash／credit total）．
－PAYROLL SYSTEM．
－QSORT PROGRAM
－CCII INTERACTIVE OPERATING SYSTEM

## MAU Corporation

5 Eldridge Street，Store North New York，New York 10002
Tel．（212）431－1277

# Assembly Language Programming 

by<br>David B．Suits

## PART IV：The Status Flags and the Stack

Well，I goofed last time．Some of you saw the mistake and made the approp－ riate corrections．The little program I gave you last time to draw a line around the screen will unfortunately only plot four points at the corners of the screen．That＇s because I left out the incremental plot mode introduction number，242，which should be inserted just after the $12,2,0,0$ in the string of bytes which the program used．In addition，there will now be 14 bytes in the string instead of 13 ，and so the counter must be initialized to 14 （ $=0 \mathrm{EH}$ ） at address 8203 H ．

## The Status Flags

Besides the normal 8080 registers （Accumulator，B，C，etc．），there is an eight bit group of Status Flags，each of which can be set（or＇on＇，＝1）or reset（or＇off＇，＝0），according to the operations of certain instructions．Al－ though the flag register is eight bits wide，there are only five flags；the other three bits are unused．The five flags are the Sign，Zero，Auxiliary Carry，Parity，and Carry Flags．Their Status Flags：

settings are often used automatically by other instructions．For example，the DCR＜reg＞instruction which we examined last time will affect all the flags except the Carry flag．Probably the most useful of the flags it affects is the Zero flag，which will be set if the result of DCR 〈reg＞is zero，and reset otherwise．Then a conditional instruc－ tion such as JNZ 〈addr＞（Jump if Not

Zero）will automatically test the status of the Zero flag．Thus，the sequence

$$
\begin{array}{ll}
\text { DCR } & \text { <reg> } \\
\text { JNZ } & \text { <address> }
\end{array}
$$

is a very common one in 8080 programming． There are other conditional jump instruc－ tions，each of which tests the status of some flag：

| Hex | Instruction |  |
| :--- | :--- | :--- |
| DA | JC | ＜address＞ |
| FA | JM | ＜address＞ |
| D2 | JNC | ＜address＞ |
| C2 | JNZ | ＜address〉 |
| F2 | JP | ＜address＞ |
| EA | JPE | ＜address＞ |
| E2 | JPO | ＜address＞ |
| CA | JZ | ＜address＞ |

## Meaning

Jump if carry（carry flag＝1）
Jump if minus（sign flag＝1）
Jump if no carry（carry flag＝－0）
Jump if not zero（zero flag＝0）
Jump if plus（sign flag＝0）
Jump if parity even（parity flag＝1）
Jump if parity odd（parity flag＝0）
Jump if zero（zero flag＝l）

In the previous installment we made use of the CALL instruction．The CALL is really a jump－to－subroutine instruc－ tion．You might expect，then，that there would be a set of conditional jump－to－ subroutine instructions．And you would be right：

| Hex | Instruction |  |
| :---: | :---: | :---: |
| DC | CC | ＜address＞ |
| FC | CM | 〈address＞ |
| D4 | CNC | ＜address＞ |
| C4 | CN2 | 〈address＞ |
| F4 | CP | ＜address＞ |
| EC | CPE | ＜address＞ |
| E4 | CPO | ＜address＞ |
| CC | Cz | ＜address＞ |

## Mearing

Call if carry（carry flag＝1）
Call if minus（sign flag＝1）
Call if no carry（carry flag＝0）
Call if not zero（zero flag＝0）
Call if plus（sign flag＝0）
Call if parity even（parity flag＝1）
Call if parity odd（parity flag＝0）
call if zero（zero flag＝1）
Finally，there are corresponding RET （return from subroutine）instructions which are executed only if a flag is of a certain value：

| Hex | Instruction | Meaning |
| :--- | :--- | :--- |
| D8 | RC | Return if carry（carry flag＝1） |
| F8 | PM | Return if minus（sign flag＝1） |
| DG | RNC | Return if no cary（carry flag＝ |
| CD | RNZ | Return if not zero（zero flag＝ |
| FQ | RP | Return if plus（sign flag＝ |
| E8 | RPE | Return if parity even（parity flag＝1） |
| ED | RPO | Return if parity odd（parity flag＝$)$ |
| C8 | RZ | Return if zero（zero Elaq＝1） |

Notice that if a flag has been affected (set or reset) by some instruction, it will retain that value until and unless some other instruction is executed which affects it. For example, the INX <reg pair> (increment a register pair) instruction affects none of the flags. So this sequence:

| DCR | C |
| :--- | :--- |
| INX | H |
| INX | H |
| JNZ | <address> |

would have exactly the same effect as

| INX | H |
| :--- | :--- |
| INX | H |
| DCR | C |
| JNZ | 〈adaress〉 |

## The Compare Instruction

The contents of the Accumulator may be compared to the contents of another register by means of the CMP <reg> instruction. The CMP <reg> instruction will not change the contents of either the Accumulator or the other register, but it will affect the statuses of all the flags. However, we will usually be concerned with only two of them: the Zero flag and the Carry flag. The CMP <reg> instruction will set the Zero flag only if the two numbers (the contents of the two registers) are equal. Otherwise, the Zero flag will be reset. Thus, a typical loop structure which we used last time,

could alternatively be re-coded in this rather cumbersome form:


Obviously, such a program, although workable, is not very clever. Here is a clever use of CMP:

## CMP A

This instruction compares register $A$ to itself. Thus, the Zero flag is necessarily set, since A is necessarily equal to A. This might be a useful way of setting the Zero flag unconditionally-in case you ever need to do that.

Not only can the contents of a register be compared to $A$, but a byte of immediate data may also be compared to A. Remember that MOV A,<reg> will move (i.e., copy) the contents of <reg> into the Accumulator, whereas MVI A, <num> will load the Accumulator with the number specified, which is called a byte of immediate data. Similarly, the instruction CPI <num> will compare <num> to the contents of the Accumulator. So we could write the loop structure above in still another (but equally cumbersome) way:


The CMP and CPI instructions affect the Zero flag. But they affect another flag as well, the Carry flag. The Carry flag is used for many purposes, and many instruction affect it. Just why and how the Carry flag is affected by the compare instructions is a more involved topic than we have space for here. But we can note one rule of thumb for making use of the result: if the contents of the register (or the byte of immediate data) being compared to the contents of the Accumulator are greater than the contents of the Accumulator, then the Carry flag is set (i.e., =1); otherwise the Carry flag is reset ( $=0$ ). Suppose, for example, that the Accumulator contains OAH (=10 decimal):

| Accum. | Instruction | Carry Flay |  |
| :---: | :---: | :---: | :--- |
| ØAH | CPI | 9 | $\emptyset$ (no carry) |
| ØAH | CPI | @AH | $\emptyset$ (no carry) |
| $\emptyset A H$ | $C P I$ | $0 B H$ | 1 (carry) |

Thus, the instruction CPI <num> (or CMP <reg>) can be interpreted as a test for greater than: is <num> greater than the contents of the Accumulator? (Or, is the number in <reg> greater than the contents of the Accumulator?) If so, the Carry flag is set. If not, then there is no carry. There is no instruction which tests for a less than relation, although we can do the same thing by a series of conditional jumps. For example:

| CMP | B | ; Compare contents of B ; with contents of A. |
| :---: | :---: | :---: |
| JC | <address> | ;Jump if $\mathrm{B}>\mathrm{A}$. |
| JZ | <address> | Jump if $\mathrm{B}=\mathrm{A}$. |
|  | ; At | nt we know that |

## Another Way to Put Stuff Onto the Screen

In the last installment we developed a handy way of putting a lot of information on the screen. We did this by constructing a loop. Here it is again:

|  | MVI | C, 0 DH | ; Put count of bytes into C . |
| :---: | :---: | :---: | :---: |
|  | LXI | H,9000H | ;HL point to start of bytes. |
| LOOP: | MDV | A, M | ; Get a byte. |
|  | CALU | 33H | ;'Plot' it. |
|  | INX | H | ; Point to next byte. |
|  | DCR | C | ;Done with all bytes? |
|  | JNZ | LOOP | ; Not yet. |
|  |  | - | ;Yes. All Done. |
|  |  |  |  |

Notice that I have taken the liberty of calling the address in the conditional jump "LOOP". If the MVI C,ODH instruction is at $8200 H$, then JNZ LOOP would translate as JNZ 8205 H . But I don't want to have to figure out all those addresses each time I give you an example program, so I'll just label a certain spot in the program and then you'll know what I mean by JMP LABEL, or JNC LABEL, and so on. This is the convenience that an assembler allows you. (The MLDP does not.) When you write something like "LOOP:" out to the left (always followed by a colon), the assembler understands that it means whatever address that is. Then, when you write somethi:g like JNZ

LOOP, it will translate "LOOP" into whatever address the label "LOOP:" was at. You can have as many labels as you wish, but no two labels may be identical.

Anyway (I got off the track), by using the compare instruction we can now generalize things a bit. Notice that, in order for the program above to operate properly, you first have to load the $C$ register (or whatever register you use for the counter) with the number of bytes in the string. But suppose you don't know (or are too lazy to count) how many bytes there are? In that case, you specify that a certain number will represent the termination byte--it will indicate the end of the string. For example, let's pick the number 239 to represent the end of the string of bytes. Now our program doesn't have to know how many bytes to print--it will just keep CALLing 33 H until it finds 239. Thus:

|  | LXI | H,9000H | - HL point to start of string. |
| :---: | :---: | :---: | :---: |
| LOOP: | MOV | A, M | ;Get a byte. |
|  | CPI | 239 | ;=239 (terminal byte)? |
|  | JZ | EXIT | ;Yes. Jump out of loop. |
|  | CALL | 33 | ;No. 'Plot' the byte. |
|  | INX | H | ; Point to next byte. |
|  | JMP | LOOP | ;Back for more. |
| EXIT: |  |  |  |
|  |  |  |  |

The data at address 9000 H will be a series of numbers to be plotted (just as in BASIC's DATA statement), and will end with the number 239. The disadvantage of this scheme, however, is that you cannot PLOT 239. Of course, if you wanted, for some reason, to plot 239 , then you could pick some other number to represent the terminal byte. Zero, perhaps:

| LOOP: | LXI | H,9000H | ;HL point to start of data. |
| :---: | :---: | :---: | :---: |
|  | MOV | A, M | ;Get a byte. |
|  | CPI | 0 | ; Is it $=\emptyset$ (terminal byte)? |
|  | JZ | EXIT | ; Yes. Jump out of loop. |
|  | CALS | 33H | ;No. 'Plot' the byte. |
|  | INX | H | ; Point to next byte. |
|  | JMP | LOOP | ;Back for more. |
| EXIT: |  | . |  |

But in this case you couldn't plot a zero.

In BASIC, the PLOT statement handles not only graphics (point plot, bar graphs, etc.), but also letters and digits. You need an ASCII table to tell you what number represents what character. (An ASCII table is provided for you in your Programming Manual. There is also one included with the Machine Language Debug Package Manual.) Thus, PLOT 65 is the same as PRINT "A". And so

$$
\begin{array}{lll}
\text { MVI } & \text { A, } 65 & \text {;Put decimal } 65 \text { into A. } \\
\text { CALL } & 33 \mathrm{H} & \text { iPrint an 'A'. }
\end{array}
$$

will be the assembly language equivalent. Since the letters, digits, punctuation marks and special characters printed on the screen range from ASCII 32 through 127, a number out of that range, such as 239, can be used as a terminal byte for a string of printable characters. It so happens that there is already such a routine in your computer's ROM, and it expects 239 as the terminal byte (which is why I chose it for my example). Get out your MLDP program and examine the routine which starts at address 33 F 4 H (for V 6.78 machines; 182AH for V8.79 and V9.80). Let's name this routine OSTR (for Out STRing). It is easy to use. Just set up HL to point to the first byte of the string of numbers to be plotted and then CALL OSTR. Thus,

$$
\begin{aligned}
& \text { LXI H,9006H } \\
& \text { CALL OSTR }
\end{aligned}
$$

If you have stored these bytes at 9000 H ,

| Address | Contents <br> (hex) |
| :---: | :---: | :---: |
| (decimal) |  |

then the result will be
IT WORKS!

Try the program using the MLDP. (Remember to set a BREAKPOINT after the CALL instruction so that the program will stop and transfer control back to the MLDP。)

## CAIL and RETurn

We've had enough experience now with the CALL instruction to feel comfortable using it. But what does it do? It does the very same thing as the JMP instruction, but with one clever addition: the address of the next instruction after the CALL will be saved in a reserved area of memory called the stack. The subroutine which is CALLed will be able to return to the part of the program which CALLed it by using the RET instruction. That instruction, too, does the same thing as a JMP with a clever addition. Un1ike JMP or CALL, RET does not specify an address to return to. How does it know where to go? Simple: it merely jumps to the address contained in the stack area. Let's look at an example.

| ALPRESS | Contients | misiraction |  |
| :---: | :---: | :---: | :---: |
| 8500 | 21 | LXI | H,7000 |
| 8501 | 010 |  |  |
| 85.1 | 90 |  |  |
| 8503 | 0 E | MVI | $\mathrm{C}, 40 \mathrm{H}$ |
| 8504 | 48 |  |  |
| 8595 | 3E | MVI | A,20H |
| 8506 | 20 |  |  |
| 8507 | ØCD | CALI | 8832H |
| 8508 | 32 |  |  |
| 8509 | 88 |  |  |
| 850A | 78 | MOV | A, B |
| - | - |  | - |
| - | - |  |  |
| $8 \cdot$ | 7 |  |  |
| 8832 | 77 | MOV | M, A |
| 8833 | 23 | INX | H |
| 8834 | 9D | DCR | C |
| 8835 | 9C2 | JNZ | 8832H |
| 8836 | 32 |  |  |
| 8837 | 88 |  |  |
| 8838 | 日C9 | RET |  |

This little program has the effect of storing 20 H ( $=32$ decimal) at 40 H ( $=64$ decimal) memory locations begining at address 7000 H . When the CALL 8832 H instruction is encountered at address 8507 H , the address of the instruction following the CALL 8832 H is stored "on the stack" (as they say). That address is 850 AH . Then a jump is made to address

8832H where the program continues. It loops around until the contents of the C register become zero, and then it RETurns; that is, a jump is made to the address contained on the stack, namely, 850 AH . Pretty neat, eh?

## The Stack and its Pointer

The stack itself is merely a designated section of memory. The programmer may designate which section of memory is to be used as the stack area. The 8080 has a number of instructions which, either explicitly or implicitly, store or retrieve data in that area. The CALL and RET instructions, for examples (plus all their conditional variations, such as CNZ, CC, RZ, etc.) make implicit use of the stack. On the other hand, some instructions, such as PUSH <reg pair> and $\operatorname{POP}$ <reg pair>, which I'll discuss in a moment, make explicit use of the stack.

The stack may hold lots of data at once: it is not limited to storing only one address (or, more generally, two bytes). The stack is sometimes more descriptively called a "push down stack", or sometimes a "first-in, lastout" stack. The common analogy is a stack of plates in a cafeteria: the cafeteria staff adds a clean plate to the stack, which sinks down a bit on a spring loaded holder. A customer comes along and takes a plate (the top plate, please; don't be difficult), and the whole stack rises up a bit. Just how many plates you can put on the stack depends on the size of the ho1der. And just how many plates can be taken off the stack depends on how many plates are on it. The nature of the physical universe determines which plate is the top one.

The analogy is not bad, but don't get carried away by it. The computer's stack doesn't move at all. Rather, there is a pointer to the most recent addition to it. And instead of rising up, like the stack of plates, the computer's stack grows downwards (where "down" means toward lower addresses; so you can't make the stack grow upwards by
turning your machine upside down). More accurately, the stack pointer is adjusted downward when a new addition to the stack is made, and it is adjusted upward when something is retrieved from the stack.

Now I'm going to confuse you. I have been writing programs (such as the example above) downwards toward the bottom of the page; but the addresses increase--they go upwards. That's quite ordinary: you do the same when writing a BASIC program:

10 REM
20. $A=A+1$

30 .
But now I want to illustrate a downward growing stack, and I am tempted to make the lower addresses actually be lower on the printed page, like this:

$$
\begin{aligned}
& 9327 \\
& 9326 \\
& 9325
\end{aligned}
$$

But I won't do that. So I'm going to illustrate a downward growing stack by a list of upward growing addresses which go downward on the printed page. Got that? Nevermind. You would probably have been better off had I not mentioned this at all. Just remember that up is down and down is up, and that there is a gross equivocation going on. [Ed note: At this point the author succumbed to a fit of conversion hysteria and had to be locked up for two weeks inside a mass storage drum filled with octal notation. Fortunately, he recovered just before this month's deadline and was able to finish this article.]

Back to the stack pointer. Only two bytes (which could represent data or an address)--no more, no less--are put onto or taken off the stack during a stack operation. That is, although the stack area might be as big as you please, the operation of adding something to or retrieving something from the stack always involves exactly two bytes. One of the 8080 's registers is called the Stack Pointer (SP). It is a two byte register which always contains the ad-
dress of the most recent two byte addition to the stack. Suppose, for example, that the stack begins at 9327 H and grows down from there:

| Stack Pointer | Address | Contents |
| :---: | :---: | :---: |
|  |  |  |
| 9327 | 9324 | x |
|  | 9325 | x |
|  | 9326 | x |
|  | 9327 | x |

If a CALL instruction is encountered in a program, the address of the next instruction is PUSHed onto the stack, and the stack pointer is adjusted down by two bytes. In the example program given earlier, the CALL 8832 H would have this effect: decrement the SP; store the high byte of the address of the next instruction after the CALL in the memory location addressed by the SP; decrement the SP again; store the low byte of that next instruction in the memory location pointed to by the $S P$; then jump to the address given in the CALL instruction.

| Stack Fointer | Address | Contents |
| :---: | :---: | :---: |
|  |  |  |
| 9325 | 9324 | x |
|  | 9325 | 9 A |
|  | 9326 | 85 |
|  | 9327 | x |

The RET instruction retrieves (POPs) an address from the stack: get the low byte from the memory location addressed by SP; increment the SP; get the high byte; increment the $S P$; jump to the address so retrieved.

| Stack Pointer | Address | Contents |
| :---: | :---: | :---: |
|  |  |  |
| 9327 | 9324 | $\mathbf{x}$ |
|  | 9325 | $0 A$ |
|  | 9326 | 85 |
|  | 9327 | $x$ |

Notice that the contents of the stack are not changed by the RET instruction. Rather, the $S P$ is merely repositioned, i.e., points to a new address.

## PUSH and POP

Explicit (as opposed to implicit) manipulation of the stack pointer and the stack's contents is possible with
some of the $8080^{\prime \prime} \mathrm{s}$ instructions. The two most common are PUSH <reg pair> and POP <reg pair>. During a PUSH instruction the stack pointer is adjusted down two bytes and the two byte register pair
is copied onto the stack. The contents of the register pair are not changed. During a POP operation, the two bytes on the stack are copied into the specified register pair and the SP is adjusted up two bytes. Thus,

$$
\begin{aligned}
& \text { LXI H,2395H } \\
& \text { FUSH H }
\end{aligned}
$$

has the effect of storing 2395 H on the stack and adjusting the SP downward two bytes. And

POP H
loads the register pair HL with the two bytes on the stack and adjusts the stack pointer back up two bytes. Notice that the PUSH and POP instructions don't care what is in a register pair or on the stack. Moreover, there is no record kept about where the information on the stack came from. Thus,

```
LXI H,2395H
PUSH H
POP B
```

will (1) load HL with 2395 H , (2) copy the contents of HL onto the stack and adjust $S P$ down, and then (3) copy the contents of the stack into $B C$ and adjust SP back up. The final result is that $S P$ will be where it was to start with and both HL and BC will contain 2395 H . If you want to get two bytes off the stack and still keep them on the stack, then

POP <reg. pair>
PUSH <reg. pair>
can be used.
One of the main functions of PUSH and POP is the temporary saving of register contents. If you want to make use of $D E$, for example, but you also do not want to lose what's presently there, then here's how to do it:

| PUSH D <br> blah | ;Save DE on stack. |
| :---: | :---: |
| blah |  |
| blah |  |
| POP D | ;Retrieve DE. |

Quite often a subroutine will make extensive use of the registers, and in order to avoid destroying important data being held in those registers, the subroutine might save some or all the registers on the stack and retrieve them just before RETurning:

(The PSW, remember, is the one byte Flags register along with the Accumulator.) Notice that the registers must be restored in the opposite order they were saved, otherwise the contents of some of the registers would be interchanged. (Sometimes that is done on purpose.) Notice also that the number of POPs must equal the number of PUSHes. Why? Because
when the subroutine is CALLed, a return address is automatically PUSHed onto the stack. Then the subroutine PUSHes some registers onto the stack. In order to get that return address back again, all those two byte additions to the stack must be cleared away first. (That's why we call it a first-in last-out stack.) If you POP too few, or POP too many, what you get as the return address with the RET instruction might be very surprising.

POP QUIZ: What will this series of instructions do?

LXI H, 8 A94H
MUSH H
RET
RET
Answer: It will have the same effect as JMP 8A94H.

NEXT TIME: Manipulating the stack pointer; more ways of getting data to and from memory; addition and subtraction; getting data from the keyboard.

## $8 \mathrm{~K} \quad \mathrm{RAM} \quad$ BOAR D

(For V6.78 Compucolors only.V8.79 available soon.)
This 8 K of additional RAM is addressed at $4000 \mathrm{H}-5 \mathrm{FFFH}$, the space unused by your compucolor and available for PROM. Now you can add additional RAM instead, allowing you to increase the maximum RAM of your machine to 40K. Installation is easy, requiring a small modification to the logic board. It is compatible with the Frepost Computers, Inc. Bank Select ROM Board (see notice in this issue).

US $\$ 65.00$ plus $\$ 2.50$ postage and handing
Tom Devlin
3809 Airport Road Waterford, MI 48095

For more information or a copy of the installation instructions, send a SASE. Also see the article in the next Colorcue.


## Back Issues Sale

Back issues of Colorcue are an excellent source of information about Compucolor computers, ISC computers, and programming in general. Interviews, interesting articles, and programs are all there with a touch of history.

The list below includes every Colorcue ever published. If it's not on the list, then there wasn't one.

RETROVIET: Vol. 3, \#1 (Dec 79/Jan 80) includes: an interview with Bill Greene; Compucolor-teletype interface; user group hotline; introduction to the Screen Editor; PEEKing at BASIC programs; talking to other computers; making programs compatible with V 6.78 and V 7.89 software; software modifications.

MULTI-ISSUES at $\$ 3.50$ each
__ Oct, Nov, Dec 1978 _ Apr, May/June 1979
_ Jan, Feb, Mar 1979 —_Aug, Sept/Oct, Nov 1979
IndIVIDUAL ISSUES at $\$ 1.50$ each

| Dec 1979/Jan 1980 | _- Feb 1980 | Mar 1980 |
| :--- | :--- | :--- |
| _ Apr 1980 | May 1980 | _ Jun Jul 1980 |

IADIVIDUAL ISSUES at $\$ 2.50$ each
_ Dec 1980/Jan 1981
_ Aug/Sep 1981
_ Oct/Nov 1981
Dec 1981/Jan 1982
postage
US and Canada -- First Class postage included.
Europe, S. America -- add $\$ 1.00$ per item for air, or $\$ .40$ per item for surface.
Asia, Africa, Middle East -- add $\$ 1.40$ per item for air, or $\$ .60$ per item for surface.

DISCOUNT
For orders of 10 or more items, subtract $25 \%$ from total after postage.

ORDER FROM:
Colorcue
Editorial Offices 161 Brookside Dr. Rochester, NY 14618


BULK RATE U.S. POSTAGE<br>PAID

## Colorcue

Editorial Offices
161 Brookside Dr.
Rochester, NY 14618
Address Correction Requested

An Publication

# Colorcue 



Apri1/May 1982

## Colorcue

## A Bi-monthly publication by and for Intecolor and Compucolor Users

Editors:
Ben Barlow
David B. Suits
Compuserve: 70045,1062

April/May, 1982
Volume 4, Number 5

## 3 Editors' Notes

5 8K RAM Board for the Compucolor II, by Tom Deviin Expand your machine to 40 K

## 12 Classified Ads

13 In and Out of the Compucolor II, by Jane Devin How to take the back off--and put it on again

17 Combine Record Documentation with Record Access, by Alan D. Matzger<br>Improve your random file routines

18 About Your Subscription
Keep COLORCUE alive!
19 Assembly Language Programming, by David B. Suits Part V: Creating an INPUT routine

Advertisers: You will find our advertising policies attractive. Write for details.

Authors: This is a user-oriented and supported publication. Your articles/tips/hints are required to make it go. Send your articles or write for information.

Colorcue is published bi-monthly by Intelligent Systems Corporation, with editorial offices in Rochester, New York. Editorial and subscription correspondence should be addressed to the Editors, Colorcue, 161 Brookside Dr., Rochester, NY 14618. Product related correspondence should be addressed to ISC, 225 Technology Park, Norcross, GA 30092, ATTN: Susan Sheridan. Opinions expressed in by-line articles are not necessarily those of the editors or of ISC. Hardware/software items are checked to the best of our abilities but are NOT guaranteed.

## Editors' Notes

## Aboard a Proud Ship

Some Compucolor/Intecolor owners, for various reasons, have abandoned their first machine and have struck up friendships with the Atari 800 or Apple II microcomputers. (I have yet to hear of an ISC computer owner switching to TRS-80, or even to the Color Computer.) One may wonder about the wisdom of such a switch: what those other machines offer that the Compucolor/Intecolor does not is, to my mind, so miniscule as to not make up for their serious deficiencies, not the least of which is their little toy keyboards. And I suspect that the greater resolution available for graphics on the Apple II is hardly a sufficient motive for discarding the otherwise superior display which the ISC machines have. Nor is the 6502 microprocessor, which the Atari and Apple II computers use, a step up from the old klunker 8080.

No, it is not, I suspect, the hardware which entices an ISC computer owner to abandon ship; it is, rather, the wealth of independent software and hardware vendors who flood the market with unimaginably diverse products for use with the Atari and Apple II computers. And if you're considering joining the ranks of the produc-ers-sellers, the market for Atari and Apple. II software is still very strong. How can a budding author hope to make that kind of money in the restricted compucolor/Intecolor market?

ISC fans have been tempted away even more by the recent introduction of the IBM Pexsona: computer. I admit that the PC appears to be a very fine machine. But does it really represent very much oí a step up from the Compucolor/Intecolor? Sure, there's an 8088, but it grinds along at a blindingly slow pace. It is instructive to read Vol $I$, No. 1 of $P C$ Magazine, a new publication dedicated to the IBM PC. Its editor, Jim Edlin,
is pictured with two computers: the new IBM PC is in the background, and in the foreground is a ... a Compucolor II! "I have", Jim says, "what I believe was the best personal computer of the pre-IBM era, though you've probably never heard of it. It's called a Compucolor II...." After expressing some doubts about the PC in its present form, Edlin says, "I'm not quite ready to put my Compucolor II away. But I can see it won't be long." He anticipates an evolution of the IBM PC, by independent vendors if not by IBM itself.

Still and all, $I$ wonder if it isn't mere hype to talk about "the pre-IBM era", as though the IBM PC is really a significant advance. For my own part, $I$ am not at all interested in the new machines which--good as they surely are--offer not very much more than popularity over the Compucolor II. I will not be ready to put away my Compucolor II until I finish the design and construction of my own "step up". Don't hold your breath; it will be a 68000 -based high resolution color graphics system. The sheer size of the thing boggles my mind: the display refresh memory alone will occupy 153 K bytes.

No, it will be quite some time before I find a better machine to move up to. In the meantime, my Compucolor II, for all its foibles and follies, does me right proud every time.
--DBS

## New Publication

The U.K. Compucolor Users' Group has launched publication of their newsletter, Complolour. April was their first issue. It was 40 pages, including: Bill Donkin's lengthy discussion of how the directory and BASIC programs (or "programmes", if you prefer) are stored on disk; notes on programs in the group's library; notes on disk drive speed and alignment adjustments; an index to colorcme from Volume 1 through Volume 3; and Dale Dewey"s article on the TMS 5501. The Secretary/Librarian is Bill Donkin. Contact the group at 19 Harwood Avenue, Bromley, Kent, BR1 3DX, England. re

```
RENAISSANCEMARKETINGMNONES
    LOWCOSTGUSINESSPROGRAMS
    FOR YOUR COMPUCOLDR MODEL 4, MODEL S
* GENERAL LEDGER ( 16K.32K) ...... 55.95
(1 ) - ACCOUNT PROGRAM : DISPLAY ACCOLNT DATA : LIST ALL ACCOUNTS , ADD
    ACCOUNT, DELETE ACTOUNT, CHANGE ACCOUNT DATA.
(2 ) - JOURNAL PROGRAM: ENTER JOURNAL DATA r DISPLAY JOURNAL DATA gY ENTRY #*
    CHANGE ENTRY DATA.
(3) - PROOF PROGRAM : DISPLAY OR PRINT JOURNAL ENTRY PROOF SHEETS WITH
    TITLES.
(4 ) - POST PROGRAM : APPLIES JOURNAL ENTRIES TO ACCOUNTS.
(5) - REPORT PROGRAM: PRINT BALANCE SHEET: PRINT INCOME STATEMENT WITH
    TITLES.
INCLUDES : EASY TO UNDERSTAND INSTRUCTIONS: SAMPLE ACCDUNTS, SAMPLE ENTRY
                        FORMS: SAMPLE PRINTOUTS.
* INUENTORY CONTROL (1GK:32K) ...... $ 34.95
GIVES THE FOLLOMING REPORTS.
```



```
PROUIDES QUANTITIES AND DATA FOR 750 ITEMS OR MODELS.
* INUOICE CONTROL ( IEK, 32K ) ....... $ 49.95
    INUOICE CONTROL PROGRAM PROUIDES DATA FOR 7GB INUOICES PER DISK PER SIDE.
GIUES ON A DAILY BASIS MONTH-TQ-DATE, YEAR-TO-DATE INUOICE INFQRMATION ON TOTAL
SALES, ACCOUNTS RECEIUABLE, AMDUNT PAID, TAX DUE, FREIGHT CHARGES, INUOICE COST
AND PROFIT.
DISPLAYS = LIST ALL INUOICES, SEARCH BY CUSTONER, INUOICE DATA, SEARCH BY P.O.
    #, LIST OF TOTALS ( TOTAL SALES, SALES TAX, FREIGHT AND COST ).
PRINTS = ALL INUOICE DATA ( MONTHLY, YEARLY), A/R DATA ( BY CUSTOMER, DR ALL
                        INUOICES MONTH-TO-DATE, YEAR-TD-DATE), INUOICE TOTALS (M-T-D, Y-T-D
                        AMOUNTS PER INUDICE AND TOTALS FOR 5ALES, TAX, FREIGHT, COST AND
                        PROFIT).
                            CALL OR WRITE FOR ADDITIONAL INFO.
OPTIONAL INUOICE PRINTING USES NEBS TRACTOR INUOICE FORMS M . .....
MAXELL MINI DISK FOR COMPUCOLOR. BOX OF 10... \(\$ 34.95\)
\(30 \%\) OFF ALL COMPUCOLOR CORP. SOFTWARE IN STOCK. CALL
[] INUOICE CONTROL \$ 49.95
[] GENERAL LEDGER \$ 59.95
[] INUENTORY CONTROL \(\$ 34.95\)
[] MAXELL DISK BOX QTY......
TERMS: PAYMENT WITH ORDER

\title{
8K RAM Board for the Compucolor II
}

\author{
by Tom Devlin \\ 3809 Airport Road \\ Waterford, MI 48095
}

For some time \(I\) had been casting a wistful eye at the 4000-5FFF "Future Space" section of the Compucolor II memory map. ISC left this area open to give us the option of installing 8 K of EPROM for special programs. They mounted connectors J9 on the logic board to mate with an EPROM board and programed PROM UB3 to supply chip select signals to this board during memory read cycles. Unfortunately very little use has been made of this feature.

My first thought upon discovering this unused area was that it would be a great spot for some RAM. This would be far more flexible than EPROM in that with RAM the program could be changed at will or it could be be used for temporary storage if the application required it. One of the nicest features of placing RAM in this area would be that BASIC wouldn't know it existed so you could (ESC)-W without worcing about your utility program.

明y search for suitable RAM chips turned up the rochlis 2016, a \(2 \mathrm{~K} \times 8\), low pover, static device. Four of these provided the Eull \(8 k\) needed and occupied the same phyetical space as the 2716 Erroms used on the Compucolor add-on board.

The only problem was the address decoding. ISC had intended this area for EPROM and had made no provision for accessing the chips during a write cycle. Rather than replace the UF3 system decoder PROM and the UB3 chip select PROM, I decided to disconnect the appropriate outputs of the UB3

PROM and run the Al2-Al5 address lines through the same J9 connector pins. This was easy since Al2-Al4 were already present on pins \(11-13\) of UB3 and Al5 was available on pin 36 of the 8080 CPU just a short distance away. These four lines and the AO-All lines already present on \(J 9\) allowed access anywhere in the 64 K memory range. While \(I\) was mainly interested in the \(4000-5 F F F\) area I left the door open for future expansion into the screen memory area via pads 0 thru \(V\) on the board. If all goes well this will be a part of another Colorcue article.

At this point I had almost everything needed. The only thing left was to pick up the active low memory read (MRD) and write (MWR) strobes. Two wires from convenient locations on the \(\operatorname{logic}\) board to the new RAM board completed the modifications.

The schematic shows the final circuit. I decided to make the board accept single 5 volt supply (Intel type) EPROMS as well as the 2016 RAM chips. Each board will accept 8 K of 2016 RAM or 2716 EPROM or 16 K of 2732 ERROM in two switch selectable 8 K banks. The schematic shows the jumper and switch settings for each type of memory chip. I also made provision for stacking these cards and selecting between them with a switch. Jumpers \(A-B\) and \(C-D\) allow for easier stacking by re-routing the MRD and MWR signals from the lowest (bottom-most) board to the upper one through connector pins 13 and 14 . If you use this
option, you MUST cut off pins 13 and 14 on the bottom of the lower board or serious damage to the computer may result because the logic board uses these pins for +12 and -5 volt power.

Address decoding and chip selection are handled by IC-1 (74LS138), a three line binary to 1 of 8 decoder. Assuming that \(\mathrm{IC}-1\) is enabled by a logic zero (ground) through one of the two diodes, one of its outputs will be low through a 2 K or 4 K block within the \(4000-7 \mathrm{FFF}\) range. The size of the block has to match the size of the memory chip installed and is set by jumpering M-L or M-N. (If IC-l is not enabled, or we are outside the 4000-7FFF range, all eight outputs will remain high.)

Pin 18 of any of the three types of memory chips is the active low chip enable (CE). If this pin is pulled low, the chip "wakes up" and awaits further orders. If pin 20 , the active low output enable (OE), is then pulled low by MRD, the memory chip will output a byte onto the data bus. In the case of the 2016 RAM chip if pin 21 , the active low write input (W), is pulled low by the MWR strobe, the enabled chip will take the byte currently on the bus and store it internally. If CE goes low, but neither the MRD or MWR follow (as might happen during an I/O cycle) the enabled chip will do nothing.

Pin 21 has other functions on the EPROM memories. The 2716 uses it for the programming voltage (VPP). For normal operation this pin must be heid at +5 volts. The 2732 uses pin 21 for the All address line to double the storage capacity to 4 k bytes.

Resistors R1 and R2 simply pull the IC-1 inputs to a logic one (+5 volts) when the switches are open. The 1 N270 diodes specified are germanium devices chosen for their lower forward voltage drop compared to the more common 1 N914 silicon types. 1 N 914 s should work fine with the switches (in fact D2 is not needed if you use switches for control) but the lower drop of the 1 N 270 s will give some extra margin for possible future card selection under software control. The capacitors are for power supply by-passing.

Installation on the V 6.78 units is straight-forward. Just follow the REV. 3 instructions. The situation is a little more complex with the V8.79 types: the REV. 3 logic board was changed to a newer REV. 4 board part way through production. Also, the initial run of V8.79 UA5 FCS ROMs must have had a bug because Compucolor patched around one section of this ROM to an EPROM located on the add-on EPROM board (the only time this board was ever used) and changed the UB3 PROM to reflect the change. If you have one of these units with the add-on board installed you will have to update to the later UA5 ROM (PN 100695) and replace the UB3 PROM so you can remove and discard this EPROM board and use the J9 connectors for the RAM card. I can supply a preprogrammed 2532 EPROM to replace the UA5 ROM and the proper UB3 PROM if needed.

The REV. 4 logic board used in the final units is mostly the same as the REV. 3 board (a few changes to the RS232 circuitry and the handshake mod added), but for some reason ISC switched from the 16 pin UB3 PROM used on the older boards to a 20 pin part. I have been unable to find a REV. 4 board to make installation drawings from. If anyone can supply a good photograph or drawing of this board I will try to get the information into a future Colorcue.

If you elect to wire wrap a RAM card, the PC layouts and assembly drawing should, with the schematic, give you all the information needed. The 14 pin header assemblies are a little tough to find, but you can make your own by cutting a 28 pin wire wrap socket in half lengthwise. Just be sure that the socket you pick has large enough holes to allow you to stack another board on top. Also, on some sockets the pins are offset slightly from the holes. If yours are like this, make sure you keep the hole center-to-center spacing at \(3.9^{\prime \prime}\) to make stacking easier. (The Compucolor J9 spacing varies a little from unit to unit. The . \(65^{\prime \prime}\) pins on the headers we use are long enough to flex slightly to make up for this. I would recommend using this length.)

If you do decide to roll your own,

would you please drop me a line and tell me how you made out?

\section*{Installing the RaM Card on Rev. 3 Logic Boards}

A word of warning. STATIC CHARGES KILL MOS ICs. Just walking across the floor can produce enough of an electrical buildup to totally destroy MOS parts. The RAM chips are just as sensitive as other parts of your machine, so remember this warning whenever you work on your computer. To avoid this potential problem touch the metal frame of the disk drive to neutralize any charge you may have accumulated. Do this as soon as the back is off.
(1) Disconect the A.C. power cord and all connectors from the rear of the computer.
(2) Remove the back (3 screws top, 1 bottom). Be careful not to hit the neck of the CRT during this step.
(3) Being careful to mark socket placement, remove the three cable connectors and the sixteen pin flat cable for the disk drive from the left side of the logic board (bottom circuit board). Remove the logic board.
(4) Locate and remove PROM UB3 (82S123). (Figure 1A.) Bend pins \(1,2,3\), and 4 out from the body of the chip. Make sure that the pins don't touch each other or anything else when the chip is replaced in the socket. If UB3 is not socketed you can cut these pins off close to the board with a sharp pair of wire cutters.
(5) Turn the logic board upside down and locate the socket pins for UB3. Using four lengths of insulated jumper wire, make the following connections (see Figure 2):
( ) pin 2 to pin 13
( ) pin 3 to pin 12
( ) pin 4 to pin 11
( ) pin 1 to UA2 socket pin 36 (8080, A15 output).
(6) Turn the logic board rightside up. Plug the RAM board into connectors J9. The end with the small parts goes to the rear (edge card connector end) of the logic board. Support the logic board from the bottom to keep it from flexing during this step.
(7) Connect a wire from the MWR pad on the RAM board to pin 8 of \(J 8\) at the front end of the logic board. You can connect directly to pin 8 of the 16 K add-on RAM card if you have it, otherwise use a wire wrap pin. See Figure 3.
(8) Locate the feed-thru pad shown in Figure 1. Connect a wire from the MRD pad on the RAM board to this point.
(9) Double check your work. Pay close attention to che UB3 and UA2 chips to be sure that you have not shorted two adjacent pins together when you soldered the wires to them.
(10) Re-assemble the computer. Remember to reconnect all of the cables removed during the disassembly. The logic board is supported at the rear by two tabs that fit into si ts in the rear cover. Be sure to yet these in place.

With the power switch off reconnect the keyboard and power cords. Take a deep breath and turn the power on. If the computer does not act normally, turn the power off immediately. Otherwise relax and wait for the screen to light up.

Problems can be localized by removing the RAM card from the computer. (You do not have to "unmodify" the logic board.) If everything then works normally, you either have a bad RAM card or else you have connected the MRD or MWR wires to the wrong places.

If the display "jitters" or the colors don't quite line up don't panic; you've probably just moved the horizontal centering control. (The small trim pot at the left rear edge of the logic board.) Re-adjust it with a small screwdriver. (The display will move in steps as the control is turned. Set the control for the best picture midway between two steps.)

If all is well so far, put a disk into the drive, enter FCS, and

\section*{FCS SAAVE RTEST \(^{0} 0-1 \mathrm{FFF}\)}
to put 8 K of fairly random data onto the disk. Now try to load this data into the RAM card with

FCS \(>\) LOAD RTEST.PRG;01 4000

If you don't get any EMEM error messages, you are home free!

\section*{Software}

Some software is already available for the \(4000 \mathrm{H}-5 \mathrm{FFFH}\) area. Quality Sottware Associates (21 Dersingham Cr., Thornhill, Ontario, Canada, L3T 4P5) has a greatly expanded version of 'THE' BASIC EDITOR for use with the RAM card. It uses the full 8 K and is a joy to work with. If you have already purchased 'THE' editor, send them your disk and \(\$ 6.00\) and they will run you a copy. This program will be included in all future disks sold. You can also relocate the Compucolor MONITOR or MLDP programs using the MENUs supplied with them. (Line 710 of the MENU program prevents ORGing below 8200 H , so just deleteit.) I have included a program to reORG version 3.20 of FREDI to 4000 H .

The RAM card is compatible with Frepost Computers' bank switching EPROM system. Frepost has instructions available. If, however, you install the RAM card and later decide to install the Frepost system, you will have to return the logic board to its original condition before proceeding.

The RAM card may also be purchased assembled and tested for \(\$ 65.00\) plus \(\$ 2.50\) postage and handling. (US funds drawn on a US bank. Make checks payable to Tom Devlin.) All ICs are socketed, the printed circuit board is tin-lead plated epoxy-glass, and all parts are of high quality. It is covered by a 90 day parts and labor warranty. It is currently available for all V6.78 Compucolor IIs. Owners of V8.79 units will have to look inside your machines before ordering. If your logic board has a 16 pin UB3 PROM, you have the same Rev. 3 logic board as in V6.78

systems. If you also have an EPROM board with one EPROM on it, you will have to replace the UA5 masked ROM and UB3 PROM to let you remove and discard this EPROM board to free up the J9 connectors for the RAM card. We have a UA5 EPROM and UB3 PROM replacement kit for \(\$ 15.00\) additional when ordered with the RAM card. V8.79 units with the Rev. 3 logic board and no EPROM board already have the cor-
rect UA5 and UB3 chips installed and are ready for installation of the RAM card. If your machine has a 20 pin UB3 PROM, you have a Rev. 4 logic board. These require no extra parts, but we have been unable to find a Rev. 4 board to make the installation drawings from. Perhaps we can make this information available in a future article.


- 1982 by Tom Dev1in
```

    \emptyset ~ R E M
    l REM
    2 REM
    3 REM
    8
    REM MOVE TOP OF MEMORY DOWN
    10 POKE 32940,191:POKE 32941,158:CLEAR 50
    19
20 PLOT 12,6,2,15
30 INPUT "INSERT YOUR BASIC EDITING DISK AND HIT RETURN";A\$
40 PRINT :PRINT "DO YOU HAVE AN <ESC> [P] JUMP TO 4000H?"
50 INPUT " ( ALU V8.79 UNITS DO ) ";AS:PRINT
6 0 ~ P R I N T ~ " I H I S ~ W I L L ~ T A K E ~ A B O U T ~ T W O ~ M I N U T E S " ~
6 8
6 9 REM GET BOTH VERSIONS INIO l6K AT THE SAME TIME
7\emptyset PLOT 27,4:PRINT "LOAD FREDl6.PRG;01":PLOT 27,27
80 PLOT 27,4:PRINT "LOAD FRED32.PRG;0l,9EC0":PLOT 27,27
88
9 REM SPOT FCS ERROR VIA CCI CODE
IF (PEEK (33231) AND 7)=1 THEN 10
9 9
100
110
120
130
138
139 REM MOVE 'FREDI' DOWN, CHANGING HIGH ADDRESS BYTES
189 REM SEE IF WE NEED TO USE <ESC> [^]
190 IF ASC(AS)<>89 THEN 250:REM YES
198
199 REM DELEIE USER JUMP SETUP (FILL WITH NOPS)
200 FOR M=16847 TO 16857
210 POKE M,\varnothing

```

230 POKE 19521,80:REM CHANGE [^] TO [P] IN SIGN-ON MESSAGE
240 GOTO 260
248
249 REM SET USER JUMP TO Ø4ø0øH IF NO 〈ESC> [P]
250 POKE 16848,ø:POKE 16849,64
257
258 REM SET UP JUMP AT Ø4000 WHEN RUN
259 REM USING OLD TOP OF MEMORY ADJUST SPACE
\(26 \emptyset\) DATA \(62,195,50,0,64,33,15,66,34,1,64,0\)
270 FOR M=16864 TO 16875
280 READ D
290 POKE M,D NEXT

309 REM CHANGE THE SIGN-ON MESSAGE TO 40ø0H, COLOR TO CYAN
310 DATA \(6,52,48,48,48,72\)
320 FOR M=19468 TO 19473
330 READ D
340 POKE M,D
350 NEXT
358
359 REM SAVE 'FRED4' TO DISK AND RUN IT. (HIM?)
360 PLOT 27,4:PRINT "SAVE FRED4 4180 105D":PLOT 27,27
\(37 \varnothing\) PLOT \(27,4:\) PRINT "RUN FRED4"

\section*{Call for Information}

An interesting use of ISC computers was devised by Mr. Shepard of Rochester (present whereabouts unknown) who marketed the machines with a horse race handicapping system. His disappearance leaves those owners completely out of the saddle. If you're one of those owners, Colorcue can serve as an information exchange for you. We will be happy to publish information (if someone will provide it for us) or, if you wish, publish names of users interested in contacting each other. Let us hear from you.

\section*{CLASSIFIED ADVERTISING}

FOR SALE: Intecolor 3651. New in original carton. \(32 \mathrm{~K}, 51 / 4\) inch disk. Disks (no games). Joystick option. \(\$ 2800\) (list \(\$ 3800\) ). George Wilson (404) 458-1431 (nights).

FOR SALE: Compucolor \(\operatorname{II}, \quad \mathrm{V} 8.79,32 \mathrm{~K}\), extended keyboard, modem, manual set, disks, games. Great condition. \$1200 or best offer. Pete Gammon (404) 393-9690.

FOR SALE: Compucolor II, V6.78, 16K, 71 key keyboard, manual set, disks, including assembler, Compuwriter, Debug, games. Asking \$1000. Charles Lovejoy ( 800 ) 225-2465 x1365, weekdays 9-5.

\title{
In and Out of the Compucolor II
}

\author{
by Jane Devlin \\ 3809 Airport Road \\ Waterford, MI 48095
}

Will someone explain what I'm doing, sitting here with a screwdriver in either hand and the Compucolor crouched in front of me like a threatening beastie? I'm more at home with knitting needles, crochet hooks or quilting blocks, so the idea seems to be that if I can get into and out of this computer without major damage to it, myself, or the immediate environment, then ANYONE can. Nervous? Who's nervous? ME! (I wonder if he'll promise to count to ten if he hears a scream and a crash? That should give me a good running start!)

First, clear a space to work (as I went along, \(I\) found it had to be about three times the size of the computer to allow the back to swing out of the way and the logic board to be pulled free). If you're working on anything that you'd like to keep scratch-free, cover it with a cotton sheet or towel, or newspapers. Nothing synthetic, please! The Compucolor II has MOS chips in it, and while with reasonable care there's no danger to them in taking the unit apart, a stray static charge might fry them. I'm told it's a good idea to get all the tools you'll need together (to just take the computer apart and put it back together, all you'll need is one regular screwdriver and one phillips head) and then stay put while you're working. (No fair scuffing your feet on the carpet first).

Disconnect the \(A C\) cord from the back of the set by grasping the connector (not the cord!) and wiggling it up and down while pulling. Find the screws holding the back on; there are three phillips heads across the top and a nut with a screwdriver slot at the very bottom. This last one is
at the end of a groove that resembles the trench in Star Wars and you need either a long screwdriver so the handle extends past the back or a lot of patience. You can tilt the set forward where it will remain fairly steady while you remove the bottom screw (Fig. 1). Remember to put the screws somewhere where you can find them again.

Now, before you remove anything else, note the two tiny green PC board tabs sticking through the back on either side of the edge connectors. These will be important in reassembling the computer. Tilt the unit back


Figure 1
down on its base and look down the ventilation slits in the neck cover (that's the part that sticks out the farthest). You'll see a small board attached to the neck of a big, black glass bottle--the CRT. As you work the back off in a straight backward motion, be careful not to put any pressure on that small board and thus on the neck of the CRT. As you pull
the back free of those two PC tabs, there will be a slight jerk as the logic board drops a little; don't panic, you didn't break anything. You now have the back in your hands, but it is still connected to the computer by an 'umbical cord' of wires. Swing the back gently to the left and it will be out of the way without disconnecting any more wires than necessary.

You can now see the interior of the computer. To the novice (like me) it's a confusing jumble of metal boxes, P.C. board, and wires running everywhere. To your left as you look into the set is a rectangular metal box with a cylindrical projection and a gray ribbon cable running off the P.C. board attached to the box's left side (Fig. 2,A). This is the working end of your disk drive. The big, black, glass 'bottle' in the center is the CRT, as I mentioned before (Fig. \(2, B)\). Leave it alone. CRTs are capable of storing jolts that can make things very uncomfortable. This one has a 'bleeder' to drain high charges, but why take chanches? On the far right (and facing to the right), partially protected by an aluminum frame that doubles as a heat sink for three power transistors (those are the little oval things that vaguely resemble a streamlined flying saucer), there is an intriguing board with all sorts of pots that seeri to beg to be twitched or tweedled. DON'T--unless you like non-working computers, sleepless nights and gray hair. This is the infamous analog board, and it's awfully easy to blow up, so leave it alone!
(At this point, it's a good idea to touch the disk drive case or some other metal inside the set so you and the computer will be at the same electrical potential and thus protect the circuitry.)

The board sitting on the bottom of the case with three edge connectors pointing right at you is the logic board. This is the one that holds all the chips that make the computer compute, and it's here that most modifications are installed. For the sake of simplicity, I'm going to refer to the side with the three edge connectors as the rear of the board
(because it's in the rear of the set) and the side disappearing into the recesses under the CRT as the front of the board.

There are three wired connectors to the left side and the other end of the ribbon cable from the disk drive in a DIP socket just beside them. These four connectors have to be disconnected before the board can be pulled out, but before you start yanking, take the time to mark the connectors with masking tape, magic marker or anything handy so you know where to put them back. It can save a lot of hassle in the long run if they're keyed for replacement. (Two of ours are, and two aren't.) Carefully work the connectors free and pull the logic board toward you. (Now see why I said you would need a lot of space to work in?)

If you have a 32 K machine, there will be a funny looking (it's upside down) small card on long spindley pins toward the right front quarter of the logic board. This is the 16 K add-on RAM card and it is placed in two connector sockets known collectively as J8. To the right of this area is another set of widely spaced connector sockets (J9). If you have an EPROM board already installed, it will be in the J9 sockets and will look about the same as the 16 K RAM card.

The chips have been designated sequentially from the right rear corner. The letter designation reads from right to left across the board and the numbers from the rear to the front. Therefore, the chip in the right rear corner is UAl, and in the front left corner you will find UG1l. It makes finding chips on the board quite simple. (But when it comes to the schematic, it's everyone for themselves!)

Now for the fun part--getting everything back the way you found it. If you look at the front edge of the logic board between J8 and J9, you will see a small notch in the P.C. board. This was designed to fit onto one of the plastic fins molded into the case just under the CRT. Reconnect the four connectors, being careful to get them in the right places, and slip the board into the case. Don't expect
to be able to see to guide that notch onto the fin; you can't, and it has to be done by feel and blind luck. When it's engaged, the board will not move sideways. Let the board rest on the two rubber knobs on the bottom of the case.

Tilt the computer forward again and take a deep breath. Getting the back on has been known to drive people crazy. I struggled with trying to get everything aligned for almost ten minutes before \(I\) gave up and asked for help. Tom made it look easy, but then he's had a lot of practice. What I hadn't noticed was that there are two small wedges on the bottom edge of the back that have to fit into two notches in the case. You can sort of
look through the ventilation slits to keep track of what you're doing. Once those wedges are engaged, the back slips easily over the neck and onto the rim. The two tabs on the logic board have to be fit into their holes. As the board is slightly raised when they are in place, you have to lift it gently with a pencil or some other non-conductive object so the back can slip the final quarter inch into place. (Don't put too much pressure on the P.C. board or you may crack it.) Find where you hid the screws, tighten them in, replace the line cord and you're done!

As for me, I have an afgan, two sweaters and a quilt to finish.


Figure 2

\title{
EASY TO USE \\ ECONOMICAL
}

\section*{LENS DESIGN PROGRAMS}

FOR THE COMPUCOLOR II

The optical design aids which we offer now for the COMPUCOLOR have evolved from 20 years of experience in developing lens programs for the IBM 704, 7094, GE 235, 400, 635, ModComp 11, and HP 1000 computers. Those of you that already have a COMPUCOLOR know that it is hard to beat in what you get for the money. These programs have been written to give you the benefit of this Computer Explosion. We currently offer a set of programs, two for the First Order Optics and one for the Rigorous Ray Tracing with spot diagrams. We plan to offer an ongoing series of programs and updates.

The current programs are written for a COMPUCOLOR with 16 K memory and disc. The programs are provided on disc with a tutorial users manual in a notebook binder. The manuals are written with the beginner in mind so that a high school student can use them. The programs are written with the professional in mind so that anyone can do practical optical evaluation and design with a small investment in the computer and software.

FIRST and EFL are first order evaluation programs which take the input of the lens description in radii, spacings, and indices and compute the image positions and magnifications and plat a representation of the optical system on the screen. EFL computes the Front, Back and Effective Focal Length and shows the nodes on the Screen-plot. Mirrors and Catadioptics are handled by all programs.

TRACE is a program to trace rays rigorously through spherical and conic section surfaces with optional decentrations. The resulting array of rays can be plotted on the screen as a spot diagram and listed by coordinates. The size and position of the best focus is computed, and images of three weighted colors can be evaluated. The inputs and outputs are simple and straightforward. In an hour of learning, you can be well on your way to constructive lens evaluation and design.

TRACE SET includes: FIRST, EFL, and TRACE programs ..... Price \(\$ 250.00\) (U.S. dollars)
ENCLOSE check or money order, Corporate Purchase Orders accepted.
Price subject to change without notice.

\footnotetext{
WIILFY CORPORA:VCAN
PO B0X 670. MELBOURNE FICRIDA 32901
PHOHE (305) 727 ? 2 :3
}

\title{
Combine Record Documentation with Record Access
}

\author{
by Alan D. Matzger \\ 960 Guerrero St. \\ San Francisco, CA 94110
}

Most business applications involve keeping records, updating them from time to time and generating reports from them. Such applications involve several to many programs, each of which does something with the fields in a record--reading or modifying or writing them back. A while ago I found \(I\) was getting lost inside a record, forgetting or confusing which field was in what location. For a time \(I\) kept a formal, handwritten list of what was what in there. Things were a bit better, but lately I've hit on a method whereby such documentation is incorporated into the body of a program and is marvelously useful.

Recall the general format for the GET and PUT statements: PUT file <,record <,first>>; nexpr, sexpr [byte count]. I guess I didn't read all that too carefully at first, because I thought that all except the "nexpr" and "sexpr" had to be constants. My early programs were overflowing with GET and PUT statements for each and every field accessed, all with the appropriate numeric constants.

But one day in a flight of fancy I played with variables in place of constants and lo!, they worked splendidly. All at once, field access became a matter of a few subroutines and lots of DATA statements. And those data statements, appropriately REMmed, provided internal documentation for the byte-by-byte contents of the record. In what follows the variable \(F L\) is the file number used when the FILE " \(R\) "... opened it; \(R\) is the record number; \(S B\) is the starting byte of the field; AS or \(A\) is the receiving
variable for string or number; and \(L\) is the length (byte count) of the field.

So now the DATA statement contains the file number (if appropriate), the starting byte and field length. Thus,

100 DATA 1,1,25:REM NAME
110 DATA 1,26,20;REM ADDRESS
\(12 \emptyset\) DATA 1,46,15:REM CITY \& STATE
-
-
2ØØ DATA 2,5,4;REM AMOUNT OWED
The accessing subroutine might be:
```

4 0 0 ~ R E A D ~ F L , S B , L : I F ~ L = 4 ~ T H E N ~
GET FL,R,SB;A:GOTO 42\emptyset
4l0 GET FL,R,SB;A\$[L]
42\emptyset REIURN

```

Be careful that none of the string. fields are four bytes long; should one be, then use a different value for \(L\) to denote a numeric field. (Then be sure to remember that the actual field length is indeed four so that the SB of the next field is correct.)

The calling program RESTOREs the proper line number for the field desired. It will also somehow have a value for \(R\), the record number. (Here is perhaps the only advantage of BASIC's having only global variables.) You then simply GOSUB 400 to get the desired field. Thus, to print the City and State field,

RESTORE 120:GOSUB 400: PRINTI A\$

The PUTting subroutine may or may not need a READ statement. Mine usually don't because I GET or "pretend" to GET whenever I work on a field. The values for \(F L, R, S B\) and \(L\) remain unchanged. By "pretending" I mean that I keep other wanted information in the DATA statement, such as the X and Y coordinates of the field's position on the screen. So even though there may be nothing in the record to GET, I'll pretend and at the same time get the several field parameters.

My updating programs do usually include the \(X\) and \(Y\) coordinates and sometimes the color code; appropriate PLOT statements in the subroutine make for a pretty screen. The report programs have appropriate values for TAB positions. These extra pieces of information are READ into other vari-
ables in the one READ statement.
Having discovered that variables work as well as constants in GET and PUT statements, I think my programs are a whole lot cleaner and surely are closer to being self-documenting. But then, I suppose you knew that all along.

\footnotetext{
P.S. As long as we're talking about GETting and PuTting, here's something I found disconcerting the first time I ran across it. The LENgth of the receiving string variable equals the byte count used in the GET or PUT. A two character string GETten by \(A \$[20]\) is 20 characters long--the two plus 18 trailing spaces. It took me a while to figure out why my print lines were overflowing. ©
}

\section*{** About Your Subscription **}

Most of our subscribers' subscriptions will come up for renewal after July. (Check your mailing label: the number indicates your last issue number. The June/July issue is issue number 6.)

\section*{We need your subscriptions to continue publication.}

Since its beginnings in 1978, Colorcue has been financed by Intelligent Systems Corporation. That financial assistance will terminate with the June/July issue. Whether we will be able to continue publishing Colorcue will depend on whether you renew your subscription. There is a critical number of subscribers, below which Colorcue will not have the funds to continue. Let us know of your desire to see Colorcue flourish. Send us your renewal now so that we may know in advance where we stand and so that you can help guarantee the continued publication of what we believe to be an outstanding magazine for Intecolor/Compucolor users.

Subscription for one year (six issues) is \(\$ 12\) in U.S., Canada and Mexico; \(\$ 24\) elsewhere. Please make check or money order in U.S. funds payable to "Colorcue".

At the same time, why not take this opportunity to let us know what kinds of information you would like to see in colorcue during the coming year? Would you prefer more hardware oriented articles? Tutorials? Programs? Applications? Games? Perhaps you have something specific in mind.

\title{
Assembly Language Programming
}

\author{
by \\ David B. Suits
}

\section*{Part V: Creating an IRPUT Routine}

We have discussed several ways of getting information from memory onto the screen. Let's take up the subject of getting information from the keyboard into memory, The very first thing we need to know is how to get one character from the keyboard.

Suppose you were working in BASIC but without the INPUT statement. How would you get data from the keyboard? If you can answer that question, then you can probably devise at least one perfectly useful method in assembly language. Under normal conditions, when a key is pressed, its code is stored in location 81 FEH ( 33278 decimal), and then it is output to the screen. The codes range from 0 (control @) and 1 (AUTO) through 17 (red) and 18 (green), and then through the printable ASCII characters, 32 (space), 33 (!), and finally up through the special function keys, 240 (F0) through 255 (F15). The routine which reads the keyboard and stores the result in 81 FEH if a key was pressed is a ROM routine which is automatically executed many times per second. That is, it interrupts your program, does its thing, and then returns control to your program. This happens in BASIC and in assembly language, and it occurs so fast that
you are usually unaware of it. There are, however, several methods of altering this routine so as to make it more adaptable to your particular needs. We'll get to that later. Right now, let's look at the quick and dirty method, which can be illustrated with the following flow chart. A simple BASIC subroutine is given in Listing

1. An equivalent assembly language program is given in Listing 2.

Last time I explained that labels may be used in place of constants (such as addresses) in an assembly language program. Thus, the instruction "JZ LOOP" is interpreted by the
\(1000 \mathrm{~KB}=33278\)
1ø1ø POKE KB, \(\varnothing\)
\(1 \varnothing 20 \mathrm{~A}=\operatorname{PEEK}(\mathrm{KB}): \operatorname{IF} \mathrm{A}=\emptyset\) THEN \(1 \varnothing 2 \emptyset\)
1030 REM GOT IT!
1040 RETURN

Listing 1. A simple BASIC program to wait for a key press.
assembler as "Jump, if the zero flag \(=1\), to the instruction which is at the address LOOP:". But suppose I had not included the "LOOP:" label off to the left in Listing 2. How would the assembler koow where to make the program jump to in response to "JZ LOOP"? It wouldn't. Similarly with KBCHAR. The instruction LXI H,KBCHAR will be translated as "put the two byte number represented by KBCHAR into the HL register pair". But what two byte number is that? It is the EQU instruction which gives the assembler that information. It says, in effect, that whenever the word "KBCHAR" appears, make it EQUa1 (or EQUivalent) to--or, better yet, EQUate it with81FEH. (It is roughly similar to an assignment statement in BASIC, as in line 1000 of Listing 1 ; whenever KB appears, 33278 is meant.) By the way, since you can load the HL pair with any two byte number, and since all addresses are two byte numbers, you can even load HL with the value of LOOP. That is, the label "LOOP:" is merely an address in the program. So if we wanted to, we could have all sorts of instructions which referred to it. For example
\[
\begin{array}{ll}
\text { JMP } & \text { LOOP } \\
\text { LXI } & \text { H,LOOP } \\
\text { CALL } & \text { LOOP }
\end{array}
\]
and so on. Although you are not allowed to have duplicate labels, you may refer to any label you please as often as you please. Notice also that labels which specify the location of an instruction in your program (such as

GETCHA and LOOP) must always be followed by a colon. But assignment (that is, EQUate) statements, such as "KBCHAR EQU 81FEH", have no colon; if you put one in, the assembler will signal an error. That's because KBCHAR does not indicate the location of an instruction in your program. Rather, it is merely a convenient way of writing the number 81 FEH (or whatever number it is which you EQUate to KBCHAR). That convenience makes your programs easier to read: "LXI H, KBCHAR" is a bit more meaningful than "LXI H,81FEH". This also makes things easier to change in case you need to do so. Most, if not all, of the EQUate statements are traditionally grouped together near the beginning of the program or routine. If you were to learn that, say, 81 FEH is not after all the correct number, then you need only change the appropriate EQUate statement, and the assembler will make all the required corrections for you.

The EQUate statement is one of several instructions which are sometimes called pseudo-ops because they don't end up in the final machine language program created by the assembler. Every time KBCHAR appears, it is translated into 81FEH. But the instruction "KBCHAR EQU 81FEH" itself disappears; it is an instruction to the assembler to act in a certain way. Thus, if the assembler were to assemble the get a character subroutine in Listing 2 starting at, say, 9000 H , the machine language code it would produce would be the string of numbers in the Contents column below:
; Subroutine to get a character from the ; keyboard. NOTE: this is a workable, ; although not very efficient, way of ; doing it.
\begin{tabular}{llll} 
KBCHAR & EQU 81 FEH & \\
& & & \\
GETCHA: & LXI & H,KBCHAR & iLoad HL with 81FEH. \\
& MVI & M, & iSet keyboard char= \\
LOOP: & MOV A,M & iSee what's there. \\
& CPI & \(\emptyset\) & if zero, keep looping. \\
& JZ & LOOP & iI.e., wait for key. \\
& RET & & iReturn with key in A.
\end{tabular}

Listing 2. An assembly language subroutine to get one character from the keyboard.
\begin{tabular}{|c|c|c|}
\hline Address & Contents & \\
\hline 9000 & 21H & LXI H \\
\hline 9001 & 0 FEH & 81FEH \\
\hline 9002 & 81H & \\
\hline 9003 & 36 H & MVI M \\
\hline 9004 & Ø0 H & \(\emptyset\) \\
\hline 9005 & 7EH & MOV A，M \\
\hline 9006 & QFEH & CPI \\
\hline 9007 & 00 H & \(\emptyset\) \\
\hline 9008 & ØC2H & JNZ \\
\hline 9009 & 05 H & 9ø05H \\
\hline \(9 \emptyset 0 \mathrm{~A}\) & 90 H & \\
\hline 900 B & 6C9H & RET \\
\hline
\end{tabular}

Another important pseudo－op is the ORG（ORiGin）statement．It in－ structs the assembler that you wish the following series of instructions to be located in memory beginning at the address specified．A popular place for \(a_{i}\) sembly language programs for ISC machines to begin is address 8200 H （because that＇s where you go with ESC T）．So you＇ 11 find a number of programs beginning with＂ORG 8200 H ＂．Often there is only one ORG in the whole program，but occasionally there are more．Perhaps，for example，you want the program to live at 8200 H ，but for some reason you want a particular subroutine located in very high memo－ ry：
\begin{tabular}{|c|c|}
\hline \multirow[t]{2}{*}{} & 8200н \\
\hline & \\
\hline \multicolumn{2}{|l|}{\multirow[t]{2}{*}{；main program}} \\
\hline & \\
\hline \multirow[t]{2}{*}{ORG \({ }^{\circ}\)} & ØfAC0H \\
\hline & \\
\hline \multirow[t]{2}{*}{} & broutine \\
\hline & \\
\hline & \\
\hline End & \\
\hline
\end{tabular}

Finally，the ERD pseudo－op merely tells the assembler to stop assem－ bling．It must be the last line in your assembly language program．More－ over，it must be followed by a carriage return（to signal the end of the line on which the END occurs）．Failure to
include the carriage return is a common source of frustration for be－ ginners；the assembler will insist that it didn＇t find the END statement， eventhough it＇s right there in front of its nose．

\section*{Logical OR}

Back to reading the keyboard． （Forgive me for these frequent，but important，digressions．）I want to make the get a character routine more elegant．Several new 8080 instructions can help．The first is the ORA＜reg＞ instruction．This takes the contents of＜reg＞and logically ORs them with the contents of the Accumulator；the result ends up in the Accumulator． （The contents of 〈reg〉 are unaltered．） What is a logical OR？It is simple． It says，for each bit in 〈reg＞and the corresponding bit in the Accumulam tor，if either one is a 1 ，then the result is a 1 for that bit．Otherwise （i．e．，if they＇re both 0 ），then the result is 0 for that bit．Some examples are given in Figure 1．Notice that the Accumulator may be logically ORed with itself．The result will be it－ self．Why bother？Ah！You see，the ORA＜reg＞instruction affects the carry，sign，parity and zero flags， and it is the zero flag which I am interested in here．If the Accumulator has all zeros，then ORA A sets the zero flag．Otherwise，the zero flag is not set．That is，we now have a nondestructive test to see if the Accumulator is zero．Since it is a one byte instruction，it can replace the more cumbersome CPI 0 instruction which was used in the original get a character routine：
\begin{tabular}{lll} 
ORA & \(A\) & \(; A=\emptyset ?\) \\
JZ & - & ；Yes． \\
－ & & ；NO． \\
－ & &
\end{tabular}

This is really no big deal，of course． We save one byte of space and negligi－ ble execution time．But \(I\) find this way of doing things to be more elegant and actually easier to read and under－ stand．（There is also an ORI＜num＞－－OR immediate－instruction which logical－ ly ORs the byte of immediate data with
the contents of the Accumulator.)
STA and LDA
A more important refinement to the get a character routine is possible by understanding alternative methods of moving data to and from the Accumulator. Instead of accessing a memory location indirectly via the address in \(H L\), we can get to a memory location directly with the instructions LDA <addr>, which loads (ise., copies) the contents of <addr> into the Accumulator, and STA <adr>, which stores (i.e., copies) the contents of
the Accumulator into memory location <addr>. The get a character routine will now look like Listing 3, assuming KBCHAR EQU 81 FEH .

\section*{Exclusive OR}

There's one more new instruction to investigate: the logical exclusive OR. The XRA <reg> instruction takes the byte in <reg> and exclusive ORs it with the byte in the Accumulator. (The contents of 〈reg> are unaltered.) What is an exclusive OR? It says for each bit in 〈reg> and the corresponding bit in the Accumulator, if exactly



Figure 2. Examples of XRA <reg>.
\begin{tabular}{llll} 
GETCHA: & MVI & A, \(\emptyset\) & ; \(=\emptyset\) \\
LOOP: & STA & KBCHAR & ; Contents of KBCHAR= \(\emptyset\). \\
& LDA & KBCHAR & ;Get byte at KBCHAR. \\
& ORA & A & ; = ? \\
& JZ & LOOP & ;Yes. Wait for a key.
\end{tabular}

Listing 3. The modified get a character routine.
one of them is 1 , then that bit of the result is 1 . Otherwise (i.e., if they're both 1 or both 0 ) the result is 0 for that bit. Examples are given in Figure 2. Notice that the Accumulator may be logically XRAed with itself. The result will necessarily be zero. Like the ORI <num> instruction, there is an XRI <nam> instruction; it exclusively ORs the byte of immediate data with the contents of the Accumulator. And, like the ORA <reg> and ORI <num> instructions, the exclusive OR instructions affect the carry, sign, parity and zero flags. (Unlike the inclusive \(O R\) instructions, the exclusive \(O R\) instructions also affect the auxiliary carry flag.) often the XRA A instruction is used instead of MVI \(A, 0\), simply as an elegant way of setting the Accumulator to zero. (But remember that it will affect the flags, whereas MVI A,0 will not.) Now we can write the get a character routine as in Listing 4.

Well, that routine is fine for getting a single character. But what about getting a whole string of characters? Let's write a routine which will act something like BASIC's INPUT statement. First, however, we must make yet another digression.

\section*{CHRTNT}

As I mentioned before, there is a keyboard scanning routine in ROM which is automatically executed a nurber of times every second. If a ke is pressed, its code is put into 81.EH, the character is sent to the screen (echoed), and then the routine returns to your program. It is the part where the character is echoed on the screen that we want to interrupt so as to be able to disallow certain keystrokes. If, for example, the AUTO key is pressed, we might want to be able to ignore it. Or perhaps we want to ignore the down arrow or erase page
key, or.... A popular routine to do just this may be found in many assembly language programs for ISC machines. The location 81 FFH holds a byte which the keyboard scanning routine looks at in order to see whether it should echo a character which has just come in frow the keyboard. Now, if we could allow a character to come in from the keyboard and be stored in 81 FEH but fool the computer at just the right time into thinking that a character hasn't come in, then we will effectively have cut off the echoing of characters. It so happens that location \(81 F F H\) is used as a keyboard character flag:if the contents of \(81 \mathrm{FFH}=0\), then the character won't be echoed on the screen. Somehow we will have to add these instructions at the appropriate spot:

\section*{XRA A STA 8IFFH}

Fortunately, there is a "jump vector" at location 81 DFH (33247 decimal). Ordinarily the contents of 81 DFH are such that the keyboard scanning routine looks at that memory location and goes off to echo the character on the screen. We can interrupt that process at this point by putting the number 1 FH ( 31 decimal) into this spot. Now the keyboard scanning routine, when it finds 1 FH , will JMP to location \(8165 H\) instead of its usual spot. What is at location 81 C 5 H ? We' 11 put still another JMP instruction there so that the routine will jmup to our own routine, namely, XRA A and STA 81 FFH.

All this must sound a bit Rube Goldberg, but it is a versatile scheme which allows us a good deal of control over what happens. It is important to remember, however, that all this is occuring many times per second--interrupting our regular piogram in order
\begin{tabular}{|c|c|c|c|}
\hline \multirow[t]{2}{*}{GFTCHA :} & XRA & A & ; \(A=0\). \\
\hline & STA & KBCHAR & : Contents of KBCHAR= \(\emptyset\). \\
\hline \multirow[t]{4}{*}{LOOP:} & LDA & KBCHAR & ; Get contents of kBCHAR. \\
\hline & ORA & A & ; Wait until \\
\hline & JZ & LOOP & ; a key is pressed. \\
\hline & RET & & : Return with char in Ac \\
\hline
\end{tabular}

Listing 4. Thrice modified get a character routine.
to do so．So we must be careful that our interrupt routine does not alter the flags or the contents of any register．We ought to make sure to save and later restore anything which will be affected．Hence，the CHaRacter INTerrupt routine，since it uses the XRA A instruction，must first save the Accumulator and flags．（See Listm ing 5．）

But first，how do we get the keyboard interrupt routine to jump to CHRINT instead of to its own echoing routine？Put the jump vector number 1FH into location 81 DFH ．That＇s easy：
\[
\begin{array}{ll}
\text { MVI } & \text { A,lFH } \\
\text { STA } & 81 \mathrm{DFH}
\end{array}
\]

Now put a JMP to the CHRINT routine into locations \(81 \mathrm{C} 5 \mathrm{H}-81 \mathrm{C} 7 \mathrm{H}\) ．That is， we want：
\begin{tabular}{lc}
81 C 5 H & JMP \\
81 C 6 H & \(?\) \\
81 C 7 H & \(?\)
\end{tabular}

The JMP is easy．The 8080 JMP instruc－ tion in 0C3H（195 decimal）．So：
\[
\begin{array}{ll}
\text { MVI } & \text { A, } \varnothing C 3 H \\
\text { STA } & 81 \mathrm{C} 5 \mathrm{H}
\end{array} \quad ;=\text { 'JMP'. }
\]

Now we need only put the low byte of CHRINT into 81 C 6 H and the high byte into 81 C 7 H ．But what are those bytes？ That is，what is the address of CHRINT？ If it begins at，say， 9000 H ，then the low byte is \(00 H\) and the high byte is 90H．But is CHRINT really to be found at 9000 H ？Not necessarily．It is wherever the assembles purs it．Wo could force the assembiex to put it there with an org statement：of course：
\begin{tabular}{lll} 
& ORG & \(9 \emptyset \emptyset \emptyset H\) \\
CHRINT： & PUSH & PSW \\
XRA & A \\
STA & \(81 F F H\) \\
POP & PSW \\
RET & \\
& ORG & \(82 \emptyset \emptyset H\) \\
& \(\cdot\) & \\
& •Your program．
\end{tabular}

But we run a terrible risk：perhaps the program at 8200 H will be long enough to overwrite 9000 H and beyond， thus destroying our CHRINT routine． We could always put CHRINT way，way up in high memory，just to make sure．．．．But there＇s a far easier wy to handle all of this．

What will LXI H，CHRINT do？It wil load the HL pair with the address of the label CHRINT，whatever that hap－ pens to be．We don＇t have to know the address；the assembler figures that out for us，just as it also figures out where to jump to with something like JMP LOOP．Now that we have CHRINT＇s address in \(H L\) ，we need to put it at \(81 \mathrm{C} 6 \mathrm{H}-81 \mathrm{C} 7 \mathrm{H}\) ．We could do it as in Listing 6．But there＇s a more elegant way．

\section*{SHLD and LHIDD}

The contents of the \(H L\) pair can be copied into memory with the SHLD ＜addr＞instruction．（Store HL Direct．） The contents of \(L\) are stored at 〈addr＞， and then the contense of \(\mathrm{l}^{2}\) are stored at 〈addr＞＋1．Similathy，LRLD 〈addr〉 （Load hy Direct）will Load I rith the contente of＜addr＞and then hoad H with the contents of＜addr＞fl．

Just what we need We have the address of CHRTNT in \(H L\) ，and we know where we want to put that address：

KBFLAG EQU BIFPH ：Location of heqboard char flag．
CHRINT：PUSH PSW ：Save Acc．and flags．
XRA A ：ACC．\(=\varnothing\) ．
STA KBFLAG ；Pretend that no key was pressed．
POP PSW ：Restore Acc．and flags．
RET
Listing 5．The character interrupt routine．

LXI H,CHRINT
SHLD 8lC6H
We can now write the routine which sets up the jump to the CHRINT routine. (Listing 7.)

Don't get confused between the LXI \(H\) and LHLD instructions. LXI H, <addr> will put <addr> into HL. LHLD <addr> will load the contents of <addr> and <addr>+1 into HL. Some examples are given in Figure 3.

\section*{The Input Routine}

Using CHRINT is easy. At the beginning of your program, set up the jump vector 1 FH and the jump to CHRINT. Once that has been done, CHRINT will have been spliced into the normal read-keyboard-echo-to-screen routine. Nothing will appear on the screen
unless we explicitly put it there. CHRINT, along with something like the get a character routine, form the basis of our input routine.

Just what do we want the input routine to do? Let's specify that it is to recognize only uppercase ASCII characters, from " " through "]". A backspace (left arrow) will cause the previous character to be erased, and a carriage return will end the input. Let's also specify that no more than 128 characters will be allowed.

\section*{Next time}

Try your hand at writing a full fledged assembly language program for the input routine. I'll show you my version in the next issue.
\begin{tabular}{lll} 
LXI & H,CHRINT & ;Get address of CHRINT. \\
MOV A,L & ;Low byte in ACC. \\
STA 81 C 6 H & \\
MOV A, H & ;High byte in ACC. \\
STA 81C7H &
\end{tabular}

Listing 6. Setting up the jump to CHRINT. First version.


\section*{Before LXI H,900@日}
\(\begin{array}{lll}\frac{\text { Address }}{900 \emptyset H} & \frac{\text { Contents }}{9 A H} \\ 9001 \mathrm{H} & 3 \mathrm{HH} & 7 \mathrm{HL} \\ & & \end{array}\)
Before LHLAD 9900H
\(\begin{array}{cc}\frac{\text { Address }}{9 \emptyset 0 \emptyset \mathrm{H}} & \frac{\text { Contents }}{9 \mathrm{AH}} \\ \begin{array}{cc}9 \emptyset \emptyset 1 \mathrm{H} & 3 \mathrm{DH}\end{array} & \\ 73 \mathrm{AL} 4 \mathrm{H}\end{array}\)
Before SHLD 9000H
\(\begin{array}{cc}\frac{\text { Address }}{9000 \mathrm{H}} \\ 9001 \mathrm{H} & \frac{\text { Contents }}{\emptyset \mathrm{AH}} \\ \begin{array}{ll}3 \mathrm{DH}\end{array} & \\ 73 \mathrm{HL}\end{array}\)

\section*{After LXI H,9000H}
\(\begin{array}{ll}\frac{\text { Address }}{9 \emptyset \emptyset \emptyset \mathrm{H}} \\ \begin{array}{c}\text { Contents } \\ 9 \emptyset \emptyset l \mathrm{H}\end{array} & \begin{array}{l}\text { 3DH } \\ \\ 3 \mathrm{DH}\end{array}\end{array}\)

\section*{After LBLD 900@H}


\section*{After SHLD 900日H}
\(\begin{array}{cc}\frac{\text { Address }}{9000 \mathrm{H}} & \frac{\text { Contents }}{9 A 4 \mathrm{H}} \\ 9001 \mathrm{H} & 73 \mathrm{H}\end{array} \quad \begin{aligned} & \mathrm{HL} \\ & \end{aligned}\)

Figure 3. Examples of the LXI H, LHLD and SHLD instructions.

CHECK THESE EXAMPLES OF HIGH QUALITY SOFTWARE LISTED IN OUR CATALOG!

CTE SCREEN EDITOR

Order number: CT833
Price: US\$ 40.00

DEBUG

Order number: CT841
PRICE: US\$ 30.00

TERMII

Order number: CT821
Price: US\$ 70.00

The new editor is a high speed machine lanquage editor progr.in which will move block, copy, delete or print text. CTE will change lower cas e characters to upper case yellow characters for those who do not have lower case installed. CTE will allow the insertion of cyan colored control characters (these characters are not counted as text, so printing will be normal) in the text for printer use. CTE has a global search-and-replace function on all ASCII characters 0-127 in both forward and reverse. CTE has an active character, line page counter on the screen at all times and is updated on every key press or use of home key. CTE has all single key commands plus any keyboard can be used. CTE will allow you to set your printer's baud rate and print all the text or marked text only. CTE has a fast typematic keyboard action. CTE will set up it's text buffer to the memory size of your computer or honor the memory size auestion from BASIC at 8OAC hex.
CTE is an excellent low cost wordprocessor
DEBUG is a machine language EDITOR, DISASSEMBLER, HEX AND ASII DUMP, MASTER DEBUGGER, and more. Use any serial printer (with user selected print format), select Baud rates, print ASII/HEX dumps (in compressed or standard mode) to the printer \& CRT with a floating address header for every 60 lines (printerlor 24 lines (CRT), move, copy, compare, verify memory blocks, disassemble programs, test all memory locations, search for a string of data (1-16) bytes long etc.

TERMI is the most powerful communications program for the CCII and Intecolor. It incluces: up/down loading of disk files (all types, with byte for byte verification): dumb terminal communications up to 9600 baud; send/receive and buffer all types of terminal outputs (binary, text, random files, pictures, etc.) at 110 to 4800 baud, print/save iil buffered data to the disk, display all in-coming \& out-going data to the CRT in HEX and/or ASCII.
You can also set upper/lower case character colors. Half/full duplex, parity ODD/EVEN/NONE, turn on/Off ECHO BACK feature, Xon/Xoff protocol, and special menu HELP command. All commands can be toggled on/off. All keyboard type-outs are in one color and port inputs in another color (both selectable). Select large or small character mode, and many more features too numerous to mention.
in Australia: For Intecolor: Color Computer Systems, ATTN: Don Sforcina 58 Valley Rosd, Hornsby, NSW 2077, Phone: 02-476-2480

INVOICE/
ORDER

Order number: EM111
Price: US\$ 38.00

\begin{abstract}
The INVOICE/ORDER program prints automatically the name and address of your company together with the personalized invoice, stating the item numer, description, number of units and the total price. This program runs together with any kind of printer, even using Epson's ability for print densities. The program is easy to operate and wrong entries can be corrected without difficulties. It allows entry of 10 lines comments, special discounts, shipping charges, taxes, etc. It is asking for the mode of payment, even for the credit card number. The totals (with customer name and address) can be stored in a data base which will be generated automatically.
The ORDER program has the same features as the invoice program, and has in addition of the data base with the order totals a data base for frequently used vendors and their addresses.
\end{abstract}

This is the first graphics checkbook on the market. You can specify your own 15 expenses and 6 deposit categories. The program will print tables, sort according to the date, display/print totals of one or more specifies categories, and generate a graphic display or printout of the monthly money flow, or the totals for each category within a given year or month. This is an excellent tool, when tax time is coming.

This program generates individual personalized labels of your family, your own company, or organizations which you have contact freqently. It can be used with the IDS "Paper Tiger" printer and all Epson printers, or printers which are able to print in different print densities. You can store the labels in different categories. This way, you can print the same label again later at any time.

This program allows you to create 7 PERT plans on one \(51 / 4\) disk with 200 activities each. You may create, change, modify, and analyze your PERT plans to detect the critical path, enter the actual dates, and calculate the ending date within the given uncertainty range. This program is a "must" for program and production managers.
Included with the PERT program are two programs which allow you to calculate the end date from a beginning date by adding a certain number of working days, and to calculate the number of working days between any two given dates.

The perfect MODEM for your communications:
HAYES SMART MODEM
fully automatic, programmable
US \$ 235.00
The perfect DISKETTE for your programs;
VERBATIM DISKETTES 10 blank diskettes US \(\$ 24.00\)
10 formatted diskettes US\$ 29.00
The perfect low-cost PRINTER for your paperwork;
C.Itoh 8300

Dot Matrix 132 columns
125 CPS, bi-directional
US \$ 375.00

\section*{Colorcue}

\title{
Editorial Offices
}

161 Brookside Dr.
Rochester, NY 14618
Address Correction Requested

\section*{Colorcue}

RUEIK＇S
CUBE

E1
ENU1F3U3
E1ERG1E1
3ロジロきた1ロふたる
E1口． 3


E1ח3E3GBELBE1
\(\square 1\)




FEU1CBFEUSL1


RUBIK＇s
CUBE
\(5=\) IHETRUGTIDHS
\(3=\) MHTCH SLAOLWE SELECT

EUEEY
OTHE
MAGAZISE
ME WE SEEH HAS
HACI NEIE S ODEE
OH ITS COWE
AHC LAE THOUIEHT
WE SHOULI：TVII．

\section*{Colorcue}

\section*{A Bi-monthly Publication by and for Intecolor and Compucolor Osers}

\section*{Editors:}

Ben Barlow
David B. Suits
Compuserve: 70045,1062

June/July, 1982
Volume 4, Number 6

\section*{3 Editors' Notes}

4 New Products from ISC
5 Rubik's Cube Demystified, by Roger Safford Have the computer solve it

9 Assembly Language Screen Dump to MX80 Printer by Steve Reddoch
A fast version of a program originally in BASIC
13 CALL Subroutine Linkage, by Ben Barlow Getting two languages together

17 Keyboard Reading in BASIC, by Steve Perrigo The INs and OUTS of keyboard scanning

20 Assembly Language Programming, by David B. Suits Part VI: The Input Routine

25 Classified Colorcue Index, by James A. Kavanagh A comprehensive CCI code

27 About Your Subscription

Advertisers: You will find our advertising policies attractive. Write for details.

Authors: This is a user-oriented and supported publication. Your articles/tips/hints are required to make it go. Send your articles or write for information.

Colorcue is published bi-monthly by Intelligent Systems Corporation, with editorial offices in Rochester, New York. Editorial and subscription correspondence should be addressed to the Editors, Colorcue, 161 Brookside Dr., Rochester, NY 14618. Product related correspondence should be addressed to ISC, 225 Technology Park, Norcross, GA 30092, ATTN: Susan Sheridan. Opinions expressed in by-line articles are not necessarily those of the editors or of ISC. Hardware/software items are checked to the best of our abilities but are NOT guaranteed.

\section*{Editors' \\ Notes}

As we mentioned in the last issue, this issue is the last issue subsidized by ISC. From this point on, Colorcue's continued existence will depend on you. We need a certain number of renewals to make its publication worthwhile-mot to make a huge profit, but to keep the enterprise from becoming a public service. Subscriptions have begun to roll in on a regular basis. We thank those early reuppers (we'll hold your checks until we reach our magic number); your response has been gratifying. The suggestions for articles that we've received along with the renewals will be an excellent source for planning next year's content. Keep them coming.

Compucolor users interested in starting a user group in the northern New Jersey area should contact Amod Patwardhan, 21 Beachmont Terr., N. Caldwel1, NJ 07006 (201) 226-8619.

The CUWEST user group has a fully working Scrabble game available for the cost of a disk and airmail postage. The game, written by Chris Teo, holds a dictionary of over 700 two and three letter words. The computer's moves take only about a minute. Send a check for \(\$ 9\) payable to CUWEST to Doug Crant, CUWEST librarian, 2 Brookside Ave., South Perth, Western Australia 6151.

Doug Van Putte, who wrote the \(3-D\) Graphics article in the Feb/Mar issue, suggests these changes to his program:
```

285 FOR L=1 TO 3:A=A(L)*3.14159/180:
IF A(L)=0 THEN 295
290 GOSUB 560: GOSUB 450
295 NEXT L
550 delete this line
500 RETURN

```

A new version of 'THE' BASIC Editor has been released by Quality Software Associates, 21 Dersingham Crescent, Thornhill, Ontario, Canada L3T 4P5. (See original review in the Dec/Jan issue of Color-
cue). The new version features typamatic keys, improved cursor movement, and many other improvements. If ordered in EPROM or for the Devlin RAM board (see Apr/May Colorcue), 'THE' Editor will also contain a full set of FCS commands, an erase line command, an origin command to move a BASIC program, and a print command for programs and variables. A super program made even better!

We'd like to have some articles on languages. We have been using BASIC and assembly language for years. But what of the less well-known languages available? A tiny-c interpreter is available; so is FORTH and, of course, FORTRAN, about which we will have a series of articles in the coming months. Whatever happened to PILOT? Do you remember rumors of a compiled language from ISC called ALGAE? It was an interesting approach to a block structured, procedure-oriented language, but apparently it was never officially released. Anyone for LISP? It is available on a few microcomputers, but is anyone running LISP on an ISC machine? (One of your anxious editors is drooling at the thought.) There are rumors of a BASIC compiler (actually, two of them) in the works. One in Germany and the other in California. Will their authors and/or users please stand up? By the way, there's a money maker! Suppose you wrote a BASIC compiler, even a simplified, stripped down version. Double precision integers instead of floating point numbers. Maybe omit trig functions and random file handling. You'd still have a very marketable product. How much would you sell it for? Let's make it. cheap, so as to encourage lots of sales. \(\$ 50\) ? \({ }^{8} 11\) take one for that price. (At \(\$ 75\) you'd be testing the patience of my poor bank account, and at \(\$ 100 \mathrm{I}\) would, reluctantly, walk away.) Now, how many can you sell? That's difficult to estimate, but at \(\$ 50\) surely you'd sell at least five or six hundred of them. Let's see... \(\$ 50 \times 600=\$ 30,000\). There's a tidy sum! We have had thoughts of just such a project for some time now, but it has always been pushed to the back burner. Perhaps, if no one else produces one (after all, a compiler is a difficult beast to create), we'll get back to work on it. Consider this a challenge. First one there wins.

\section*{NEW PRODUCTS FROH ISC}

\section*{TERMILAL SOFTWARE}

ISC has released FlexWare, a terminal firmware package for ISC 8000, 8300 and 8900 machines. FlexWare's three asynchronous communications protocol modes with parameters definable from the keyboard or host allow ISC terminals to be interfaced to a variety of hosts. FlexWare's editing functions are accessed through ANSI sequences, and color via ISO standard color sequences. FlexWare's built-in expansion capabilities allow the user to interface his own firmware extensions and special functions by simply adding EPROMs.

Optional function keys can consist of keystrokes in local or on-line modes. Function keys can also be loaded with executable machine language patches. FlexWare has a disk \(I / 0\) option that allows simultaneous capture of incoming data. The CRT display may be subdivided into as many as 32 discrete regions, each with its own roll mode, colors and character size. Single prece price is \(\$ 250\).

\section*{GRAPHICS SOFTWARE}

ISC has released IGS (Intelligent Graphics System), an advanced color graphics programming language with over 100 grapnics commands which can be issued by any host computer. IGS may be used alone or driven by other computer languages. Features include seven type fonts, automatic labelling, scaling and rotation of grids and axes, rotate and zoom, and real-time plotting. IGS runs on ISC 8000R and 8300 R series terminals and computers. Introductory price for IGS and 8301 R terminal is \$3490.

\section*{TEKTRONIX 4014 EMULATOR}

ICS's 8001R terminal with 8 color dot-addressable display is available with Tekcronix 4014 emulation package. The ISC terminal emulates the Tektronix storage tube display by using the "shrink mode", in which the terminal receives vectors and fits them to its own 480 x 384 resolution. Color is added to the data by using the intensity bit in the 4014 data string. The emulation package also supports the screen dump to Printronix printers and Printacolor 8-color graphics printers. The 100 piece price is \(\$ 3275\).

\section*{DEC VT52/VT100 EMULATOR}

This software package allows the ISC terminal to behave like VT52 or VT100 in DEC environments. Color coding is used to emulate the VTl00's reverse video, underscore and bold. 100 piece price, including terminal, is \(\$ 2465\).

\section*{IBM 3275 EMULATOR}

ISC's 8001 terminal with \(13^{\prime \prime}\) or \(19^{\prime \prime}\) display allows for an 80 character by 24 line display with a 25 th line for status. Soitware is now available so that the terminal may be attached to a host with binary synchronous communication (BSC). Color control is provided by the 3270 data stream. 100 piece price for terminal and software is \(\$ 4055\).

\section*{UL-APPROVED TERMIEAL}

ISC has received Underwriters' Laboratory approval for a new version of its model 8001G \(19^{\prime \prime}\) color terminal. The new terminal, the \(8001 \mathrm{G} / 82\), meets UL 478 (electronic data processing) and FCC Class A regulations. While the digital circuitry of the \(8001 \mathrm{G} / 82\) remains unchanged from the original \(8001 G\), the analog circuitry has been completely redesigned, including a new switching power supply and switchselectable 115 or 220 input voltage. 100 Piece price is \(\$ 2870\).

\section*{NEM TERMILAL}

The IS/2405 has an 80 column, 24 line color alphanumeric and graphics display with a single-board design based on the 8085 microprocessor. There are 12 function keys providing 36 programmable functions, each of which can store up to 40 characters. The terminal has a resident set-up mode for selecting physical I/O parameters which is retained by 2 K of CMOS RAM even during power off. The IS/2405 also has 4 K of screen RAM and 8 K of EPROM for the operating system and an RS-232C printer port and a communications port that can be configured for RS-232C or current loop interface at 19.2 K baud. The standard single-piece price is \(\$ 1995\), but before October I, 1982, the price is \(\$ 1200\).

For information on any of the above products, contact Marketing Services, Intelligent Systems Corporation, 225 Technology Park, Norcross, GA 30092, (404) 449-5961.

\title{
Rubik's Cube Demystified
}

\author{
by Roger Safford \\ 15 Safford Ave. \\ Perry, NY l453ø
}

The Rubik's Cube fad may be on the down cycle, but fortunately programming for fun is not. The program listing which follows is not just for the eight year old down the street who has no problem solving the cube in 37.5 seconds. He assuredly will find it a challenge trying to match up the colors with the handicap of not being able to tumble the Rubik's Cube in his hands. This program is also for those who don't have an eight year old to show them how to solve the Cube, and for those who followed the array of patterns described in the book but ended up with a colorful mess not unlike the
one they started with. Unlike the book, however, this is not a set of patterns. You enter the color orientation matching that of the Cube in your hands, and the computer will print a list of moves which will solve the cube. If there is no solution (if someone has moved the stickers or disassembled it), the program will let you know. The program will also let you try to solve it yourself by entering single or multiple move commands.

The program RUBIKS uses approximately 13 K . Have fun. [Yes, and have fun typing it in! We suggest a team of eight year olds be hired for the job. --Eds.]

\section*{}



310 DATA \(6,39,4,18,40,5,42,38,3,54,41\) ， 2
315 DATA \(1,14,53,13,17,52,49,16,51,37,15,50\)
318
REM \(\qquad\) EDGIE IDENT．DATA
320 DATA \(8,47,9,44,101,45,11,46,36,23,12,22\)
325 DATA \(24,21,48,24,7,35,31,34,43,33,19,32\)
325
328
329
329 REM
REM－\(\quad\)－－ \(\qquad\)
\(349 \mathrm{D} \$=\mathrm{D} \$+" \mathrm{DE}\)
345 \(D \$=D \$+F \$+" 3 D 1 "+F \$+" 1 ":\) RETURN
\(350 \mathrm{D}=\mathrm{F}=\mathrm{F}+\)＂ \(3 \mathrm{D} 3 "+\mathrm{F} \$+\)＂1＂
355 D \(\ddagger=D \pm+\)＂D 2 ＂
\(360 \mathrm{D} \$=\mathrm{D} \$+\mathrm{I} \$+\)＋1D3＂＋I\＄＋＂3＂：RETURN
365 D \(5=" D 3 "\)
370． \(\mathrm{D} \$=\mathrm{D} \phi+\mathrm{I} \$+{ }^{\prime \prime}{ }^{\prime \prime}\)
375 D\＄＝D\＄＋＂U1＂＋H\＄＋＂3U3＂：RETURN
380 D \(\$=" D 3^{\prime \prime}\)
\(385 \mathrm{D} \$=\mathrm{D} \$+\mathrm{I} \$+" \mathrm{E}\)＂：RETURN
390 D \(=\)＝＂D1＂：GOTD 360
395 D ＝＂D3＂：GOTO 360
\(409 \mathrm{D} \$=F \$+" 3^{\prime \prime}\)

\(410 \mathrm{D} \Phi=\)＂U3＂＋F\＄＋＂3U1＂：RETURN
\(415 \mathrm{D}=\mathrm{G} \$+{ }^{(1)}\)
420 D\＄＝D\＄＋＂U2＂＋G\＄＋＂3Uこ＂：RETURN
425 D \(=\)＝＂D1＂：GOTO 385

4．

\(45 \mathrm{D} \$=1 \$+" 3 \mathrm{D} \mathrm{D}^{2}+1 \$+\cdots\)
\(455 \mathrm{D} \$=F \Phi+{ }^{2} 3^{\prime \prime}\)

\(465 \mathrm{D} \phi=" \mathrm{DE}\)＂：GOTO 370
465 D \(==\) DE＂：GOTO 370
470 D \(==\) D1＂：GOTO 370
\(475 \mathrm{D}=\)＝＂DE＂：GOTD 385
\(480 \mathrm{D} \$=\mathrm{H} \$+111\) ：GOTO 375
\(485 \mathrm{D}=\mathrm{G} \$\)＋＂\(^{\prime \prime}{ }^{\prime \prime}\)
\(490 \mathrm{D} \$=\mathrm{D} \$+\mathrm{HU} 1 \mathrm{H}+\mathrm{H} \$+\)＂1U3＂：RETURN
\(495 \mathrm{D}=\mathrm{H}=\mathrm{H}+" 1 "\)

\(505 \mathrm{D} \$=\mathrm{I} \$+" 3 ":\) GOTO 375
S10 D\＄＝G\＄＋＂3Dご＋Gゅ＋＂1＂：GOTO 345
515 Dक＝＂D1＂：GOTO 345
520 Dぁ＝＂D3＂：GOTO 345
525 D \(\$=\)＂D3＂：GOTO 335
53＠Dq＝＂Uこ＂＋Gゅ＋＂1しご
535 RETURN
540 D\＄＝G\＄＋＂1D1＂＋G\＄＋＂3＂：GOTO 355
\(545 \mathrm{D} \$=F \$+" 1 D \mathrm{~S}^{2}+\mathrm{F} \ddagger+" 3 ":\) GOTO 345
\(550 \mathrm{D} \Phi=\mathrm{G} \$+" 1 \mathrm{D} 1 "+6 \$+" 3 ":\) GOTO 345
\(555 D \$=H \$+" 1 D D^{\prime \prime}+H \$+" 3 ": G O T O\)
\(560 \mathrm{D} \$=1 \$+" 101 "+1 \$+" 3 ": G 0 T 0 ~ 340\)
\(560 \mathrm{D} \$=1 \$+" 1 D 1 "+1 \$+" 3^{\prime \prime}: G 0 T 0340\)
\(565 \mathrm{D}=\mathrm{DH}:\) ON CM GOTO 535， \(390,3 \mathrm{EQ}, 395,355,330,535,405,410,420\) 570 ON CM－10 GOTO \(375,425,430,350,435,440,445,450,455,365,455\)
 580 ON CM－ 2 C GOTO \(485,495,505,385,510,340,515,345,520,525,415\) \(590 \mathrm{D}=\cdots ": 0 \mathrm{NCM}\) GOTD 430 ． \(355,390,360,3 \exists 5,525,480,375,405,410\)


600 ON CM－E1 GOTO \(4 E 5,470,390,535,535,535,535,535,535,415,505\) 605 ON CM－32 GOTO 400，485，495，455，535，520，340，515，345，335，455
 \(615 \mathrm{D}=4 \cdot 4: 0 \mathrm{~N} \mathrm{CM}\) GOTO \(540,395,355,390,360,335,415,420,375,405\) \(615 \mathrm{D} \Phi=\cdots \cdots: O N\) CM GOTO \(540,395,355,390,3 E 0,335,415,420,375,405\)
620 ON CM－10 GOTO \(410,380,519,440,445,350,435,525,480,470,370\) 625 ON CM－21 GOTO \(365,465,395,535,535,535,535,535,535,455,495\) 625 ON CM－21 GOTO \(365,465,395,535,535,535,535,535,535,455,495\) 630 ON CM－32 GOTO \(505,4001,485,475,430,345,5 \approx 0,340,515,450,535\) 635 ON CM－43 GOTO 5010，4E0， \(530,490,4 E 5,535,555,5 E 0,545,550,330\) 640 Dq＝＂M：ON CM GOTO \(510,360,395,355,390,450,455,410,420,375\)
\(645 \mathrm{ON} \mathrm{CM}-10 \mathrm{GOTO} 405,385,535,435,440,445,350,355,415,465,470\) 645 ON CM－18 GOTO \(405,385,535,435,442,445,350,335,415,465,470\)
650 N CM－21 GOTO \(37 \pi, 365.435,535,535,535,535,535,535,535,485\) \(655 \mathrm{ON} \mathrm{CM}-32 \mathrm{GOTO} 495,505,400,380,540,515,345,550,340,330,490\) 560 ON CM－43 GOTO 490，500，4E0，530，475，430，550，555，560，545，535 \(665 \mathrm{D}={ }^{6}=\mathrm{M}: 0 \mathrm{~N} 5 \mathrm{GOTO} 690,690,695,685,695,695,695,695\)

675 UN S－17 GOTO 685， \(685,695,695,695,685,690,695,680\)
680 F \(\$=\)＂E＂：Gs＝＂＇－＂：GOTO 700
685 F \(==" L ": G \$=" F ": G 0 T 07\) 7el
690 F \(=\)＂F＂：G\＄＝＂R＂：EDTO 7 700



710 RETURN

\(720 \mathrm{D} \phi=\because \mathrm{H}: \mathrm{R}=\mathrm{E}: \mathrm{ON}\) CM－7 GOTO \(765,770,775,780,760\)
\(725 \mathrm{R}=1\) ：ON CM－19 GOTO 750，755，760，745
\(730 \mathrm{R}=2: I F \quad \mathrm{CM}=24 \mathrm{GOTO} 75\)
735 IF CM＝36GOTO
790，795，800， 750
\(745 \mathrm{D} \$=" L 1 U 1 D 3 E 1 U 1 D 3 R 3 U E D E L 3 ": O N\) R GOTO 775，790
 760 D＝＂耳1





\(790 \mathrm{D} \Phi=\mathrm{D} \$+\mathrm{J} \$+" \mathrm{~F}\) 1U3D1EEUSD1LRUこDEF3＂＋K\＄：RETURN


805 ON R－1 GOTD \(310,815,820,895,830,835,840,845,850,855,860\) \(810 \mathrm{D}=\)＝＂REUSDIEEUIDコ＂：RETURN
815 D \＄＝＂BEUID3REU3D1＂：RETURN
8こせ D\＄＝＂UこDEFこUこDこEこ＂：RETURN
825 D\＄＝＂FeU3D1R2U1D3＂：RETURN
83＠D\＄＝＂LE゙U3D1FEU1D3＂：RETURN
835 D \(\$=" R E U 103 F E U 3 D 1 ":\) RETURN
84e1 D\＄＝＂REU3D1LEU1D3＂：RETURN
845 D\＄＝＂RELSU1DSFEERU3D1＂：RETURN
850 D \(\$=\)＂F2U1D3LEU3D1＂：RETURN
855 D \(\ddagger=\)＂LeU1D3EEUSD1＂：RETURN
860 D\＄＝＂RミUきDミLこUこDき＂：RETURN
865 D\＄＝＂R3L1U1R3L1F1RSL1D1RELEU1R3L1F1R3L1D1R3L1B2＂：RETURN 870 D \(\ddagger=" R 1 L 1 U 1 R 3 L 1 F 1 R 3 L 1 D 1 R E L 2 U 1 R 3 L 1 F 1 R 3 L 1 D 1 R 3 L 1 R E R 2 ": R E T U R N\) 875 D \(\$=\)＂ESF1DLEJF1LIBJF1U1BEFED1BJF1LIEJF1U1BJF1RE＂：RETURN \(880 \mathrm{D} \$=\)＂R3F1U1E3F1R1E3F1D1E2FこU1B3F1R1E3F1D1E3F1LE＂：RETURN 895 Ds＝＂R3L 3 U1R3L1F1R3L1D1RELZU1R3L1F1R3L1D1R3L1B2L2＂：RETURN 890 Dt \(=\)＂＂
\(895 \mathrm{D} \$=\mathrm{D} \$+" \mathrm{R} 3 \mathrm{~L} 1 \mathrm{D} 1\) R3L1B1R3L1U1R2L2D1R3L1B1＂
909 D \(\$=D \$+" R J L 1 U 1 R 3 L 1 F こ ゙: R E T U R N\)
905 G0SUB 865：PLOT 3，0，WW，6，7：WW＝WW＋1：PRINT D\＄：GOTO 895
\begin{tabular}{|c|c|}
\hline \multicolumn{2}{|l|}{\multirow[t]{2}{*}{}} \\
\hline & \\
\hline \multicolumn{2}{|l|}{10015 A \(=\) CHR （ 240 ）＋＂eeeDVVV＂＋CHR\＄（255）} \\
\hline 1010 &  \\
\hline \multicolumn{2}{|l|}{} \\
\hline \multicolumn{2}{|l|}{1020 FOR \(Z=1\) TO 47 STEP 6：C \((z)=3\)} \\
\hline \multicolumn{2}{|l|}{\(10225 C(z+1)=1: C(z+2)=7: C(z+3)=5: C(z+4)=4\)} \\
\hline 1030 & \(C(Z+5)=2:\) NEXT \(Z\) \\
\hline \multicolumn{2}{|l|}{10355 PLOT 27，24，6，6，14，12} \\
\hline \multicolumn{2}{|l|}{104\％PLOT 3，35，1：PRINT＂UP BACK LEFT＂} \\
\hline 10145 & PLOT 3，21，14：PRINT＂FRONT RIGHT DOWN＂ \\
\hline \multicolumn{2}{|l|}{1050 PLOT 3，2，2， \(6,3: P R I N T\)＂R U E I K＇S＂} \\
\hline \multicolumn{2}{|l|}{1055 PLOT 3，11，4，6，4：PRINT＂C U E E＂：GOSUB 115} \\
\hline \multicolumn{2}{|l|}{1060 PLOT 15， \(6,2,3,0,7:\) PRINT＂1－INSTRUCTIONS＂} \\
\hline 1065 & PRINT＂E－MIX \＆PLAY＂：PRINT＂3－MRTCH \＆SOLVE＂ \\
\hline \multicolumn{2}{|l|}{1070 PRINT ：INPUT＂SELECT＂；D\＄} \\
\hline \multicolumn{2}{|l|}{10175 IF Ds＝＂1＂GOTO ご囚め} \\
\hline \multicolumn{2}{|l|}{1080 IF D\＄¢＂ご AND D\＄く＂3＂GOTO 1035} \\
\hline \multicolumn{2}{|l|}{1085 PLOT 3， 0 ，16} \\
\hline \multicolumn{2}{|l|}{1090 INPUT＂DISPLAY ERCH MOVE（ Y／N ）？＂；E} \\
\hline \multicolumn{2}{|l|}{\(1095 \mathrm{E}=1: \mathrm{IF} \mathrm{E} \$=\)＂Y＂THEN E＝2} \\
\hline \multicolumn{2}{|l|}{1100 IF D \(=\) \＃＂3＂GOTO 1300} \\
\hline \multicolumn{2}{|l|}{1198} \\
\hline \multicolumn{2}{|l|}{1199 REM－－－－－－－－－－－－－－mIX AND PLAY ROUTINE} \\
\hline \multicolumn{2}{|l|}{12a¢ Input＂how many Mixes＂；U} \\
\hline \multicolumn{2}{|l|}{} \\
\hline \multicolumn{2}{|l|}{121 FOR J＝1 Tロ U： \(\mathrm{D}=\mathrm{INT}(\operatorname{RND}(8) * 6+1\) ）} \\
\hline \multicolumn{2}{|l|}{} \\
\hline \multicolumn{2}{|l|}{} \\
\hline \multicolumn{2}{|l|}{1225 NEXT J：GOSUE E45：G0SUE 115} \\
\hline \multicolumn{2}{|l|}{1230 PLOT 3， \(2,17,11,28,11\)} \\
\hline \multicolumn{2}{|l|}{1235 PLOT 3，4，1E：INPUT＂ENTER MOVE（ TYPE PQUIT＇TO END ）＂；D\＄} \\
\hline \multicolumn{2}{|l|}{\(1240 \mathrm{~F}=\mathrm{E}: \mathrm{GOSUB} 245: \mathrm{IF} F=1\) THEN GOSUB 115} \\
\hline \multicolumn{2}{|l|}{1245 GOTO 1230} \\
\hline \multicolumn{2}{|l|}{1298} \\
\hline \multicolumn{2}{|l|}{1299 REM－－－－－－－－－－－－－－－－MTCH AND SOLVE RDUTINE} \\
\hline \multicolumn{2}{|l|}{1300 GOSUE 1805：WW＝16} \\
\hline \multicolumn{2}{|l|}{1303} \\
\hline \multicolumn{2}{|l|}{1394 REM－－－－－－－－－－－－－－－－－SOLVE TOP EDGIES} \\
\hline \multicolumn{2}{|l|}{} \\
\hline \multicolumn{2}{|l|}{1310 RESTORE 320 ：EX＝C（25＋EG）＊C（25）：FOR EY＝1 TO 12：RERD R， 5} \\
\hline \multicolumn{2}{|l|}{1315 IF EX（）C（R）＊C（S）THEN NEXT EY：EY＝C（ES）：GOTO EOM} \\
\hline \multicolumn{2}{|l|}{\(1320 \mathrm{EY}=1 \mathrm{E}:\) NEXT EY：CM＝R：IF C （S）\(=\mathrm{C}(\mathrm{ES}\) ）THEN CM＝5} \\
\hline \multicolumn{2}{|l|}{1325 ON EG GOSUB 590，Є15，640，565：IF Di＝＂＂GOTD 1335} \\
\hline \multicolumn{2}{|l|}{1330 PLOT 3，D2，7＋EG，6，3：PRINT D\＄：GOSUE E45：GOSUB 115} \\
\hline \multicolumn{2}{|l|}{} \\
\hline \multicolumn{2}{|l|}{1338} \\
\hline \multicolumn{2}{|l|}{1339 REM－－－－－－－－－－－－－－－－SOLVE TOP CORNERS} \\
\hline \multicolumn{2}{|l|}{1340 FOR EG＝1 TO 4：RESTORE 310：EX＝C（ 35\() * C(26) * C(29)\)} \\
\hline \multicolumn{2}{|l|}{1345 IF EG（ 4 THEN EX＝C（25）＊C（2S＋EG）＊ \(2(26+E G)\)} \\
\hline \multicolumn{2}{|l|}{1350 FOR EY＝1 TO 8：READ \(R, S, T: Q=C(R) * C(S) * C(T)\)} \\
\hline \multicolumn{2}{|l|}{1355 IF EX（）Q THEN NEXT EY：EY＝C（25）：GOTD E®45} \\
\hline \multicolumn{2}{|l|}{1360 EY＝9： NEXT EY：CM＝T：IF C（R）\(=\mathrm{C}(25)\) THEN \(\mathrm{CM}=\) R} \\
\hline \multicolumn{2}{|l|}{1365 IF C（S）\(=\mathrm{C}\)（25）THEN CM＝S} \\
\hline \multicolumn{2}{|l|}{1370 ON EG GOSU日 590，615， 540 ，565：IF D\＄＝＂＂GOT0 1390} \\
\hline \multicolumn{2}{|l|}{1375 PLOT 3，Q，11＋EG，6，E：PRINT D\＄：GOSUE E45：GOSUE i15} \\
\hline \multicolumn{2}{|l|}{1380 IF C（EG＋25）（）C（EG＋49）THEN EY＝C（25）：G0T0 2040} \\
\hline \multicolumn{2}{|l|}{} \\
\hline \multicolumn{2}{|l|}{1388} \\
\hline \multicolumn{2}{|l|}{1389 REM－－－－－－－－－－－－－CHECK FOR VALID CORNER COLORS} \\
\hline \multicolumn{2}{|l|}{1390 FDR EG＝1 T0 4：RESTORE 310：EX＝C（30）＊C（26）＊C（29）} \\
\hline \multicolumn{2}{|l|}{1395 IF EG（ 4 THEN EX＝C（30）＊ 2 \｛2ら＋EG\}*C(26+EG)} \\
\hline 1480 & FOR EY＝1 TO 4：READ R，S，T： \(\mathrm{Q}=\mathrm{C}(\mathrm{R}) * \mathrm{C}(\mathrm{S}) * \mathrm{C}(\mathrm{T})\) \\
\hline
\end{tabular}

1405 IF EX（\}Q THEN NEXT EY:EY=C(30):GOTU EDES
410 EY＝4：NEXT EY，EG
1413 \(\qquad\) －PATTERN TO FLIP EQTTOM CDRNERS
1415 REM－
142 R＝0：RESTAR \(\quad\) R：IF \(C(R)=C(J Q)\) THEN \(S=S+3 \wedge I * J\)
425 NEXT J，I：IF \(S=0\) OOTO 1445
1430 PLOT 3 ， 0 ，Wh， \(6, E: W W=W W+1: G 05 U E\) 665：PRINT D \({ }^{1}\)
1435 IF WW）EQ THEN PLOT 3，D，E8：PRINT＂NO SOLUTION＂：GOTO EDS0
1440 GOSUE \(245: G 05 U E 115: G 0 T 01415\)
1445 FOR CM＝2 TO S：IF C \((\mathrm{CM})=\mathrm{C}(\mathrm{CM}+36)\) GOTO 146 CD
1450 NEXT CM：GOSUE 715：PLOT 3，D，WW，E，E：WW＝WW＋1：PRINT DE
1455 GOSUB 245：GOSUB 115：GOTO 1445
1460 FOR \(J=26\) TO EЭ：IF C（CM）认C（J）THEN NEXT J
\(1465 \mathrm{EG}=\mathrm{J}-\mathrm{CM}-\) 24：IF EG THEN EG＝EG＋4
1470 ON CM－E GOTO ：485，149 14,1495
1475 IF C \((4)=C(40)\) GOTO 1510
\(1480 \mathrm{~F}=\)＝＂B＂：G\＄＝＂L＂：GDTO 1500
\(1485 \mathrm{~F} \Phi=4 \mathrm{~L}\)＂：G\＄＝＂F＂：GOTO 15018
490 F ＝＂F＂：G\＄＝＂R＂：GOTD 1501
1495 F \＄＝＂R＂：G\＄＝＂B
500 GOSUE 715：PLDT 3，0，WW，E，こ：WW＝WW＋1：PRINT D\＄
1505 GOSUB 245：GOSUE 115
1510 IF EG＝VGOTO 1525

520 PRINT D\＄：GOSUB 245：GOSUE 115
1523
1524 REM－－－－－－－－－－－－－－－－－－－FINAL CORNER CUEIE CHECK
1525 FOR EG＝1 TO 4：IF C（EG＋37）＝C（25＋EG）THEN NEXT ：GOTO 1540
\(1530 \mathrm{EY}=\mathrm{C}(26+E G): I F E G=4\) THEN \(E Y=C(26)\)
\(1535 \mathrm{EX}=\mathrm{C}(30) * C(25+E G) * E Y: E Y=C(307): G O T O\) EQ20
1538
1539 REM－－～－－－－－－－－－－－－－－SOLVE EOTTOM EDGIES

1545 FOR EY＝1 TO 8：READ R，\(S\)
\(5 S 0\) IF EX（）C（R）＊C（S）THEN NEXT EY：EY＝C（30）：GOTO בand
\(1555 E Y=8: N E X T E Y: C M=R: I F \quad C(S)=C(30)\) THEN CM＝S
\(1560 \mathrm{R}=I N T\)（EG＊EG＊EG／3－EG＊EG＊ \(1.5+E G / 6+5.1\) ）
565 IF CM＝12＊RGOTO 1590


589 GOSU 5 ，
1585 PLOT 3 ，D，WW，E，5：WW＝WW＋1：PRINT D\＄：GOSUE 245：GOSUB 115
1590 NEXT EG
59

1595 RESTORE \(320: Q=C(26) * C(29): F O R \quad I=Q\) TO \(3: R E R D R, S\)

\(1605 \mathrm{Q}=\mathrm{I}: \mathrm{I}=3: \mathrm{NEXT}\) I：FOR EG＝1 TO \(\overline{3}:\) RESTORE 3 ED
\(615 \mathrm{~F}=\mathrm{C}(E G+2 S) * C(E G+26): F D R\) EY＝1 TO 4 ：READ R，
615 IF P《C（R）＊C（S）
620 EY＝4：NEXT EY，EG
1623
1624 REM－－－－－－－－－－－－－－－－－－－－POSITION CENTER EDGIES
1625 RESTORE \(3 E D: P=C(26) * C(27): F O R I=0\) TO \(3: R E A D R, S\)
1630 IF \(P(\) ）\(C(R) * C(S)\) THEN NEXT I ：END
\(1635 R=Q * 3+I\) ：IF I \(Q\) Q THEN \(R=\tilde{R}+1\)
640 I＝3：NEXT I：IF R＝1GOTD 1655
1645 GOSUB 805
1650 PLOT 3，B，WW，6，7：WW＝WW＋1：PRINT D\＄：GOSUR 245：GOSUB 115
653
654 REM－－－－－－－－－－－－－－－－－－－－FLIP CENTER EDGIES
\(1655 \mathrm{R}=1\) ：IF C （ 8 ）（） \(\mathrm{C}(26)\) THEN \(\mathrm{R}=\mathrm{R}+4\)

1660 IF \(C\)（9）（）C（E7）THEN \(R=R+2\)
1565 IF \(C(10)\)（）C（28）THEN \(R=R+1\)
1670 IF R＝1GOTO 16Э5
1675 ON R－1 GOSUE \(865,870,875,880,885,890,905\)
1680 PLOT 3， 0 ，WW：IF LEN（D \(\ddagger\) ） 5 S THEN PRINT RIGHT \(\$\)（D \(\$, 42\)
1685 PLOT 6，7：IF LEN（D．p）（5\％THEN PRINT D
1690 GOSUB E45：GUSUB 115
\(1695 \mathrm{FOR} E G=26\) TO 29：IF C（EG）（）C（EG＋18）GOTO 2068
1700 NEXT EG
1705 PLOT 3，1a，36：PRINT＂ALL DGNE＂
1800 INPUT＂HIT RETURN WHEN READY＂；WW\＄：RUN
1803
1804 REM－－－－－－－－－－－－－－－－－SUBROUTINE TO SET CUBIE COLORS
1805 PLOT 3，4，17：PRINT＂ENTER COLORS VIA COLORED KEYS＂
1810 FOR S＝R TO 5：FOR \(T=0\) TO 2：FOR U＝0 TO \(2: Z=S+U * E+T * 18\)
1815 ON S GLTO 18E5，1830，1835，1845， 1855
\(18 E \cdot X=(T+U) * E+51: Y=11 \Xi+(U-T) * G: G O T O 1960\)
\(1825 X=T * 6+46: Y=85+\cup * 10-T * E: S O T O 1850\)
\(1835 X=T * 5+92: Y=94+U * 10+T * E\)
\(1835 \mathrm{X}=\mathrm{T}\) ．
\(1845 X=T * 5+110: Y=103+U * 12-T * 6\)
1850 GOSUB 1970：GOSUB 185：PRIN
850
8Js XOSU
1855 NEXT \(\cup, T, 5: G 0 T 01900\)

1875 FLOT E，Q，BE，2G：INPUT＂＂；WW\＄
1880 IF WW \(="\)＂
1885 WWH＝ASC \｛LEFTक（WW中，1））－16：IF WW） 7 THEN WW＝WW－64

1950 IF WW＜ 1 E THEN \(C(Z+1)=W W\) ：RETURN
1895 GOTO 1870
1900 PLOT 3，4，17：INPUT＂ANY FURTHER CHANGES（ Y／N）＂；WW\＄
1905 IF LEFTक（WWゅ，1）＝＂Y＂THEN PLOT 11，ЕS，11：GOTO 1895
1910 FOR EG＝7 TO 15：PLOT 3，R，EG：PRINT SPC（EQ）：DRINT ：NEXT EG
1915 FOR EG＝16 TO 31：PLOT 3， 0, EG， \(11:\) NEXT EG：RETURN
1998
1999 REM－－．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．
2のMO PLOT 3， 0,29, E，EG：PRINT＂CANNOT FIND＂；
2ด® GOSUF EQ70：EY＝C（ESTEG）：GOSUE 2070：PRINT＂EDGIE＂
2010 PLOT E，2：PRINT＂SOMEONE SWITCHED THE STICKERS＂
2015 EOTO 1902
2020 PLOT 3，Љ૬，29，6，56：PRINT＂SWAPPED COLORS AT＂；：GOTO 2030 2Q2S PLOT \(\overline{3}, Q_{3}, E \exists, E, E 6: F R I N T\)＂CANNDT FIND＂：：EY＝C \((3 \Omega)\)
2030 GOSUB E070：EY＝C（ES＋EG）：GOSUB EOTQ
2035 EY＝EX／C（30）／EY＋．1：GOSLIE 2Q7日：PRINT＂CORNER＂：GOTO 2010 2040 PLOT \(3,0,2 \exists, 6,65:\) PRINT＂GWAPPED COLORS AT＂；：GOTO EOSO
2045 MLOT 3，
2050 GOSU L
COS CORNER＂：GOTO EO1Q 20E 2070 ON EY－1 GOTO ED80，E085，2050，E095， 2100,2105
2075 PRINT＂RED＂：：RETURN
2Q日日 PRINT＂GFEEN＂．：RETURN
RQ85 PRINT＂YELLOW＂：：RETURN
2Q80 PRINT＂FLUE＂：RETURN
2095 PRINT＂PURPLE＂：：RETURN
210Q PRINT＂INVALID＂：RETURN
2105 PRINT＂WHITE＂；：RETURN

2198
INSTRUCTIONS
2200 PLOT 3，a，1E：PRINT＂MODES GF PLAY＂
2ᄅI PRINT ：PRINT＂SELECT MIX \＆PLAY AND I＇LL MIX THE CUBE FD R YOU TO SQLVE．＂SELECT MATCH \＆SOLVE AND YOU CAN MATCH MY DISPLAYS TO＂
2e30 PRINT Y YUR HAND HELD CUBE THEN I＇LL GIVE YOU A LIST OF MOVES＂
2240 PRINT＂WHICH WILL SOLVE THE CUEE FROM THAT POSITION OR \(L\) ET YOU＂
225 PRINT＂KNOW YOU HAVE AN INVALID CUEE：
2260 PLOT 3，12， \(30:\) INPUT＂HIT RETURN TO CONTINUE＂；WW
227 PLOT 3， \(0,16:\) PRINT＂HOW TO ENTER MOVES＂
2280 PRINT ：PRINT＂MOVES ARE ENCODED BY USING THE FIRST LETTE R DF FACE TO EE＂
2290 PRINT＂ROTATED．THIS IS FOLLOWED BY A NUMBER FROM ONE 1）TO
2300 PRINT＂＂．THREE（3）TO INDICATE THE NUMBER OF QUARTER TURNS IN A PLOT 11：PRINT ， CLDCKWISE DIRECTION AS YOU FACE THAT SIDE

2330 PRINT＂FOR EXAMPLE：F3 TURNS THE FRONT FACE THREE（3）QU ARTER＂
2340 PRINT＂MULTIPLE MOVES MAY EE WRITTEN AS FOLLOWS：＂
2350 PRINT＂MLLTIDLE MOVES MAY GE WRITIEN AS FOLLOW＂
2378 PLOT \(3,12,30:\) INPUT＂HIT RETURN TO CONTINUE＂；WW末
23AQ FOR I＝3日 TO 16 STEP－1：PLOT 3，D，I，11：NEXT I
2380 FOR I＝30 TO 16 STEP－ \(1:\) PLOT 3 ，
24RO PRINT＂A YYE＇RESPONSE TO THE DISPLAY EACH MOVE QUERY W
ILL DO＂
2410 PRINT：
R，WHEN＂
2420 PRINT＂
2420 S ＂
2430 PRINT＂
2440 PRINT THE＂
245 PRINT \(\because\) COLOR KEY PAD（OR THE COLDRED KEYS OF THE STD．KE YEOARD），＂
2455 PRINT＂ E THE＂
2460 PRINT CHANGE LOCATION．
2470 PLOT \(3,12,30:\) INPUT＂HIT RETURN TO START＂；WW\＄：RUN

\title{
Assembly Language Screen Dump To MX80 Printer
}

\author{
by Steve Reddoch \\ 1158 Via Bolzano
}

Santa Barbara, CA 93111

After typing in the BASIC language version of the screen dump program (Aug./ Sept. Colorcue), I typed RUN and waited. And waited. BASIC was just too slow, or I'm too impatient. In any case, I sat down right then and rewrote the program in Assembly Language, and after working out the bugs, it's as fast as the original was slow.

The program follows the original to the letter (almost). The BASIC statements are included as comments, and serve to guide the reader, even one unfamiliar with the 8080 instructions, through the code. The translation looks easy. It was, but it's misleading. Much use was made of the fact that many of the numbers never get greater than 127 (one byte) and there are no floating point numbers. Not all BASIC programs would lend themselves to such a simple translation. Since most of the multiply instructions are by powers of 2 , rotate instructions are used. The only tricky instruction is at line 9520, but
some study will lead to an understanding of its logic, if you're so inclined.

The program is set up to be CALLed from a BASIC program, and it occupies the top of memory. When RUN, it is brought into memory and protects itself by changing BASIC's end of memory pointer to just ahead of itself. It sets the CALL vector to point to the CALL entry point, sets the baud rate and returns to FCS, and then to the BASIC CALLer. When CALLed, the label named CALL: is given control, and the program begins to translate the screen (plot characters only) into Epson graphics characters in the standard character set, and print them line by line. Screen plot blocks are 2 dots wide and 4 dots high; Epson blocks are 2 dots wide and three dots high. When printed, the image looks a little higher than when on the screen, but that shouldn't prove objectionable. A sample BASIC test program is shown which exercises the routine.
```

9 ~ R E M
BRING IN ASSEMBLY LANGUAGE ROUTINE
10 PLOT 27,4:PRINT "RUN EPSPLT":PLOT 27,27:CLEAR 1000
1 1 REM
PUT INTO VECTOR PLOT MODE
20 PLOT 12,2,242
21 REM
PLOT 5O RANDOM VECTORS
30 FOR X= OTO 50
40 PLOT RND (2)* 127,RND (2)* 127
5 0 ~ N E X T
51 REM
PRINT SCREEN
60 Z= CALL (O)

```

ASSEmbly Language screen dump for mx-bo frinter

; BALJD RATES:
;1 - 110 EAUD
i2- 150
;3-300
: 4 - 1200
;5-2400
is - 2400
;7-9600
;
CALLED FOUTINE, ENTERED FROM BASIC BY i=CALL (0) ......CALL VALUE NOT USED
\begin{tabular}{|c|c|c|c|}
\hline \multirow[t]{10}{*}{CALL:} & PUSH & PSW & \\
\hline & PUSH & H & \\
\hline & PUSH & E & : SAVE NEEDED REGS (AND UNEEDED AS WELL) \\
\hline & LDA & LOFL & \\
\hline & FuSH & FSW & ; SAVE Lofl value \\
\hline & MVI & As, 14 & \\
\hline & STA & LAFL & ; NEW l loflag to direct all outfut to ptr \\
\hline & MVI & A, 199 & \\
\hline & OUT & 8 & (LOCK OUT KEYEOARD (LOWER CASE Y*S) \\
\hline & JMP & EEGIN & \\
\hline \multirow[t]{8}{*}{EXIT:} & POP & F.SW & \\
\hline & STA & LOFL & ;REstore lofl to frev value \\
\hline & mVI & A, 207 & \\
\hline & OUT & E & UUNLOCK THE KEYBOAFD \\
\hline & FOF & E & \\
\hline & POP & H & \\
\hline & POP & PSW & ; RESTORE REGISTERS \\
\hline & FET & & ORETURN TO EASIC FGM \\
\hline
\end{tabular}
; MAIN ROUTINE
; 9020
FOR Y \(=0\) (TO 126 STEF \(\Xi\)
EEGIN:

\(: 9030\) FLOT 14
\(\begin{array}{ll}\text { MVI } \\ \text { CALL } & \text { LO }\end{array}\)
99040 FOR \(x=0\) (TO 126 STEF 2)
XFA A
59040
STA
EQU
CALL CI
CFI 10
JZ \(\quad\) EHECK FOF INTERRUFT
\(\because 9050 \mathrm{CH}=160\)
MVI A. 160
: M=1 VCH
STA VM
MUI A,2
9060 IF \(Y=126\) THEN IM \(=1\)
LDA VY
INZ S9060E
MVI \(\quad A_{8} 1\)


;*** a table of a flot souafe ***
SQUTAB:
\begin{tabular}{ll} 
DE & 1 \\
DB & 2 \\
DB & 4 \\
DB & 8 \\
DB & 16 \\
DE & 32 \\
\(D B\) & 64 \\
DB & 128
\end{tabular}
;*** FROGRAM VARIABLES ***
\begin{tabular}{|c|c|c|c|}
\hline UY: & DS & 1 & ; VARIABLE FOR ACROSS THE SCFEEN LOOP CTL \\
\hline \(v x\) : & DS & 1 & : VARIAEIEE FOR DOWI THE SCREEN LODF CTL \\
\hline VI: & DS & 1 & ; VAR FOR VERTICAL SFOTS FEE EFSON FLIOT \\
\hline ; & & & BLOCK. § FDF ALL EUT LAST FASS, WHEN 1. \\
\hline VJ: & DS & 1 & : LOOP VARIAELE FOF HORIZ SFOTS FEFi ELDCE \\
\hline UM: & DS & 1 & FEPSON FWR OF 2 SFOT CODES. REL. TO 160 \\
\hline VIM: & DS & 1 & ; VERTICAL SPOTS LOOF LIMIT. WOULD EE 3 \\
\hline ; & & & EXCEFT LAST FASS REQUIFES EUT \\
\hline VDA: & DS & 1 & ; DATA CHAR FFICM SEREEN \\
\hline VCH: & DS & 1 & :EFSON FLOT CHARACTER \\
\hline VSY: & DS & 1 & ; CCII Y FLOT FOS. (0 TO 127) \\
\hline vsx: & DS & 1 & ;CCII X FLOT FQS (OTO127) \\
\hline UFT: & DS & 1 & ;RETURN VARIAELE FRDM "IS DOT ON?" SUBR. \\
\hline ; & & & - 1 IF NOT ON, \\
\hline ; & & & >0 IF DOT FRESENT \\
\hline
\end{tabular}

Cueties

PLOT 12,2:F=.3:X=64:Y=X:FOR J=ø TO 2øø: \(A=F * J: X=X+S I N(J) * A: Y=Y+C O S(J) * A: P L O T X, Y:\) NEXT:PLOT 255

\title{
CALL Subroutine Linkage
}

\author{
by Ben Barlow
}

Volume 1, Number 1 of Colorcue contained the first assembly language CALLable subroutine (the famous "Scrolling Patch"), and many have appeared since. Recently we have published CALLable subroutines to sort strings and handle protected fields on the screen. It's about time to explain what they are, why they are, and how they work.

First, what are they and why are they used? A CALLable subroutine, or assembly language patch, is a piece of code similar to a GOSUB subroutine, but coded in 8080 assembly (i.e., machine) language. It is connected to a BASIC program, and the BASIC program CALLs it much like it GOSUBs to BASIC subroutines. Figure 1 illustrates this type of linkage. (Some computer languages offer the ability to "drop into" assembly language from a high level language, which would bring the assembly language routine "inline", but BASIC is not one of these.

Such routines are most frequently used when speed is required or desired. The scrolling patch used in the Star Trek program, for example, had a need to manipuate a small portion of the screen as if it were a scrolling CRT terminal screen. BASIC could have done the job,

PEERing and POREing, but would have been just too slow to be effective. The sort routine (Oct/Nov Colorcue, p. 21) has been written in BASIC, but was speeded up by a factor of 50 when recoded into assembly language. The Soundware routines, on the other hand, wouldn't work if they were coded in BASIC; fast execution is needed to control the sound.

Figure 2 shows how CALLable subroutines are most often positioned in memory, stuffed away almost at the top. (Although the routines can be placed in other locations, the top of memory is the easiest and safest place. Other locations are left as an "exercise for the reader".) One obvious question is how to get the subroutine up there. The most common methods are:
(1) A series of DATA statements, with READs and POREs from BASIC, suitable for short routines, but tedious, very tedious, for something the size of the sort routine. This type of placement is illustrated by the scrolling patch, but is useful only for small routines.
(2) Using FCS to LOAD a previously assembled routine and establishing the linkages through BASIC POKEs. This method saves most of the trouble of converting

BASIC PGM

statements, but it is still a non-optimum mix of function between the BASIC and the assembly language program.
(3) Using FCS to RUN a previously assembled program which then proceeds to set up its own linkages. This is the method used in recent Colorcue articles; it keeps the BASIC program small, reduces the chance for error (and the need to do hex to decimal conversions), and is simple to program.

Linkage to the subroutine is established by putting a JMP (jump) instruction into the CALL(x) jump vector location ( 33282 or 8202 H for the hex fanatics) to


Figure 2
cause a jump to the beginning of the assembly language routine. Normally, this location contains a JMP to an error routine in BASIC. When BASIC interprets a CALL, it sets up registers as directed, and then JMPs to the CALL vector location. If the vector contains a JMP to your patch, it's entered. If not, the CALL immediately returns to the interpreter, and acts as a NOP (No OPeration) statement. There are two ways to get the JMP to your routine into the CALL vector:
(1) POKE it in from BASIC. Assuming your routine begins at location OFEODH in memory, the following code in Listing 1 will set up the vector.
(2) Store it from the assembly language routine. Given the same assumptions as above, the code would be as given in Listing 2.

Loading the routine into memory and establishing linkage to it is not the end of the inititalization phase. Without taking precautions to protect the routine, BASIC would overwrite (read destroy) it at the first opportunity. By putting the routine at the top of memory, however, it is a simple matter to "fake out" BASIC and protect our routine. This is done by changing the system's pointer to the top of usable memory, making it point just in front of the routine we wish to protect. When BASIC subsequently allocates string space, it will check the "top of memory" location (32940, or 80 ACH ) and not go beyond. Figure 3 shows what this looks like. There are two ways to set this


Figure 3
pointer (wouldn't you guess?):
(1) By executing BASIC POKEs, as in Listing 3.
(2) By executing 8080 instructions in the setup section of the routine (Listing 4).

Once the routine is loaded into memory and linked to the BASIC program, it is used by executing the CALL statement in BASIC, e.g.:
\[
60 \mathrm{z}=\operatorname{CALL}(\mathrm{F})
\]

As one might expect from an understanding of BASIC, the variable \(F\) is given to the subroutine, and the value returned by the subroutine is assigned to the variable 2 . There are restrictions on \(F\) and \(Z\); \(F\) will be converted to a two byte integer value (giving it a range from 0 to 65535), and Z will be similarly limited.

On entry to the subroutine, BASIC will have placed the value of \(F\) (after converting to an integer) into the \(8080^{\prime} \mathrm{s}\) D and E registers ( E is the low order byte). To return a value, which the interpreter will
place into \(Z\), the subroutine must exit (RET) with the value in the \(D\) and \(E\) registers. During the routine, \(H\) and \(L\) should be preserved.

Now that we understand the concepts, let's look at a small assembly language routine (Listing 5) that illustrates them. There are several interesting topics related to the ones discussed in this article:

Implied parameter passing, i.e. using BASIC variable values directly in an assembly language routine.

Linking to an assembly language routine from things other than CALLs, e.g. interrupt routines, such as User Timer no. 2 or the keyboard, the serial port, or BASIC output.

Implementing multiple functions in a single subroutine without POKEing new addresses.

These are topics for a future article, though you can find examples of some of them in past Colorcue pieces. Write your editors and let them know what you'd like to see.
\begin{tabular}{|c|}
\hline \begin{tabular}{l}
\(15 \mathrm{RT}=65037: \mathrm{REM}\) FEODH \\
16 POKE 33282,195:REM The JMP op code \\
17 POKE 33283,RT AND 255:REM The low order byte of the address \\
18 POKE 33284 ,INT (RT/256):REM The high order byte of the address \\
Listing 1.
\end{tabular} \\
\hline \begin{tabular}{l}
MVI A, 195 ;The JMP op-code. Can be (JMP) if Macro-Assembler. \\
STA 33282 ;or 8202 H \\
LXI H,CALL ; CALL is the entry point of the routine. \\
Listing 2.
\end{tabular} \\
\hline \begin{tabular}{l}
\(19 \mathrm{TM}=61439:\) REM \(\emptyset\) FDFFH, one location in front of routine \\
20 POKE 32940,TM AND 255:REM The low order byte \\
21 POKE 32941, INT(TM/256):REM The high order byte \\
Listing 3.
\end{tabular} \\
\hline LXI H, START-1 ;One byte in front of code SHLD 80ACH ;Stuff into top of mem ptr (32940) Listing 4. \\
\hline
\end{tabular}


\section*{Listing 5 .}
```

5 REM THE FIRST THING TO DO IS LOAD AND LINK TO SUGROUTINE
10 FLOT 27.4:PRINT "RUN QALP":PLOT 27,27:REM LINK UP TO AL SUER
OUTINE
15 CLEAR 100:REM MUST BE HERE, RIGHT AFTER THE RUN CMD.
25 FLOT 27,24,6,2,15
30 FOR Y= OTO ड1
35 PLOT S,O,Y
40 FOR X=OTO RND (2): 63
45 PLOT 32+ X
SO NEXT : NEXT
5 5 ~ F R I N T
tO PLOT 3,0,0,6,2,11:INFUT "ENTER CCI CODE: ";CCI
65 Z= CALL (CCI)
70 IF CCI< 25STHEN GO
80 FOR CCI= OTO 255: Z= CALL (CCI):NEXT :END

```

Listing 6.

\title{
Keyboard Reading In BASIC
}

\author{
by Steve Perrigo \\ c/o Harding Lawson Associates \\ P.O. Box 3885 \\ Bellevue, WA 98009
}

Several months ago, an article in FORUM (1) by B.F. Muldowney of Australia caught my eye. The article was titled "Advanced Keyboard Reading" and outlined a method to interrogate the keyboard for single or multiple key closures. At the same time I was looking for a simple joystick modification that could provide more than up, down, right and left. The solution was in the use of the Atari joystick and a keyboard scanning technique. An article describing keyboard scanning for the joystick appeared in the March/April issue of FORUM. An update was published in the May/June issue.

Since the publication of articles regarding joystick modification, software using joysticks has begun to appear on the market. Most notable (as of this writing) is CHOMP (2), a version of Pac-Man. As a result of the article explaining a method to install and use the joystick and keyboard scanning technique, I received several requests to explain a simplified version of how to read the keyboard in BASIC. It should be noted that most of this material is a representation of material from Muldowney's article.

To understand keyboard reading you must look at the keyboard from the computer's point of view. Figure 1 is a schematic drawing of just that. To understand this article you should rely heavily on this Figure. In addition, it is very important to note that this method only determines whether or not a particular key is physically depressed at any one time. It does not matter whether that key is in a shift, control or command status. That is why on the schematic you may not see some keyboard characters that appear on your key caps. In all cases I have listed the non-shifted, default characters represented by a key closure.

As depicted in Figure 1, the keyboard is accessed by 16 lines in, any of which may be interrogated individually, and one line out which will tell you the status of the keys along the interrogated line. The keyboard is interrogated through port 7 using the OUT 7, "line number" statement. For example, the statement OUT 7,10 will send a signal through line 10 to interrogate the closure/non-closure status of the \(5, \mathrm{E}, \mathrm{U}, \mathrm{F} 5\), AUTO and MAGENTA keys. Note that only one of the 16 lines may be interrogated at any one time.

After the keyboard is interrogated with OUT 7,"1ine number", you must then check for an answer from the keyboard. This is done by checking the value in port number 1 by means of the \(\operatorname{INP}(1)\) statement. The value is retrieved using the form \(X=\operatorname{INP}(1)\), and \(X\) will be in the range 128-255. Note that in Figure 1 each of the horizontal grid lines is assigned a value in the box on the far right. These values are \(1,2,4,8,16,32\) and 64. (Notice any relationship among those numbers?) When a key is closed along an interrogated vertical line (i.e., oUT 7,"line number"), the ultimate result is that the value 255 is decremented by the value assigned to the corresponding horizontal grid line. Now an example. Suppose you interrogate line 10 as described above with the OUT 7,10 statement. If, for example, the E and F5 keys are both closed, then the corresponding value of \(X\) that you retrieve with the \(X=\operatorname{INP}(1)\) statement will be 245. That is, 2 for the E closure and 8 for the F 5 closure for a sum of 10 , which is subtracted from 255, resulting in 245 . I'll talk more about simultaneous closures later.

Now for another example. Suppose you wanted to detect the closure of any of the keys \(2,4,6\) or 8 . How would you go about it? Answer: scan for each closure
independently. First scan line 13 for the closure of 2 , then line 11 for the closure of 4 , then line 9 for the closure of 6 , and finally line 7 for the closure of 8. Immediately after executing each OUT 7, "line number" statement, check the value of \(\operatorname{INP}(1)\) for a value of 254 ( \(=255-1\) ) indicating that the key is closed.

Prior to scanning the keyboard, you will want to lock out the keyboard (and some other built in routines) by an out 8,0 statement in your program. In this mode all the keys, including AUTO, ERASE PAGE, ATTN/BREAK, etc. may be depressed and their intended functions will be ignored by the computer. With this lock out command your program will run, but no characters from the keyboard will be echoed to the screen, thus keeping a "clean screen" during program execution. At the end of program or function execution you will want to reenable the keyboard with OUT 8,255.

Try the routine in Listing 1 to check the values of the keys depressed. Make sure you understand why it operates the way it does before you proceed.

After using this short routine you should have a feel for how this system monitors the keyboard for single key closures. Now try the routine with multiple closures and examine the results. Depress the 4 key and keep it depressed. Now add the \(D\) key, then the \(T\) key and so forth, pressing all of the keys along one of the lines into the keyboard. You will see that the same line is being interrogated and that the decremented value changes as the sum of the multiple closures changes.

Of course a slick method like this can't exist without a hitch. And here it is. There are times when the keyboard can signal the closure of a key that is not really depressed. Refer to the schematic. Suppose that you simultaneously close the H, G and W keys while you are monitoring along lines 7 and 8 for key closures. As usual, while interrogating along lin'. 8 for closures, the value of \(\operatorname{INP}(1)\) will be 249. (If you don't understand this yet, return to go and start over.) However, when you interrogate along line 7 , the value of \(\operatorname{INP}(1)\) will also be 249 , indicat-

Interrogate one of these lines with the Out \((7,-)\) Command


Figure 1. Keyboard scanning lines and decrement values.
ing that both the \(H\) and \(X\) keys are depressed. Why? Because when you interrogate the keyboard with OUT 7,7, the signal travels down line 7 and then, because the \(H\) key is depressed, the signal travels along the horizontal line which decrements the value by 2. But when it reaches the closed \(G\) key, the signal also goes down line 8 to the closed \(W\) and out along the next horizontal line, which decrements the value by 4. Simple? Not really; but it does make sense. Needless to say, when designing software using this technique, be sure to avoid this problem.

Now you should have a general idea of how the technique works. As you can see, it is a cumbersome method of getting a character from the keyboard. The real value of keyboard scanning is when you are monitoring for multiple closures, especially in game situations.

I will make a quick digression into why this system works so well for the Atari joystick modification. Basically, the Atari joystick has five switches in it. For ease of programming, the primary joystick is wired into the four main directions of the number keypad, namely, the \(2,4,6\) and 8 keys. The joystick fire button is wired to the 5 key. The joystick is designed so that an angle closure will close two keys simultaneously. For example, the diagonal vector of up and to the right corresponds to the closure of the 6 and 8 keys. With the keyboard scanning technique we can detect the double closure and with a short routine interpret that
intended direction. Listing 2 is a short routine to do that. It returns a value of 0 (no key closure) to 9 , corresponding to the positions of the numbers on the number key pad. If you come up with a more efficient routine in BASIC to do this, let me know.

By now you have probably noticed that the system as described will not detect the closure of certain keys, namely, the SHIFT, CONTROL, COMMAND and REPEAT keys. (The only key closure that cannot be detected is the CPU RESET because it has a direct line to the processor.) Those key closures and several others are detectable using the following similar method. Use OUT 7,128 and interrogate INP(1) for a value of \(127-\) SUM (instead of \(255-\) SUM as before). The decrement values are given in the Table. You will notice that the COMMAND key is not listed. Remember that it is actually the simultaneous closure of the SHIFT and CONTROL keys. -
\begin{tabular}{|clc|}
\hline Key & & Decrement \\
\hline\(/\) & (divide) & 1 \\
0 & (letter 0) & 2 \\
\(\bar{F} 15\) & (Function key) & 4 \\
Control & 16 \\
Shift & 32 \\
Repeat & 64 \\
Table of Keys and Decrement Values \\
Used with the OUT 7,128 Command \\
\hline
\end{tabular}

Table of Keys and Decrement Values Used with the OUT 7,128 Command
\begin{tabular}{|c|c|c|}
\hline 10 & REM & KEYBOARD SCANNING PROGRAM \\
\hline 20 & PLOT 29,27,11,6,3,12 & \\
\hline 30 & OUT 8,0: REM & DISABLE KEYBOARD \\
\hline 40 & FOR \(\mathrm{I}=0\) TO 15 & \\
\hline 50 & OUT 7,I: REM & SCAN EACH LINE INTO KEYP \\
\hline 60 & \(\mathrm{X}=\mathrm{INP}(1)\) : REM & GET A VALUE FROM THE KEYBO\&RD \\
\hline 78 & IF \(X<>255\) GOTO 190: REM & CHECK FOR ANY KEY CloSURE \\
\hline 80 & NEXT I & \\
\hline 90 & GOTO 40 & \\
\hline 100 & IF \(X=X 1\) AND \(I=I 1\) GOTO 80 & \\
\hline 110 & \(\mathrm{XI}=\mathrm{X}: \mathrm{Il}=\mathrm{I}\) & \\
\hline 120 & PRINT I, 255-X, X: REM & PRINT CLOSURE VALUES \\
\hline 130 & IF \(X=251\) AND \(I=15\) GOTO 15 & B: REM QUIT IF "Q* \\
\hline 140 & GOTO 86 & \\
\hline 150 & OUT 8,255; END: REM & REENABLE INTERRUPTS \& END \\
\hline
\end{tabular}

Listing 1. Keyboard Scanning Routine.
```

    10 OUT 8,0
    20 \(\mathrm{T}=0\)
    30 OUT 7,13: IF INP (1) \(=254\) THEN \(T=T-3:\) REM DOWN (2)
    40 OUT 7,11: IF INP \((1)=254\) THEN \(T=T+5\) : REM LEFT (4)
    50 OUT 7,9: IF INP(1)=254 THEN T=T+7: REM RIGHT (6)
    60 OUT 7.7: IF INP(1) \(=254\) THEN \(T=T+9:\) REM UP (8)
    70 IF \(T>9\) THEN \(T=T-6\)
    \(80 \mathrm{~T}=\mathrm{ABS}(\mathrm{T})\)
    98 IF \(T=0\) THEN \(T=1\)
    10日 OUT 7,10: IF \(\operatorname{INP}(1)=254\) THEN \(T=6:\) REM FIRE (5)
    $110 \mathrm{~T}=\mathrm{T}-1$
120 OUT 8,255
130 RETURN

```

Listing 2. Scanning the Number Keypad.

\section*{REFEREHCES}
(1) B.F. Muldowney, "Advanced Keyboard Reading", FORDM International, Vol. I, No. 5 (Nov/Dec, 1981), pp. 64-65. And see S . Perrigo, "Joysticks Standard for the Compucolor", FORDM, Vol. II, No. 2 (May/June, 1982). (FORTMM is the publication of the Canadian Compucolor Users' Group. c/o Editor, 21 Dersingham Cres., Thornhill, Ontario, Canada L3T 4P5. I highly recommend that any serious user join this group.)
(2) "CHOMP", a highly recommended program for the joystick, available from ICS, 12117 Comanche Trail, Huntsville, Alabama 35803, and other CCII software dealers for \$29.95.

\title{
Assembly Language Programming
}

\author{
by \\ David B. Suits
}

\section*{PART VI: The Input Routine}

In this installment we'll finish up the input routine begun last time, and we' 11 be taking another look at using the assembler.

The program in the listing is my version of the input routine, fancied up a bit just for purposes of demonstration. It not only accepts input (both upper and lower case, contrary to my proposal in the last issue), but, just for fun, reprints the input in a different color. This verifies that the input string did indeed get stored in memory. In case you still have doubts, after you've assembled and run the program, you can use the Machine Language Debug Package to look at what's in the input buffer (starting at 822 EH , if you use my program). After the input is reprinted, the program loops back and asks for another line of input. It will loop forever this way, and you'll have to push CPU RESET to get out of it.

The two columns of numbers at the left of the listing are the hexadecimal numbers which are the machine language code which the assembler generated from my assembly language text. (The rest of the material is the source code which I typed in using the Screen Editor.) The column at the far left is the address (the program starts at 8200 H ), and the next group of numbers indicate the contents of the memory locations starting at that address. For example, at 8202 H the 3 EC 3 H represent two bytes. The first byte, 3 EH , is the code for the MVI A instruction, and 0C3H is the byte of immediate data to be moved into A. That is, it is a translation of the assembly language line MVI A,0C3H.
(Remember that the assembler takes all numbers as decimal unless followed by \(H\) for hexadecimal or \(B\) for binary or 0 or Q for (ugh!) octal. Moreover, all numieers must begin with a digit; that's why JC3H must be written the way it is.)

Somewhere before the first statement in the program, I must tell the assembler where the program is to reside in memory. The ORG 8200 H tells it that.

Notice that I've probably gone overboard in my use of EQUate statements. For example, instead of using the number 13 throughout as the code for a carriage return, \(I^{\prime}\) ve used \(C R\) and \(I^{\prime} v e\) told the assembler that \(C R\) is to be EQUated with 0 DH (=13 decimal). I've used hexadecimal everywhere, merely for the sake of adventure.

After the EQUates comes the label INBUFF. You can tell it's a label of an address because it has a colon after it. The DS means Define Storage, and the number to the right of that indicates how many bytes, namely, LIMIT+1. Since LIMIT is EQUated to 40 H ( \(=65\) decimal), the assembler knows to reserve, or skip over, 41 H bytes and to assign the label INBUFF to the location of the first byte (822EH). When characters are read from the keyboard, they will go here. Next come two bytes reserved for holding the address of (i.e. pointing to) the next available location for a new character in INBUFF. The program really doesn't make any use of that pointer (although it looks like it is going to at \(828 \mathrm{DH}-8290 \mathrm{H}\) ), but if you decide to add a lot of bells and whistles to the program, you might have need for such a thing. Anyway, I put it in just in case. (Or was it to confuse you? I can't remember now.) As it stands,
the instruction at 8290 H is not necessary; the contents of HL always keep track of the next spot in the buffer, and so HL is used as the buffer pointer.

Next come the labels SETUP, PROMPT and REPRNT. These define addresses at which a string of Defined Bytes (sort of like BASIC's DATA) can be found for use with the OSTR routine. If you translate the hex numbers on the left into ASCII, you'11 see that the assembler has translated YELLOW into 13 H ( \(=19\) decimal), the \(T\) in TYPE SOMETHING into 54 H ( \(=84\) decimal), and so on. In BASIC if you PLOTted these numbers (in decimal) you would get the words "TYPE SOMETHING:" in yellow. But you knew all that from our previous work, right?

The very last line of my program is END TEST (followed by a carriage return). The END signifies the end. (Natch!) The TEST tells the assembler that the starting address for the program is at the label TEST. Didn't it know that already? Well, yes; that is, it assumed it. But while the ORG statement told the assembler where the program is to be loaded, the TEST at the end tells it where to begin execution. It is assumed that the two are the same unless told otherwise, so technically the TEST is not necessary. But if I had written END INPUT, then the program would load at 8200 H , but execution would start at 828 DH . But I (almost) always make the start address the same as the load address; you will too. For now, anyway.

\section*{Dealing with the Assembler}

It would be a tax on anyone's patience to have to plunk in the TEST program one byte at a time. Even using the MLDP would not lessen the burden very much. So it's time to use the assembler.

I use ISC's assembler (not the Macro Assembler, please), but if you're using someone else's assembler, the same, or similar, procedures will probably apply. First, write your program using a suitable text editor. (The Screen Editor is perfect. If you have the old Text Editor, throw it away and buy the Screen Editor.) Just type it right in and save it on disk. Now run the assembler. It will sit there blinking stupidly at you. You have several choices. You can quit and go watch Rockford Files re-suns. Or you can put your text disk in the drive and tell the assembler to assemble your program:

\section*{>ASM TEST}

Or you can have the assembler assemble the program and save the result on disk:

\section*{>ASM TEST TO CDO:}
(Depending on your disk drive type, you will use CDO: or MDO: or whatever you call your drives. Actually, just the drive number and a colon should be necessary.) If you have two drives, you can take the text off one and save the assembled program onto another. For example:

\section*{>ASM 1:TEST TO 0:}

When the assembler is assembling your program, the assembled listing (as in Listing 1) will flash by on the screen. Be prepared to hit the ATTN/BREAK key to have it pause, especially if you see a line in red flash by, because that will mean that the assembler has picked out an error and you will want to re-edit the source file accordingly and then give it to the assembler for another try. If you want only the errors, if any, to be displayed, then add /E when you tell the assembler what to do. For example:

\section*{>ASM TEST/E}

Usually I have the assembler just assemble the program without saving the result, so that I can be assured that the program is error free. When \(I\) am satisfied, \(I\) run the assembler again and have the object file (i.e., the assembled code) written to disk (and printed on the printer at the same time).

\section*{. LDA and .PRG Files}

The result which is saved on disk is called a load file (.LDA). It is not quite the final, machine language program which you can run. The . LDA file can be loaded into memory with

\section*{FCS \(>\) LOA TEST}
(The .LDA extension is not necessary, since it is the default type for the LOAD command.) Now you can run the MLDP and test the program. When you're confident that it is the way you want it, leave the MLDP, return to FCS and type:

FCS>SAV TEST.PRG <;version, if you wisn> startaddr-endaddr

For my program, I would type:
FCS>SAV TEST.PRG 8200-82FE
The result is a copy of the memory contents from 8200 to 82 FE and is a machine language program which can be run just like any other . PRG file.

Try my input routine. Then try adding variations. (Have the input string reprinted backwards, for example.) Fiddle around. We've come a long way since last year. There are many tricks still to be learned, but you ought by now to have a fundamental grasp of 8080 assembly language programm ming. What you need now is practice. Why not try converting some simple BASIC programs into A.L.? That's often an easy
task. Or how about a machine language patch to a BASIC program? (See Ben's article in this issue.)

In the issues ahead we'll explore tables and vectors, animation, and simple number crunching. And we'll start looking at some of the more esoteric features of the routines in your machine's ROM: disk file handling, talking to your printer, and so on. In the meantime, if you've specific areas of interest you'd like to see covered, please drop me a note.

See you next issue! (Unless your subscription expires.)

\begin{tabular}{|c|c|}
\hline 2080 & 3EC3 \\
\hline 8202 & 32 \\
\hline 05 & \\
\hline 2 B & 2EC6 \\
\hline Beas & 3E \\
\hline \(1{ }^{\text {a }}\) & \\
\hline & \\
\hline & \\
\hline
\end{tabular}


\begin{tabular}{|c|c|c|c|c|}
\hline ( \(Q Q D D\) & CR & EQU & QDH & - Car^iage return. \\
\hline (QQAA) & LF & EQU & OAH & ; Lirie feed. \\
\hline (281A) & LFTARO & EQU & 1 AH & ; Left arrow for backspace. \\
\hline (00088) & ERSLIN & EQU & QEH & RASE LINE Ke \\
\hline
\end{tabular}


8こAE Dこ958こ
\begin{tabular}{|c|c|}
\hline 828 & 47 \\
\hline B2H？ & 79 \\
\hline 8 883 & FE40 \\
\hline 82 ES & CAGE82 \\
\hline 8288 & 70 \\
\hline 9283 & 23 \\
\hline 8こВ & （2C \\
\hline 82BE & 78 \\
\hline 8евС & CD923 \\
\hline 82нF & C395 \\
\hline
\end{tabular}

Bece CDD4Be
8こC5 C3958こ

BECB C2CB8：
B2CE C39582

82D 1 36EF 82D3 59

JNC INPUT1
；Nc．Iqnore the character．
；The input character（still in A）is valid．But
；see if there is still racm in the buffer for it．
MOV B，A Save character temporarily
MOV A，C \(\quad\) Get courit cif characters．
CPI LIMIT ；Has it reached its limit？
JZ INPUTI ；Yes，so igriare the character．
MOV M，B ；Put character irito buffer．
\(\begin{array}{ll}\operatorname{MOV} \mathrm{M}, \mathrm{B} & \text { ；Put character irito b } \\ \text { INX } H\end{array} \quad\) ；Bump buffer paipter．
INR C I Incremerit courit of characters．
MOV \(A, B \quad\) Get character again for echoirg．
CALL LO ；Echg character on screen．
JMP INPUT 1 ；Back for more．
；Come here upan firdirig a left arrow．
INPUTE：CALL ECKSP ；Delete previous character．
IMP INPUT1 ；That was easy！
；Come here with ERASE LINE．

；Come here with carriage return．
INPUT4：MVI M，EOLINE ；Put end af lirie byte irito buffer．
RET ；Returr tg calling rcutire．
；BCKSP－－Subroutire to delete the previous character（if
there is ore）．
ENTRY： （C）\(=\) preserit courit of characters．
（HL）- s spat in buffer of previous char +1 ．
EXIT：〈C〉＝new courit of characters．
（HL）adjusted accordingly．
（Z）if buffer empty．
（NZ）if buffer rigit emoty．
\(\begin{array}{ll}\text { 日2D4 } & 79 \\ \text { BEDS } & \text { R7 } \\ \text { 日2D6 } & \text { CAEABE }\end{array}\)

8こD 3 3E1A
B2DR CD9ころ3
BこDE 3EEQ
82E0 CDGころ3
GEE3 3E1A
8こE CD9233
日2E日 2ロ
E2E9（10
82EA CF
\begin{tabular}{|c|c|}
\hline & \\
\hline 82EB & AF \\
\hline 8こEC & 3EFE81 \\
\hline 8こEF & 3AFE日1 \\
\hline BEFき & \\
\hline B2F3 & CAEFB2 \\
\hline g2F6 & C9 \\
\hline
\end{tabular}

；Keyboard character interrupt routine．

－ERRDRS

\title{
Classified Colorcue Index
}

\author{
by James A. Kavanagh \\ Gnostech, Inc. \\ 222 S. Easton Road, Suite 15 \\ Glenside, PA 19038
}

There have been two Colorcue indices published: the first appeared in the January, 1980 issue, and the second in the June, 1980 issue. They are arranged by article names only and not by topic or subject.

Th1s index includes all issues of Colorcue, including the present issue. It is arranged by subject, although author names are included if the name appeared in two or more articles.

I compiled this index for my own convenience, and so I cannot guarantee its accuracy or completeness. It is intended more for the intermediate programmer and less for the beginner and the non-programmer. There is little reference to elementary topics or to application programs such as games.

The format used is \(\mathbf{\nabla} . \mathbf{I} \mathbf{P}\), where \(\mathbf{V}\) is the volume number, \(I\) is the issue number, and \(P\) is the page number. The December 1980/January 1981 Colorcue appeared without Volume/Issue numbers. For the purposes of this index, that issue will bear the designation III. 7.
```

Appending (programs) II.3-2, II.5-6
Apple II.5-13, II.5-14
Arrays II.6-6, III.4-7
ASCII codes III.4-14
Assembly Language
Add A to HL (ADHLA) III.6-12
AND DE to HL (ANDHD) III.6-12
ASCII to binary III.3-16
Binary to ASCII III.3-16
Binary to 1 hex nibble (B2HEX) III.4-20
Binary to 2 hex chars (LBYT) III.4-19
Block move (MOVDH) III.4-20
Carriage return/line feed (CRLF) III.4-19
Communication III.6-4
Compare, dble precision (CMPHD) III.4-19
Compatibility 6.78 vs 8.79 III.1-22
Cursor movement III.2-19
Debug IV.2-3, IV.2-6
Divide DE by HL (DIVHD) III.3-16, III.6-12
Error codes II.4-2

```
```

    Exclusive OR DE with HL (XORHD) III.6-12
    FCS from assembler III.5-13
    Graphics (cursor movement) III.2-19
    IBM I/O terminal (USC 1035) III.5-15
    Interfacing with Teletype III.1-7
    I/O II.7-14, II.8-9, III.1-18 III.2-19,
        III.3-16, III.3-18, III.3-26, III.4-18,
        III.4-26, III.5-13, III.7-3, IV.3-5,
        IV.5-19, IV.6-00
    Keyboard input III.7-3, IV.5-19, IV.6-00, IV.6-20
    Multiply DE by HL (MULHD) III.6-12
    NOR (l's complement) HL (NOTH) III.6-12
    OR DE with HL (ORHD) III.6-12
    Printer program III.3-18, III.4-26
    Programming II.7-14, IV.1-17, IV.2-6, IV.3-19,
            IV.4-19, IV.5-19, IV.6~20
    Protected fields (input) IV.3-5
    Shift DE left HL times (SHLHD) III.6-12
    Shift DE right HL times (SHRHD) III.6-12
    Sort IV.2-21
    Subtraction, dble precision (SUBHD) III.4-20
    2's complement HL (NEGH) III.6-12
    Otility routines III.4-18, III.6-12
    Wait routines
        20 milliseconds (WATL) III.4-19
        0.5 milliseconds (WATS) III.4-19
    Barlow, Ben \& IV.2-4, IV.3-13
BASIC
Chaining (Menu) II.l-6
Changing directory name II.5-3
Code line format III.1-16
Editing (FREDI) II.5-8, II.8-3, II.8-4
Editor ('THE') IV.3-25, IV.6-3
Fixing sequence numbers III.1-16
Format of code line III.1-16
Formatting numeric fields I.7-18, II.8-6,
III.4-8
Input flag, input table I.3-2
Reyboard input II.2-3, IV.6-17
Random files (see Random files)
ReORGing FREDI IV.5-11
Sequence numbers III.1-16
String manipulation I.3-4
Structure of code line III.1-16
Tokens II.4-4, III.1-16, III.4-16
Variable listings III.3-24
Bell III.7-16
Binary to ASCII III.3-16
Book reviews III.7-18
Break (generation) III.3-12, III.6-21
CALL function III.7-9
Chaining (Menu) II.1-6
Characters
Large special II.8-7
String manipulation I.3-4

```

Clarke, de France III.1-27, III.2-9, III. 2-24
Clock, real time I.2-2
Color
Codes II.3-11, III.2-16
Misc. II.3-3, III.2-18, III.3-25, III.6-19
Printer, ink jet III.7-16
'Comments \& Corrections', etc. I.3-6, II.2-6, II.3-3, II.4-9, II.5-6, II.7-17 II. 7-19, II.8-15, III. 3-26, III.6-22, III. \(7-13\), IV. 2-3, IV. \(2-20\), IV. 4-3

Communication III.1-14, III.1-18, III.3-26, III.6-4
Community access bulletin board III.5-22
Compatibility: 6.78 vs 8.79 III.1-22
Comp-U-Writer III.7-15, IV.4-15
COPY, FCS command III.5-21
COPY (screen display) II.6-13
Cross reference 1 isting, V6.78 III.7-6
Curnin, Peter III.3-3, III.4-3
Cursor movement III.2-19, III.5-4
Costom character sets III.4-12
Date routine II.7-17
Debug IV.2-3, IV.2-6, IV.4-23
Devlin, Tom IV.1-13, IV.5-5
Directory
Changing name II.5-3
Format II.5-11, III.5-9
Printout I.8-15
Supplemental (non-FCS) III.2-4
Disk dup II.1-4
Disk recovery III.6-20
Dotted 1ines II.1-2
Dup II.1-4
Dust cover II.4-2
Editing
Basic, FREDI II.5-8, I.8-3 II.8-4, IV.5-11
Basic, 'THE' IV.3-25, IV.5-9, IV.6-3
Screen, text III.1-13
Epson Printer IV.1-14, IV.2-4, IV.2-15, IV.3-13
Factoring numbers II.5-5
FCC (on Compucolor) IV.1-3
FCS
Copy III.5-21
From assembler III.5-13
Input flag, input table I.3-2
Printing directory II.8-15
(See also Directory)
Files
Creating data files II. \(2-4\)
Random (see Random files)
Recover Erum errors 1IT.6-20
Transferring from other computers III.6-4
WDA and ERE between disks YuT. \(5-21\)
 IIL.4-8
Fortana III. 7 m 17
EmDOT (BASIC Editot) \(\quad\) U, IT, \(8-3\), II, \(8-4\), IW. \(\mathrm{F}_{1} 1\)
GLine II.7-3
Graphica
8as graphs III.7-13
Circular piots I. 2-3
CRT Mode plotting IV.4-17
Cursor movement III.2-19, III.5-4
Dotted lines II.l-2
Re-entrant plot submodes III. 2-11
Scaling III.7-13
Sphere IV.2-15
3-D III.2-6, IV. 2-15, IV.4-7
Green, Bill II.4-2, II.8-3, III.1-4, III.1-6, III. 4-23, III.5-15, III.6-9, III.6-22

Handshake III. 2-26, III.3-26, III.5-21, IV, 1-10
Hardware mods III.2-26, III.3-12, III.3-26, III.4-21, III. 5-21, III.7-16, IV.1-5, IV.1-13, IV.3-13, IV.5-5, IV. 5-13

Hogan, Brian III.1-26, III.2-7
Hudson, Tom II.5-14, III.4-7, III.5-3

IBM I/O terminal (USC 1035) III.5-15
Index (Compucolor) II.3-12
Input flag, input table I.3-2
I/0 I. 2-6, I. 2-7, I.3-2, II.2-3, II. 7-14, II.8-9, III. 1-6, III.1-18, III.2-19, III.2-24, III. 2-25, III.3-12, III.3-16, III.3-18, III. 3-26, III. 4-18, III.4-26, III.5-21, III.6-4, IV.1-5, IV.1-14, IV.3-5

I/O controller (TMS 5501) III.1-6
Interface, serial to parallel IV.3-16, IV.4.3
Interfacing with Teletype III.1-6
Interrupts III.2-25
Reyboard
Input 1.2-7, III.2-24, III.4-11, III.7-3, IV.5-19, IV.6-17, IV.6-20

Lockout II.2-3, III.2-25
Reywords (see Tokens)
Light pen III.3-14
Line length II.5-15
Linked lists II.3-5, II.4-5
Lisbajous figures IV.2-18
Literature II.3-10, III.2-11, III.4-13, III.6-7

Manazir, Rıchard III.2-19, III.4-17, III.5-21
Map, system memory III.1-21
Martin, Dennis II.3-3, III.1-16, III.4-15

\section*{III.5-9}

Matzger, Alan IV.2-21, IV.5-17
Memory map III.1-21, III.7-6, III. \(7-10\)
Menu (chaining) II.1-6
MX-80 (see Epson)
Networking III.1-14
Newcombe, F. Lee II.2-4, II.8-6
Noise, power line III.5-16
Numeric base conversion III.5-11
Numeric field formatting II.8-6, III.4-8
OP code table III.6-9
Pascal triangle III.2-7
Personal software II.4-7, II.7-12
Photographing screen II.6-2
Plot table IV.2~17
Power line noise III.S-16
Printer (color ink jet) III.7-16
Printer (Epson) (see Epson)
Printer interfacing III.2-25, IV.1-5
Printer program IIL.3-18, III.4-26
Printer, screen to MX80 IV.1-14
Priating directory II.8-15
Publishing/selling programs I.2-5, II.6-5
Raffe, Bernie IV.3-5
Eam, edd on II.4-9, IV.5-5
Random files II.6-9, II. 7-12, XV.5-17
Ramdog muber III. 1-26
Reti time clock I. 2-2
Roseu, Howard IV, 1-27, IV. 2-27, IV. Row \(^{2} 15\)
ES232 I. 2-6, II.5-7, III.1-6, LII. 1.18. III.3-26. III.S-15

Screen character position IIT. \(2-13\)
Screen display copy II.5-13
Screen editor III.l-13
Screen save III. 5-3
Screen to MX80 IV.1-14, IV.2-15, IV.6-9
Scrolling patch I.1-2
Selling/publishing programs I. 2-5, II.6-5, III.7-3
Serial port (see RS232)
Sbank, Bill II.7-19, III.1-9
Smith, Bob V. IV.2-17, IV.4-17
Software (for Compucolor) I.8-14, II.4-7, II.5-7,
(see also II.6-4, II.7-6, III.1-13,

Personal software)
III.1-27, III.2-14,
III.5-18, IV.4-3, IV.4-18

Sort routine (CALLable) IV.2-21
Sound board II.5-7
Source network III.1-14
Space saving in arrays III.4-7
Steffy, Myron II.5-13, III.3-24, III.7-9
String manipulation I.3-4
```

Stroop phenomenon II.2-2
Suits, David II.2-6, III.2-11, III.3-25, III.6-19
IV.1-3, IV.1-19, IV.2-3, IV.2--6,
IV.3-19, IV.4-19, IV.5-19, IV.6-20
Taylor, Denise III.2-25, III.3-25
Taylor, Trevor III.2-25, III.3-12, III.3-25,
III.4-21, III.6-4, IV.2-18
Teletype (interfacing) III,1-6
Texr editor III.1-13
'THE' Editor (see Editing, BASIC)
TMS 5501 (see I/O controller)
Tokens (BASIC) II.4-4, III.1-16, III.4-15
TRS-80 II.5-13, II.5-14, III.2-23, III.6-4
Ungerman, Mike III.2-23, III.6-21

```

\section*{Cueties}

PLOT 12:FOR Z=Ø TO l:FOR X=ø TO 255:A=28672+X*4+ 128*INT (X/32) +Z*2ø48: POKE A,X: POKE A+l, 128*Z+2: NEXT:NEXT

\section*{** About Your Subscription **}

Most of our subscribers' subscriptions will come up for renewal after July. (Check your mailing label: the number indicates your last issue number. The June/July issue is issue number 6.)

We need your subscriptions to continue publication.
Since its beginnings in 1978, Colorcue has been financed by Intelligent Systems Corporation. That financial assistance will terminate with the June/July issue. Whether we will be able to continue publishing Colorcue will depend on whether you renew your subscription. There is a critical number of subscribers, below which Colorcue will not have the funds to continue. Let us know of your desire to see Colorcue flourish. Send us your renewal now so that we may know in advance where we stand and so that you can help guarantee the continued publication of what we believe to be an outstanding magazine for Intecolor/Compucolor users.

Subscription for one year (six issues) is \$l2 in U.S., Canada and Mexico; \(\$ 24\) elsewhere. Please make check or money order in U.S. funds payable to "Colorcue".

At the same time, why not take this opportunity to let us know what kinds of information you would like to see in colorcue during the coming year? would you prefer more hardware oriented articles? Tutorials? Programs? Applications? Games? Perhaps you have something specific in mind.

\section*{Colorcue}

Editorial Offices
161 Brookside Dr.
Rochester, NY 14618

\section*{Address Correction Requested}```


[^0]:    10070 519
    100 REA [10] [10] PTETD HANDLING SUBROUTINE[19] [10]
    $185 \mathrm{REM}[22] \mathrm{X}=\mathrm{X}$ CO-CRDNNALE [18]
    110 REM [22]Y Y Y CO-CRDINATE[18]
    115 REM [22]ML $=$ MAXTMMM INPTT LENGTH [18]
    129 REA [22]NAS= 'N' FOR MANERIC OR 'A' FOR RLPGANOMERIC[18]
    125 REM [22] KB = KEYBCARD CHARACTER [18]
    139 REM [22]PISMPINAL SIRNNG[18]
    135 RESH [22]SA = SCREEN ADDPRSS OF FTELD[18]
    149 POZPRER (32946) + PEER ( 32941 ) *256+1:REM EOINT TO PARAMETER PASS AREA
    145 SAm $28672+128 \star Y+X+X$ : REM SCREEN ADORESS
    $150 \mathrm{ZZ}=\mathrm{INT}(\mathrm{SA} / 256): \mathrm{POKE} \mathrm{PO}+3$, $\mathrm{SA}-256$ ²Z $\mathrm{ZOOKE} \mathrm{PO}+4, \mathrm{ZZ}$
    155 SAMSA-2
    169 PCKE PO+1, ASC(NAS):REM ALHHANUMERIC OR NOMERIC
    165 FIS" ${ }^{\text {" }}$
    170 PLOT $3, X, Y$
    $175 \mathrm{~KB}-\mathrm{CALL}$ (ML)
    180 IF RB-9 OR PEEX (PO) $=0$ THEN RETUKN
    185 FOR I=2 TO PEER (PO) ${ }^{2} 2$ STEP 2
    199 FIS=PIS+CARS (PEER (SA+I) ) :NEXTI I
    199 REITHR
    509 REM [10] [10] PROCRAM INTTIALITAATION[10] [10]
    529 IIRRR 190
    538 हLOT 27,4
    $54 B^{5 R}$ RRIT TRON PTELDS"
    545 REM [19]SET UP SCPERN[10]
    559 PLOT $27,27,27,24$
    560 RLOT $12,3,5,1,14$
     ${ }^{81}{ }^{\text {® }}$
    
    
    606 TP ASO "I"
    [18]*
    620 FOR $I=1$ 20 20:PLOT 10,11:NEXT I
    630 PLOT 3,1,10:PRINT "[23]NAME[18]"
    659

